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A channel-based framework for steering, non-locality and beyond

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Abstract

Non-locality and steering are both non-classical phenomena witnessed in nature as a result of quantum entanglement. It is now well-established that one can study non-locality independently of the formalism of quantum mechanics, in the so-called device-independent framework. With regards to steering, although one cannot study it completely independently of the quantum formalism, 'post-quantum steering' has been described, which is steering that cannot be reproduced by measurements on entangled states but does not lead to superluminal signalling. In this work we present a framework based on the study of quantum channels in which one can study steering (and non-locality) in quantum theory and beyond. In this framework, we show that kinds of steering, whether quantum or post-quantum, are directly related to particular families of quantum channels that have been previously introduced by Beckman *et al* (2001 *Phys. Rev. A* **64** 052309). Utilizing this connection we also demonstrate new analytical examples of post-quantum steering, give a quantum channel interpretation of almost quantum non-locality and steering, easily recover and generalize the celebrated Gisin–Hughston–Jozsa–Wootters theorem, and initiate the study of post-quantum Buscemi non-locality and non-classical teleportation. In this way, we see post-quantum non-locality and steering as just two aspects of a more general phenomenon.

Entanglement is one of the most striking non-classical features of quantum mechanics. Given appropriately chosen measurements certain, but not all, entangled states can exhibit a violation of local realism (local causality), called 'non-locality' [1]. Apart from its fundamental interest, non-locality has also turned into a key resource for certain information-theoretic tasks, such as key distribution [2] or certified quantum randomness generation [3], and has been witnessed experimentally in a loophole-free manner [4–6].

The non-classical implications of entanglement also manifest as a phenomenon called 'Einstein–Podolsky–Rosen steering', henceforth referred to as solely 'steering'. There, one party, Alice, by performing appropriately chosen measurements on one half of an entangled state, remotely 'steers' the states held by a distant party, Bob, in a way which has no local explanation [7]. A modern approach to steering describes it as a way to certify entanglement in cryptographic situations where some devices in the protocol are not characterized [8]. Steering hence allows for a 'one-sided device-independent' implementation of several information-theoretic tasks, such as quantum key distribution [9], randomness certification [10, 11], measurement incompatibility certification [12–14], and self-testing of quantum states [15, 16].

Even though these phenomena arise naturally within quantum mechanics, they are not restricted to it. Non-local correlations and steering beyond what quantum theory allows are conceivable while still complying with natural physical assumptions, such as relativistic causality [17, 18]. By 'post-quantum' we mean non-locality or steering that cannot be realized with local measurements made on an entangled quantum state[4]. Post-quantum non-locality has been vastly explored, especially its implications in information-theoretic tasks

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[4] We do not mean post-quantum in the sense of post-quantum cryptography, where one designs cryptographic protocols that cannot be efficiently broken by quantum computers. Post-quantum in our sense could refer to non-locality and steering in generalized probabilistic theories.

[19]. Little is known, however, about post-quantum steering, mainly due to the lack of a clear formalism for studying this phenomenon beyond quantum theory.

It may be unclear why one would be interested in steering in theories beyond quantum theory, since it is a phenomenon that is defined within the quantum formalism. Indeed, if we are testing quantum theory against all possible, sensible classical descriptions of reality, a local hidden variable is the most general starting point. One may however ask in which sensible ways nature may differ from a world described by quantum theory. Here we argue that it makes sense to consider the picture where locally in our own laboratory everything is described according to quantum theory, however, the global process governing the interactions between laboratories is not, analogous to the study of indefinite causal order in [20]. The existence of post-quantum steering demonstrates that the global theory can deviate from quantum theory in intriguing ways, even if our own laboratory is restricted to quantum theory. In fact, because of this, we would argue that post-quantum steering is of more foundational interest than local hidden state (LHS) models. We also note that in quantum information, bounding the set of quantum assemblages from the post-quantum set has also been studied in the guise of extended non-local games by Johnston *et al* [21].

To rectify the lack of a clear formalism for post-quantum steering, we present a framework to study both non-locality and steering complying with the no-signalling principle. Our formalism is based on quantum channels, i.e. completely-positive trace-preserving maps on density matrices. More specifically, we consider channels on multipartite systems that satisfy a form of the no-signalling principle, introduced first by Beckman *et al* [22] in bipartite setups. Indeed, they defined two families of channels. On the one hand, 'causal channels', that do not permit superluminal quantum (and classical) communication between two parties. On the other, 'localizable channels', that can be described by parties sharing a quantum (entangled) state and performing local operations with respect to each party. Furthermore, the set of localizable channels is a strict subset of the causal channels [22].

In this work, a given conditional probability distribution (correlations) in a non-locality scenario or a set of conditional quantum states (assemblage) in a steering scenario, is associated to a causal channel, and vice versa. We identify the nature of the correlations, or assemblages, with the properties of the channels that may give rise to them. In particular, if correlations or assemblages are post-quantum then they can be associated with a causal, but not localizable, channel. Utilizing this connection we derive results in both the study of quantum channels and steering.

We also show that our framework is not limited to the study of non-locality and steering. We show that non-locality studied from the perspective of channels can be expanded to other kinds of non-locality studied in the literature. In particular, Buscemi introduced the scenario of the semi-quantum non-local games [23], in which we can demonstrate a form of non-locality, denoted as 'Buscemi non-locality'. Buscemi showed that an entangled state can be used as a resource for demonstrating this form of non-locality. Here, we expand upon this original work to introduce post-quantum Buscemi non-locality, and show how it can be understood through quantum channels. Finally, we consider the analogue of steering for Buscemi non-locality, which is the study of non-classical teleportation, as initiated by Cavalcanti *et al* [24].

Summary of results

This manuscript presents a variety of results which, to guide its more comprehensive reading, we now briefly outline.

First, in the study of quantum channels, we define a novel class of quantum channels called the 'almost localizable channels' in definition 1, which are a generalization of the set of localizable channels in [22]. We show in theorem 15 that the set of almost quantum assemblages (as defined in [18]) result from almost localizable channels, and almost localizable channels only give rise to *almost quantum correlations*[5] [25] or assemblages. This is the first time that almost quantum assemblages are given a physical definition, rather than just being defined in terms of semi-definite programmes.

Second, our framework provides a connection between the study of quantum channels and post-quantum steering, which is itself a novel observation. Starting from this connection, in section 3.3 we give new analytical examples of post-quantum steering constructed from non-localizable, yet causal, channels. In addition, section 3.2 shows that a consequence of post-quantum steering is the existence of non-localizable channels that cannot be used to violate a Bell inequality through any local operations whatsoever. We moreover give a characterization of non-signalling assemblages in terms of quantum states and unitary operations, which results in a diagramatic proof of the Gisin–Hughston–Jozsa–Wootters (GHJW) theorem in corollary 14. We show in section 4.4 that this proof of the GHJW theorem can be generalized to the study of non-classical teleportation,

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[5] Almost quantum correlations are defined as a particular relaxation of the set of quantum correlations in Bell scenarios. That is, the set of almost quantum correlations strictly contains those that are achievable by quantum mechanics. Almost quantum assemblages are defines as a particular relaxation of the set of quantum assemblages in steering scenarios. We revise the rigorous definition of these concepts in the next sections.

and we show in corollary 36 that post-quantum non-classical teleportation can only be witnessed if there are multiple black boxes in your network.

Finally, we are the first to highlight the possibilities of studying forms of post-quantum Buscemi non-locality and post-quantum non-classical teleportation. Our framework further outlines how to approach these through the study of quantum channels.

The paper is structured as follows. In section 1 we introduce a new family of quantum channels of utmost relevance in this work, while we review relevant known classes of channels in appendix A. In sections 2 and 3 we (i) discuss the interpretation of Bell and steering scenarios in terms of quantum channels, and (iii) present some results that follow when looking at these non-classical phenomena from the scope of quantum channels. The traditional scope to these phenomena is briefly reviewed in appendices B and C. Finally, in section 4 we discuss how our framework further includes the above mentioned Buscemi non-locality [23] and non-classical teleportation [24]. For clarity in the presentation, some of the proofs of results in the main body of the paper are presented in the appendix.

A quick note on notation. A Hilbert space will typically be denoted by $\mathcal{H}$, unless otherwise stated, and the set of positive operators acting on $\mathcal{H}$ with trace at most 1 will be denoted as $\mathcal{D}(\mathcal{H})$. Furthermore, for the more general set of linear operators acting on $\mathcal{H}$, we will use the notation $\mathcal{L}(\mathcal{H})$.

1. Quantum channels

In the study of non-locality in quantum physics and beyond, a common approach is to have the fundamental objects being a black box associated with some stochastic behaviour: for a given set of inputs for each party, an output is generated stochastically. A stochastic process should be suitably normalized, i.e. the sum over all outcomes for a given input is 1. The quantum analogue of such a process is a quantum channel. Recall that a channel Λ is a trace-preserving, completely-positive (CPTP) map. That is, given an input quantum state described by the density matrix ρ_i, a channel Λ acts on this system producing an output state with density matrix $\rho_o := \Lambda(\rho_i)$. The suitable normalization condition is then that the trace of ρ_o is 1 whenever $\mathrm{tr}\{\rho_i\} = 1$. A classical stochastic process can be encoded into a channel with respect to some orthonormal basis of the respective Hilbert space. To retrieve the probabilities in the stochastic process one only needs to prepare states in that basis as input, and then only measure in that basis.

Given these simple observations, one can readily relate quantum channels to the study of conditional probability distributions, and thus quantum non-locality. For example, we can ask which channels give rise to correlations that are compatible with a local hidden variable model, or otherwise. Such non-local properties of quantum channels have been observed and utilized in previous works [22, 26]. There, the relevant objects of study are semicausal and causal multipartite quantum channels, in particular the subset of localizable ones, which we formally review in appendix A. To sketch their definitions now, the causal channels are those where one party's output quantum state is the same for all input states for another party, and the localizable channels are those that are generated by local operations and shared entanglement between the parties. In this section we introduce a new class of channels, called the *almost localizable channels*, which will be pertinent when discussing non-locality and steering.

The general scenario we consider is that of multiple space-like separated parties such that they cannot use any particular physical system in their respective laboratories that could result in communication. In this way, the parties are subjected to the same conditions as in a Bell test. We can model the parties' global resources as a device with multiple input and output ports: an input and output port associated with each party. Therefore, each party can produce a local input quantum system, put it into their respective input port, and receive a quantum system from the output port. The global device can contain resources that are shared between distant parties, such as entanglement. For example, if we have two parties, and they each input a system into their respective devices, the output of both devices could be associated with an entangled quantum state. We will now make this picture more formal.

We have N parties labelled by an index $j \in \{1, 2, \ldots, N\}$, and each party has an input and output Hilbert space, $\mathcal{H}_{\mathrm{in}}^j$ and $\mathcal{H}_{\mathrm{out}}^j$, respectively, associated with the input and output ports of the parties' device[6]. The input quantum systems have states that are associated with the density matrix $\rho_{\mathrm{in}}^j \in \mathcal{D}(\mathcal{H}_{\mathrm{in}}^j)$. The N-partite device is then associated with a quantum channel $\Lambda_{1\ldots N} : \bigotimes_j^N \mathcal{L}(\mathcal{H}_{\mathrm{in}}^j) \to \bigotimes_j^N \mathcal{L}(\mathcal{H}_{\mathrm{out}}^j)$ taking the input state $\rho_{\mathrm{in}}^1 \otimes \rho_{\mathrm{in}}^2 \otimes \ldots \otimes \rho_{\mathrm{in}}^N$ to $\rho_{\mathrm{out}} := \Lambda_{1\ldots N}(\rho_{\mathrm{in}}^1 \otimes \rho_{\mathrm{in}}^2 \otimes \ldots \otimes \rho_{\mathrm{in}}^N)$. It will be convenient at times to take bipartitions $S_A, S_B \subseteq \{1, \ldots, N\}$ of the N parties, such that $S_A \cup S_B = \{1, \ldots, N\}$. With these bipartitions, we can then consider Hilbert spaces $\mathcal{H}_{\mathrm{in}}^{S_A}$, $\mathcal{H}_{\mathrm{in}}^{S_B}$ and $\mathcal{H}_{\mathrm{out}}^{S_A}$, $\mathcal{H}_{\mathrm{out}}^{S_B}$ associated with the input and output Hilbert spaces of S_A and S_B, respectively.

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[6] In this paper, all Hilbert spaces are assumed to be finite dimensional, unless otherwise specified.
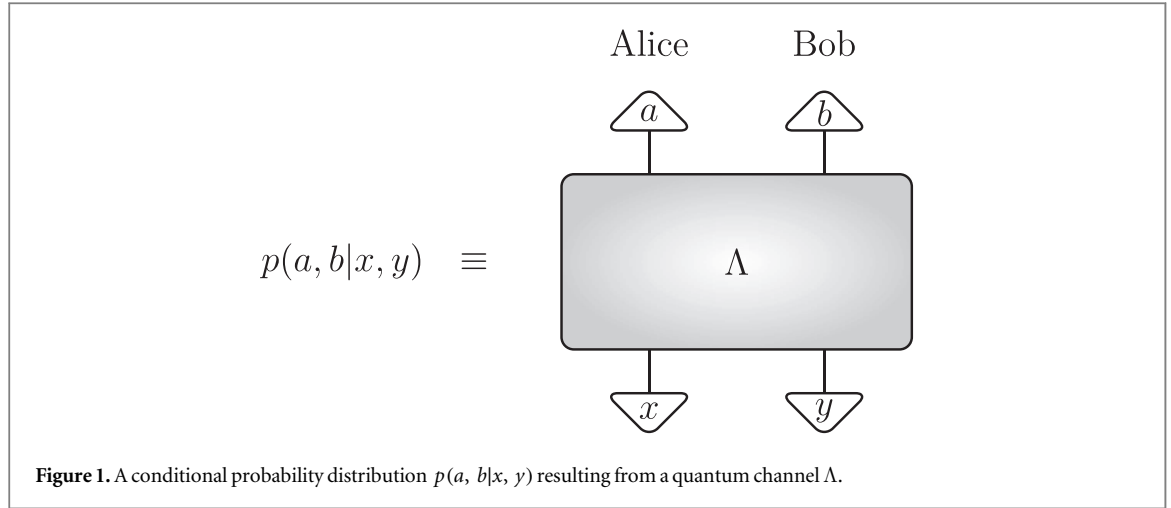
Figure 1. A conditional probability distribution $p(a, b|x, y)$ resulting from a quantum channel Λ.

Given this set-up, we can informally sketch the definition of semicausal and localizable channels in the bipartite case (i.e. $N = 2$). The formal definitions can be found in appendix A. A *semicausal* channel is one where the output state for one particular party is independent of the input state of the other party. In other words, the reduced quantum state for one of the parties is well-defined since it is independent of the other party's input. For example, if a channel Λ_{12} is semicausal from 1 to 2, denoted $1 \nrightarrow 2$, then the output state is $\rho_{\text{out}} = \Lambda_{12}(\rho_{\text{in}}^1 \otimes \rho_{\text{in}}^2)$ and if we trace out party 1, the output state of party 2 is $\rho_{\text{out}}^2 = \text{tr}_1(\Lambda_{12}(\rho_{\text{in}}^1 \otimes \rho_{\text{in}}^2)) = \Upsilon(\rho_{\text{in}}^2)$, for $\Upsilon : \mathcal{L}(\mathcal{H}_{\text{in}}^2) \rightarrow \mathcal{L}(\mathcal{H}_{\text{out}}^2)$ being a channel. A bipartite channel is *causal* if it is semicausal in both directions, i.e. $1 \nrightarrow 2$ and $2 \nrightarrow 1$. A bipartite channel is *localizable* if there exists a joint quantum state shared between the two parties such all the parties' maps are only from the jth party's input and their share of the entangled state to the jth party's output.

In this work, we will use diagrammatic representations of quantum channels where input and output systems to a channel are represented by wires, and the channels as boxes connecting inputs and outputs. One can see an example of such a diagram in figure 1, where Λ is the channel, and time (the flow from inputs to outputs) goes from bottom to top. Furthermore, later on, we will denote the preparation of states as triangles at the beginning of input wires, and measurements as triangles at the end of output wires.

Within this scenario we define a new class of channels called the *almost localizable channels* as follows:

Definition 1 (**Almost localizable channels**). A causal channel $\Lambda_{1\ldots N}$ is almost localizable if and only if there exists a global ancilla system E with Hilbert space $\mathcal{H}_E$, a local ancilla system E_k for each kth party, with input and output Hilbert spaces $\mathcal{H}_{\text{in}}^{E_k}$ and $\mathcal{H}_{\text{out}}^{E_k}$, respectively, and state $|\psi\rangle_E \in \mathcal{H}_E \otimes \mathcal{H}_{\text{in}}^{E_1} \otimes \mathcal{H}_{\text{in}}^{E_2} \otimes \ldots \otimes \mathcal{H}_{\text{in}}^{E_N}$ such that, for all states $\rho \in \mathcal{D}(\bigotimes_{j=1}^N \mathcal{H}_{\text{in}}^j)$,

$$\Lambda_{1\ldots N}[\rho] = \text{tr}_{EE_1\ldots E_N}\left\{ \prod_{j=1}^N U_{jE} \left(\rho \otimes |\psi\rangle \langle\psi|_E \right) \prod_{k=0}^{N-1} U_{N-kE}^\dagger \right\},$$

where $U_{kE} : \mathcal{H}_{\text{in}}^k \otimes \mathcal{H}_E \otimes \mathcal{H}_{\text{in}}^{E_k} \rightarrow \mathcal{H}_{\text{out}}^k \otimes \mathcal{H}_E \otimes \mathcal{H}_{\text{out}}^{E_k}$ is a unitary operator for all k, such that, for any permutation π on the set $\{1, 2, \ldots, N\}$,

$$\prod_{j=1}^N U_{jE} \left(\rho \otimes |\psi\rangle \langle\psi|_E \right) \prod_{k=0}^{N-1} U_{(N-k)E}^\dagger = \prod_{j=1}^N U_{\pi(j)E} \left(\rho \otimes |\psi\rangle \langle\psi|_E \right) \prod_{k=0}^{N-1} U_{\pi(N-k)E}^\dagger.$$

Notice that in this definition for localizable channels, the ancilla σ_R is the same for all inputs to the channel $\Lambda_{1\ldots N}$. If we compare this definition with that of localizable channels as given by definition 42, we see that almost localizable channels are a natural generalization of the localizable ones. Indeed, in definition 42, the condition of the representation that for all permutations π, $\prod_{k=1}^N U_{kE} = \prod_{k=1}^N U_{\pi(k)E}$ is equivalent to the constraint that $\prod_{k=1}^N U_{kE}|\psi\rangle = \prod_{k=1}^N U_{\pi(k)E}|\psi\rangle$ for all possible states $|\psi\rangle \in \mathcal{H}_{\text{in}}^E \bigotimes_{j=1}^N \mathcal{H}_{\text{in}}^j$. This last universal quantifier over all ancilla states can be relaxed further to an existential quantifier, i.e. that there exists a state $|\psi\rangle$ such that the unitary operators' ordering is invariant under permutations of the parties. This relaxation precisely gives the set of almost localizable channels.

Note that localizable channels are by definition almost localizable, as well as causal. However, as we will show in section 2.2, there exist almost localizable channels that are not localizable. In showing this, we use the close connection between the so-called *almost quantum correlations* defined in [25] (see appendix B), and the almost

localizable channels. Indeed, the motivation for the name almost localizable comes from this connection. In this direction we also generalize this connection to the study of steering in section 3.1.

2. Non-locality from the scope of quantum channels

In this section, we reinterpret the traditional Bell scenario [1] in terms of quantum channels. In particular, we connect every quantum channel to a family of correlations in a Bell test. We emphasize that non-locality can, in a sense, be studied independently of the quantum formalism, so considering all processes as fundamentally quantum may seem *excessive*. Instead, one can see our review of non-locality from the point-of-view of quantum channels as just the beginning of a bigger story, as will hopefully become clear. We review the traditional notion of a Bell scenario and its relevant sets of correlations in appendix B.

2.1. Non-locality via quantum channels

A Bell scenario is characterized by the parameters (N, m, d), where N denotes the number of parties, m the number of measurements each party can choose from, each with d possible outcomes. Consider now the parties to have input and output Hilbert spaces given by $\mathcal{H}_{\mathrm{in}}^{j} = \mathcal{H}_{\mathrm{in}}^{j'} = \mathcal{H}_{m}$ and $\mathcal{H}_{\mathrm{out}}^{j} = \mathcal{H}_{\mathrm{out}}^{j'} = \mathcal{H}_{d}$ for all $j \neq j'$, where $\mathcal{H}_{m}$ has dimension m and $\mathcal{H}_{d}$ dimension d. Denote by $\{|x\rangle\}_{x=1:m}$ an orthonormal basis of $\mathcal{H}_{m}$, and by $\{|a\rangle\}_{a=1:d}$ an orthonormal basis of $\mathcal{H}_{d}$. In what follows, we relate channels of the form $\Lambda_{1\ldots N} : \mathcal{L}(\mathcal{H}_{m}^{\otimes N}) \to \mathcal{L}(\mathcal{H}_{d}^{\otimes N})$ to correlations in a Bell scenario.

Definition 2. A conditional probability distribution $p(a_1 \ldots a_N | x_1 \ldots x_N)$ in a Bell scenario is channel-defined if there exists a channel $\Lambda_{1\ldots N} : \mathcal{L}(\mathcal{H}_{m}^{\otimes N}) \to \mathcal{L}(\mathcal{H}_{d}^{\otimes N})$, and some choice of orthonormal bases $\{|x_j\rangle \in \mathcal{H}_{m}\}_{x=1:m}$ and $\{|a_j\rangle \in \mathcal{H}_{d}\}_{a=1:d}$ for each jth party, such that

$$p(a_1 \ldots a_N | x_1 \ldots x_N) = \mathrm{tr}\{\otimes_{k=1}^{N} |a_k\rangle \langle a_k| \, \Lambda_{1\ldots N}(\otimes_{k=1}^{N} |x_k\rangle \langle x_k|)\}. \tag{1}$$

Given a channel $\Lambda_{1\ldots N}$, it is always possible to define correlations resulting from it for a given choice of input and output orthonormal bases. Figure 1 sketches (in the bipartite case) this construction of correlations schematically. Given this connection, we can now directly relate the families of correlations presented earlier to families of channels presented in section 1 and appendix A. Although the results pertinent to non-signalling, quantum and classical correlations were noticed in previous works [22, 26, 27], we present all proofs in appendix E.1.
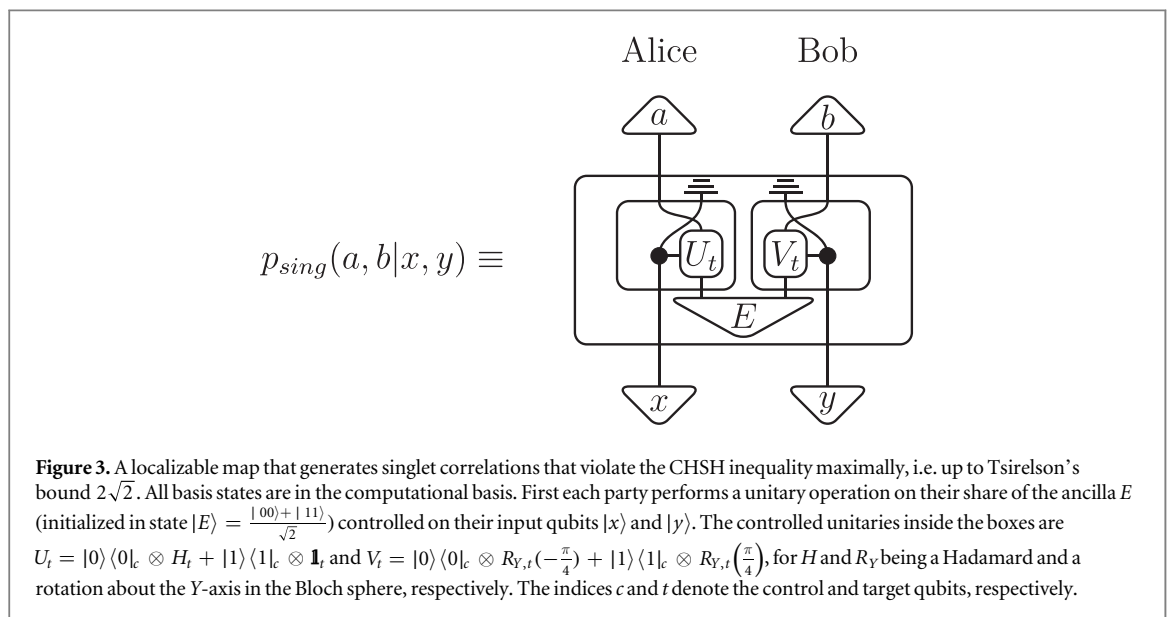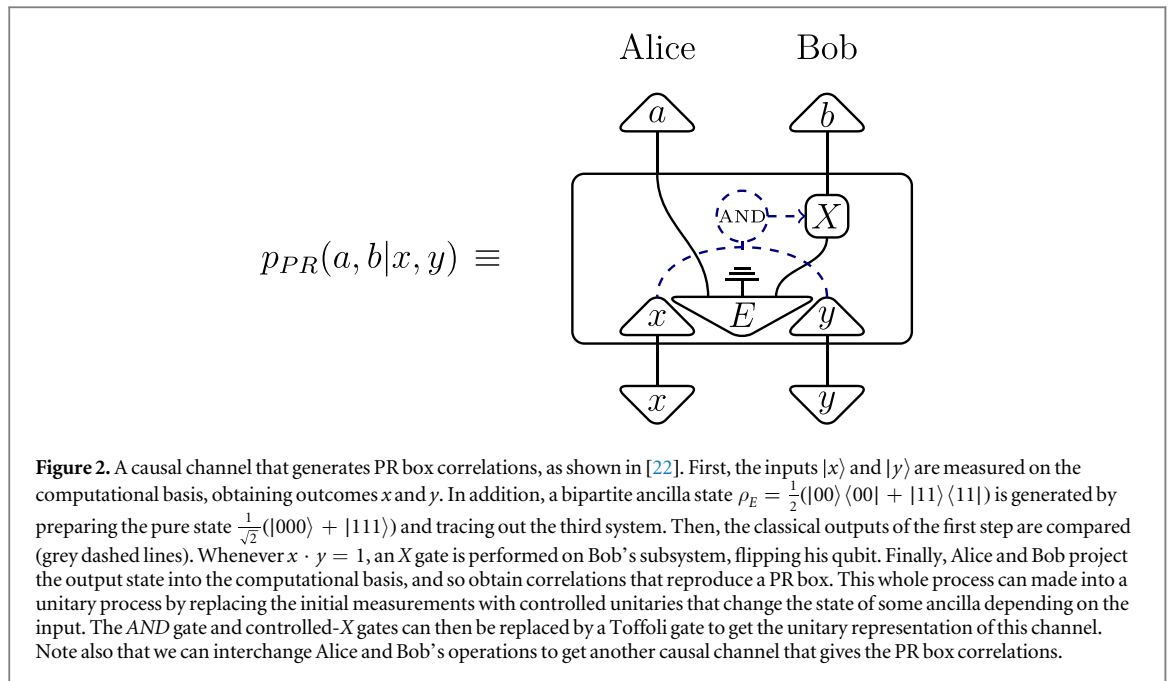
Proposition 3. *A conditional probability distribution $p(a_1 \ldots a_N | x_1 \ldots x_N)$ is non-signalling if and only if there exists a causal channel $\Lambda_{1\ldots N}^{\mathsf{C}} : \mathcal{L}(\mathcal{H}_{m}^{\otimes N}) \to \mathcal{L}(\mathcal{H}_{d}^{\otimes N})$ such that the distribution is channel-defined by $\Lambda_{1\ldots N}^{\mathsf{C}}$.*

In figure 2 we show how the example of a Popescu–Rohrlich (PR) non-local box can be realized by a causal channel. The PR non-local box is a device that can violate the Clauser–Horne–Shimony–Holt (CHSH) inequality beyond Tsirelson's bound, and thus cannot be realized by local measurements on an entangled state [17]. The statistics produced by a PR box, for binary inputs and outputs, are $p(a, b|x, y) = \frac{1}{2}\delta_{a\oplus b}^{xy}$, where $\oplus$ is addition modulo 2. The channel in this figure is an entanglement-breaking channel [27], and thus its Choi state Ω is separable across the partition of Alice's and Bob's input and output Hilbert spaces. However, non-localizable causal channels that are not entanglement-breaking have been constructed in the literature [26], and we will refer to one such channel later. How can one detect non-localizability in a particular channel? One possible approach is through the correlations that are channel-defined by that channel, as described in the following result.

Proposition 4. *A conditional probability distribution $p(a_1 \ldots a_N | x_1 \ldots x_N)$ is quantum if and only if there exists a localizable channel $\Lambda_{1\ldots N}^{\mathsf{Q}} : \mathcal{L}(\mathcal{H}_{m}^{\otimes N}) \to \mathcal{L}(\mathcal{H}_{d}^{\otimes N})$ such that the distribution is channel-defined by $\Lambda_{1\ldots N}^{\mathsf{Q}}$.*

In figure 3 we present the example of a localizable channel that channel-defines the correlations $p_{\mathrm{sing}}(a, b|x, y)$ which give Tsirelson's bound for the CHSH inequality [28], i.e. the maximal violation for local measurements on an entangled state. We present the channel in terms of its unitary representation.

Given a channel Λ, if the correlations that are channel-defined by it are not compatible with quantum correlations, i.e. they are post-quantum correlations, then the channel was not localizable. For example, if one obtains correlations that are channel-defined by a channel Λ, and then computes their CHSH value, if this exceeds Tsirelson's bound, the channel Λ is non-localizable. Indeed, this is how it is shown that the channel in figure 2 is non-localizable, as well as the channel given in [26].

Figure 2. A causal channel that generates PR box correlations, as shown in [22]. First, the inputs $|x\rangle$ and $|y\rangle$ are measured on the computational basis, obtaining outcomes x and y. In addition, a bipartite ancilla state $\rho_E = \frac{1}{2}(|00\rangle\langle00| + |11\rangle\langle11|)$ is generated by preparing the pure state $\frac{1}{\sqrt{2}}(|000\rangle + |111\rangle)$ and tracing out the third system. Then, the classical outputs of the first step are compared (grey dashed lines). Whenever $x \cdot y = 1$, an X gate is performed on Bob's subsystem, flipping his qubit. Finally, Alice and Bob project the output state into the computational basis, and so obtain correlations that reproduce a PR box. This whole process can made into a unitary process by replacing the initial measurements with controlled unitaries that change the state of some ancilla depending on the input. The *AND* gate and controlled-*X* gates can then be replaced by a Toffoli gate to get the unitary representation of this channel. Note also that we can interchange Alice and Bob's operations to get another causal channel that gives the PR box correlations.

Figure 3. A localizable map that generates singlet correlations that violate the CHSH inequality maximally, i.e. up to Tsirelson's bound $2\sqrt{2}$. All basis states are in the computational basis. First each party performs a unitary operation on their share of the ancilla E (initialized in state $|E\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$) controlled on their input qubits $|x\rangle$ and $|y\rangle$. The controlled unitaries inside the boxes are $U_t = |0\rangle\langle0|_c \otimes H_t + |1\rangle\langle1|_c \otimes \mathbb{1}_t$ and $V_t = |0\rangle\langle0|_c \otimes R_{Y,t}(-\frac{\pi}{4}) + |1\rangle\langle1|_c \otimes R_{Y,t}(\frac{\pi}{4})$, for H and R_Y being a Hadamard and a rotation about the Y-axis in the Bloch sphere, respectively. The indices c and t denote the control and target qubits, respectively.

Proposition 5. *A conditional probability distribution $p(a_1 \ldots a_N | x_1 \ldots x_N)$ is classical if and only if there exists a local channel $\Lambda^{\mathsf{L}}_{1 \ldots N} \colon \mathcal{L}(\mathcal{H}_m^{\otimes N}) \to \mathcal{L}(\mathcal{H}_d^{\otimes N})$ such that the distribution is channel-defined by $\Lambda^{\mathsf{L}}_{1 \ldots N}$.*

It should be noted that there can exist non-local but localizable channels that will only channel-define classical correlations. A simple bipartite example of such a channel is one where the maximally entangled two-qubit state is prepared in the ancilla register, the input systems are discarded (or traced out), and each party's output is one half of the two-qubit register. For this channel, correlations are generated by each party is measuring one half of a maximally entangled state in a fixed basis, which can be reproduced by classical correlations.

Finally, we now address the set of almost quantum correlations. We have included the proof of the following result, since it will be useful for our subsequent discussion.

Proposition 6. *A conditional probability distribution $p(a_1 \ldots a_N | x_1 \ldots x_N)$ is almost quantum if and only if there exists an almost localizable channel $\Lambda^{\bar{\mathsf{Q}}}_{1 \ldots N} \colon \mathcal{L}(\mathcal{H}_m^{\otimes N}) \to \mathcal{L}(\mathcal{H}_d^{\otimes N})$ such that the distribution is channel-defined by $\Lambda^{\bar{\mathsf{Q}}}_{1 \ldots N}$.*
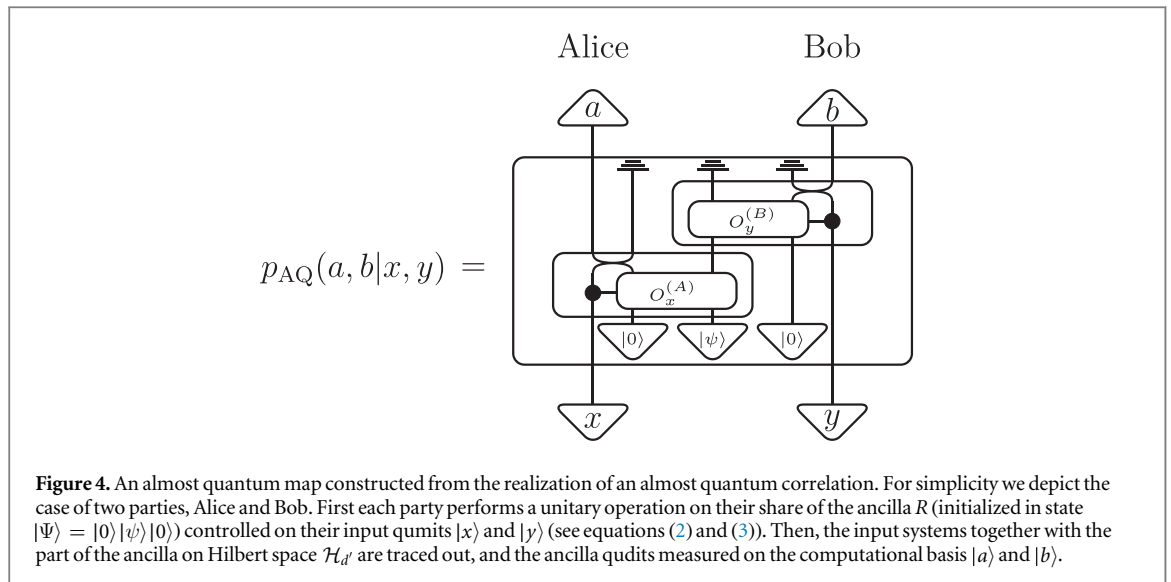
Figure 4. An almost quantum map constructed from the realization of an almost quantum correlation. For simplicity we depict the case of two parties, Alice and Bob. First each party performs a unitary operation on their share of the ancilla R (initialized in state $|\Psi\rangle = |0\rangle |\psi\rangle |0\rangle$) controlled on their input qumits $|x\rangle$ and $|y\rangle$ (see equations (2) and (3)). Then, the input systems together with the part of the ancilla on Hilbert space $\mathcal{H}_{d'}$ are traced out, and the ancilla qudits measured on the computational basis $|a\rangle$ and $|b\rangle$.

Proof. First we take a probability distribution $p(a_1 \ldots a_N | x_1 \ldots x_N)$ that is in the set of almost quantum correlations. Let $\Pi^{(i)}_{a_i|x_i}$ and $|\psi\rangle$ be the projectors and state that realize[7] this distribution, which have an associated Hilbert space $\mathcal{H}_{d'}$ of dimension d'. From these we will define an almost localizable channel $\Lambda^{\tilde{Q}}_{1 \ldots N} \colon \mathcal{L}(\mathcal{H}_m^{\otimes N}) \to \mathcal{L}(\mathcal{H}_d^{\otimes N})$ such that the correlations are channel-defined for this channel. Let the ancilla be a quantum system initialized on the state $|\Psi\rangle = |\psi\rangle \otimes |0\rangle^{\otimes N}$, where these extra N systems are qudits in the Hilbert space $\mathcal{H}_d$, one for each party, initialized in $|0\rangle$. Now define the following operators for each party:

$$O_x^{(i)} = \sum_{j=1}^{d} \Pi^{(i)}_{j|x} \otimes A_j, \tag{2}$$

where $A_1 = \mathbf{1}_d$ and $A_j = |j\rangle \langle 0| + |0\rangle \langle j| + \sum_{\alpha=1, \alpha \neq j}^{d-1} |\alpha\rangle \langle \alpha|$ for $j = 1 : d - 1$. The operator $\Pi^{(i)}_{j|x}$ acts on $\mathcal{H}_{d'}$, A_j on the ancillary qudit that corresponds to party i. Then define, for each party, controlled unitary operations on their input qumit in Hilbert space $\mathcal{H}_m$ together with their ancillary system as follows:

$$U^{(i)} = \sum_{x=1}^{m} |x\rangle \langle x| \otimes O_x^{(i)}. \tag{3}$$

Now the almost localizable channel can be defined, as in figure 4. Each party has as input system a qumit, the unitary representation of the channel is given by the $U^{(i)}$ followed by a swap on the ancilla qudit and the input qumit for each party. Finally, the output system for each party is their corresponding qudit, and the input qumits and the ancillary subsystem on $\mathcal{H}_{d'}$ are traced out. The commutation relations of the $\Pi^{(i)}_{j|x}$ on the state $|\psi\rangle$ imply that the unitaries U^i associated with the different parties commute on the ancilla, thus implying the channel is almost localizable. It is straightforward to check that the correlation $p(a_1 \ldots a_N | x_1 \ldots x_N)$ is recovered by the parties inputting their measurement settings $|x_i\rangle$, and measuring their output systems in the basis $\{|a_i\rangle\}$.

So far we have seen that an almost localizable channel can be constructed from almost quantum correlations. Given an almost localizable channel, it is relatively straightforward to see that the channel-defined correlations will be almost quantum. Note that the action at each jth party of preparing an input state, followed by a unitary and then a projective measurement can be simulated by the projective measurements $\Pi_{a_j|x_j}$ on the ancilla state $|E\rangle$ as per the definition of almost localizable channels. These projectors will then satisfy the properties required to produce almost quantum correlations given the definition of an almost localizable channel. □

2.2. Connections between channels and correlations
In this section, we first comment on how, given some correlations in a Bell scenario, one can find a canonical channel that channel-defines them. Then, we elaborate on further ways one may use a channel to generate correlations.

We have previously considered how correlations result from channels. One can then readily ask how channels can be constructed once we are given a set of correlations. Given correlations $p(a_1 \ldots a_N | x_1 \ldots x_N)$, there is a *canonical channel* that channel-defines them, which amounts to a controlled preparation of a quantum system. In particular, for a given choice of input and output orthonormal bases $\{|a_j\rangle\}$, $\{|x_j\rangle\}$ for all N parties,

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[7] The state ρ that realizes an almost quantum correlation can be taken to be a pure state, without loss of generality when there are no restrictions on the dimension of the Hilbert space.

such a channel is defined as

$$\Lambda_{1\,\ldots\,N}^c(\cdot) = \sum_{x_1,\,\ldots,\,x_N} \sum_{a_1,\,\ldots,\,a_N} p(a_1\ldots a_N|x_1\ldots x_N)|a_1\ldots a_N\rangle\langle x_1\ldots x_N|(\cdot)|x_1\ldots x_N\rangle\langle a_1\ldots a_N|. \tag{4}$$

It can be readily seen that $\Lambda_{1\,\ldots\,N}^c$ channel-defines the correlations $p(a_1\ldots a_N|x_1\ldots x_N)$ for the choice of input and output orthonormal bases $\{|a_j\rangle\}$, $\{|x_j\rangle\}$ for all N parties. We also remark that one can take any channel Λ that channel-defines some correlations with given preparations and measurements, and then construct the canonical channel Λ^c from Λ with those preparations and measurements. This intuitively amounts of the taking the original channel and applying fully decoherent channels to the inputs and outputs.

Now, let us elaborate further on how correlations may arise from the use of quantum channels. If we are given a particular channel $\Lambda_{1\ldots N}$, indeed, choosing a set of orthonormal bases such that correlations are channel-defined by $\Lambda_{1\ldots N}$ may not be *optimal* for witnessing non-locality. That is, given access to a channel, correlations can be generated through more elaborate means than just preparing a state from an orthonormal basis, plugging it into a local port of the channel, and then measuring in another basis. For example, one party could prepare a bipartite system and send one half of it into the channel, then after the system has emerged from the channel, one can jointly measure this output and the remaining half of the bipartite system[8]. More formally, with each party j, in addition to the Hilbert spaces associated with the jth input and output ports of the channel, we associate an auxiliary Hilbert space $\mathcal{H}_{aux}^j$. Then for a given input x_j for the jth party, without loss of generality, this party can prepare a state $\rho_{x_j} \in \mathcal{D}(\mathcal{H}_{in}^j \otimes \mathcal{H}_{aux}^j)$, and then the output a_j is associated with some POVM element $M_{a_j} \in \mathcal{L}(\mathcal{H}_{out}^j \otimes \mathcal{H}_{aux}^j)$, such that $\sum_{a_j} M_{a_j} = \mathbf{1}$ and every element M_{a_j} is a positive operator[9]. Putting this together, given a channel $\Lambda_{1\ldots N}$, correlations are generated as:

$$p(a_1, a_2, \ldots, a_N|x_1, x_2, \ldots, x_N) = \mathrm{tr}\left\{ \bigotimes_{j=1}^N M_{a_j}\left(\Lambda_{1\ldots N} \otimes \mathbf{1}_{aux}(\bigotimes_{j=1}^N \rho_{x_j})\right)\right\}, \tag{5}$$

where $\mathbf{1}_{aux}$ is the identity operator acting on all Hilbert spaces $\mathcal{H}_{aux}^j$. This allows us to explore whether a particular channel may result in non-local correlations, as we now formalize in the following definition.

Definition 7 (Local-limited channels). If for all states ρ_{x_j} and measurements $\{M_{a_j}\}$, a channel $\Lambda_{1\ldots N}$ never produces non-local correlations via equation (5), then the channel is local-limited.

There are several channels which are local-limited. As an example, consider the entangled quantum states that can never produce non-local correlations for all general measurements [29, 30]. These quantum states can give rise to localizable channels that are not local yet are local-limited. The construction goes as follows. Take a localizable channel where the ancillary system is initiated in such an entangled quantum state. In addition, the 'unitary operations' between the input and ancillary ports of each party simply trace out the input states. For all practical purposes then, this channel only prepares a fixed quantum state among the parties, which then goes to the output ports. It follows that event though such channel is not local, it is however local-limited.

In general, if $\Lambda_{1\ldots N}^c$ is a *canonical* channel for local correlations $p(a_1\ldots a_N|x_1\ldots x_N)$, we have the following result, which is proven in section E.3 of the appendix.

Proposition 8. *Given $\Lambda_{1\ldots N}^c(\cdot)$ from $p(a_1\ldots a_N|x_1\ldots x_N)$, for all measurements $M_{a_j'}$, and all states $\rho_{x_j'}$, the correlations*

$$p(a_1', a_2', \ldots, a_N'|x_1', x_2', \ldots, x_N') = \mathrm{tr}\left\{ \bigotimes_{j=1}^N M_{a_j'}\left(\Lambda_{1\ldots N}^c \otimes \mathbf{1}_{aux}(\bigotimes_{j=1}^N \rho_{x_j'})\right)\right\}$$

are local if the correlations $p(a_1\ldots a_N|x_1\ldots x_N)$ are local.

Another interesting question is that of constructing almost localizable channels that are non-localizable. The following method works for any general Bell scenario as a starting point, depending on which type of channel one wishes to construct, and goes beyond the canonical form previously discussed. For the sake of simplicity, however, we focus on a bipartite Bell scenario with two dichotomic measurements per party.

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[8] A more general strategy would be to apply an instrument with a quantum memory to the channel. That is, preparing a bipartite state, and then sequentially using the channel, in between each use a party applies an operation to the output of the channel and the other half of the bipartite system (stored in a memory). This would be in analogy to performing a Bell test through collective measurements on a number of quantum states.
[9] We do not need to explicitly consider choices of different measurements for $\{M_a\}$, since the state ρ_x carries the information about the input.

First, take an almost quantum correlation $p(a, b|x, y)$ with no quantum realization. Such correlations can be found by taking those that violate Bell inequalities beyond a Tsirelson-like bound, as presented in [25]. Then, obtain a state and measurements that reproduce the correlations as outlined in [31]. Using the protocol described in the proof of proposition 6, also depicted in figure 4, an almost localizable channel that channel-defines these correlations can be hence constructed from these 'state and measurements'. This almost localizable channel is hence provably non-localizable, since it channel-defines Bell correlations beyond what quantum theory allows, and completes the picture of the hierarchy of channels in theorem 44.

Finally, while proposition 6 tells us that almost localizable channels channel-define the almost quantum correlations, does this mean the correlations in equation (5) that are generated by an almost localizable channel $\Lambda_{1...N}$ will necessarily be almost quantum correlations? The answer does not follow immediately from the statement of proposition 6, but the proof of this theorem can be slightly extended to give an answer in the affirmative. To sketch this extension, first note that all states ρ_{x_j} and measurements M_{a_j} can be made pure and projective, respectively, by introducing a large enough auxiliary system for each party. That is, ρ_{x_j} can be replaced by a pure state $|\psi_{x_j}\rangle$ in a larger space, and then we can rewrite these states to be $|\psi_{x_j}\rangle = V_{x_j}|\psi_0^j\rangle$ for some fixed state $|\psi_0\rangle$. Now if we apply an almost localizable channel to (part of) these input states, the whole process can be modelled as preparing the state $|\psi\rangle_E \bigotimes |\psi_0^j\rangle$, then applying V_{x_j} to the input states, followed by the unitaries in the almost localizable channel. Finally, a projective measurement is made on the output qubits. This whole process is equivalent to applying the inverse of the unitaries to these projective measurements to form new projective measurements which act on the state $|\psi\rangle_E \bigotimes |\psi_0^j\rangle$. These new projective measurements, due to the definition of the almost localizable channel, will 'commute' for the particular state $|\psi\rangle_E \bigotimes |\psi_0^j\rangle$, and thus will generate almost quantum correlations by definition. Note that due to proposition 6, given almost quantum correlations, we can always find states and measurements and an almost localizable channel that reproduce these correlations.

3. Steering from the scope of quantum channels

Steering refers to the phenomenon where one party, Alice, by performing measurements on one half of a shared state, seemingly remotely 'steers' the states held by a distant party, Bob, in a way which has no classical explanation [8]. This resembles the phenomenon of non-locality presented in last section, but with a slight change: now one party describes its system as a quantum system. In this section, we discuss an approach to studying steering via quantum channels. Here, we review the traditional notion of a steering scenario, while its relevant sets of assemblages are presented in appendix C.

In a bipartite steering scenario, the actions of one party (here Alice, also referred to as 'untrusted' or 'uncharacterized'[10]) are described solely by m possible classical inputs to her system, labelled by $x \in \{1 \ldots m\}$, each of which results in one of d possible classical outputs, labelled $a \in \in \{1 \ldots d\}$. The second party (Bob, also referred to as 'trusted' or 'characterized') fully describes the state of his share of the system by a subnormalised quantum state $\sigma_{a|x} \in \mathcal{L}(\mathcal{H}_B)$, where $\mathcal{H}_B$ is the Hilbert space associated with Bob's quantum system with dimension d_B. The set of subnormalised conditional states Alice prepares on Bob's side $\{\sigma_{a|x}\}_{a,x}$ is usually called *assemblage*, and $p(a|x) = \mathrm{tr}\{\sigma_{a|x}\}$ denotes the probability that such a subnormalised state is prepared, i.e. the probability that Alice obtains a when measuring x.

In this work we go beyond the bipartite definition of steering, and consider a setting with N untrusted parties and a single trusted party, still called Bob, who has some associated Hilbert space $\mathcal{H}_B$. Now, we have N Alices, where for the jth Alice, her input is $x_j \in \{1 \ldots m\}$ and output is $a_j \in \{1 \ldots d\}$. As a result Bob obtains an assemblage $\{\sigma_{a_1 \ldots a_N|x_1 \ldots x_N}\}_{a_1 \ldots a_N, x_1, \ldots x_N}$ with elements $\sigma_{a_1 \ldots a_N|x_1 \ldots x_N} \in \mathcal{L}(\mathcal{H}_B)$ such that $p(a_1 \ldots a_N|x_1 \ldots x_N) = \mathrm{tr}\{\sigma_{a_1 \ldots a_N|x_1 \ldots x_N}\}$. As with Bell scenarios, for the case of $N \leqslant 2$ we will use the same notation of inputs being x and y and outputs being a and b.

3.1. Steering via quantum channels

Here we extend the ideas of section 2.1 to steering scenarios, which provides a novel way to understanding the phenomenon. First we introduce the formalism and then characterize the channels that give rise to each set of assemblages.

A steering scenario is characterized by N untrusted parties, each of which generate one of m possible inputs, of d possible outcomes each, and a trusted party Bob with Hilbert space $\mathcal{H}_B$ with dimension d_B. Consider now all $(N + 1)$ parties (including Bob) to have input and output Hilbert spaces. For the N untrusted parties, these Hilbert spaces are $\mathcal{H}_{\mathrm{in}}^j = \mathcal{H}_{\mathrm{in}}^{j'} = \mathcal{H}_m$ and $\mathcal{H}_{\mathrm{out}}^j = \mathcal{H}_{\mathrm{out}}^{j'} = \mathcal{H}_d$ for all $j \neq j'$, where $\mathcal{H}_m$ has dimension m and $\mathcal{H}_d$

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[10] The sense in which the parties are untrusted is that whatever is used to produce classical data is some black box.
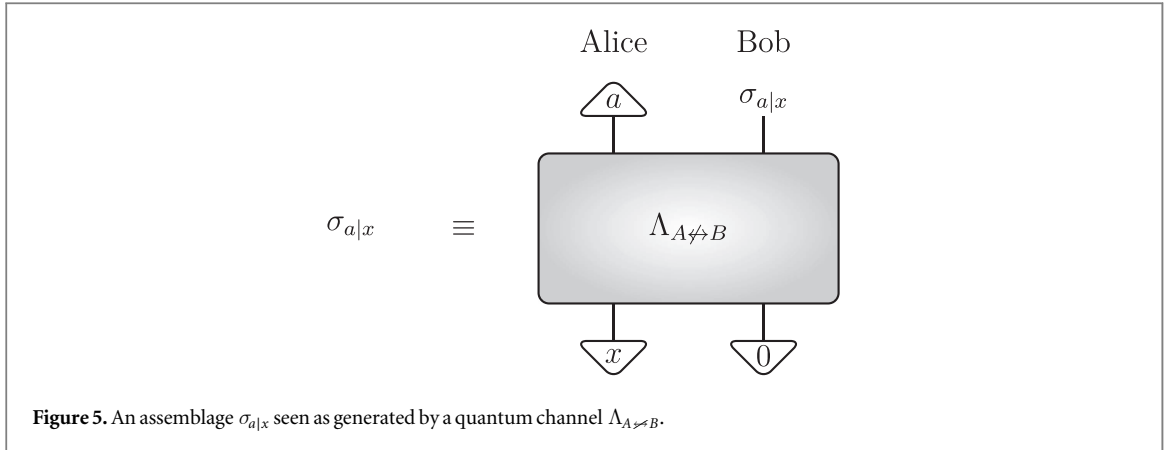
Figure 5. An assemblage $\sigma_{a|x}$ seen as generated by a quantum channel $\Lambda_{A \not\rightarrow B}$.

dimension d. Denote by $\{|x\rangle\}_{x=1:m}$ an orthonormal basis of $\mathcal{H}_m$, and by $\{|a\rangle\}_{a=1:d}$ an orthonormal basis of $\mathcal{H}_d$. For Bob, he has Hilbert spaces $\mathcal{H}_{B_{\text{in}}}$ and $\mathcal{H}_{B_{\text{out}}}$, which are taken to be equal[11]. In what follows, we relate channels of the form $\Lambda_{1\ldots N,B}: \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{\text{in}}}) \rightarrow \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{\text{out}}})$ to assemblages in a steering scenario as in figure 5.

Definition 9. An assemblage $\{\sigma_{a_1 \ldots a_N|x_1 \ldots x_N} \in \mathcal{D}(\mathcal{H}_B)\}$ in a steering scenario is channel-defined if there exists a channel $\Lambda_{1 \ldots N,B}: \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{\text{in}}}) \rightarrow \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{\text{out}}})$, some choice of orthonormal bases $\{|x_j\rangle \in \mathcal{H}_m\}_{x=1:m}$ and $\{|a_j\rangle \in \mathcal{H}_d\}_{a=1:d}$ for each party j, and a state $|0\rangle \in \mathcal{H}_{B_{\text{in}}}$, such that

$$\sigma_{a_1 \ldots a_N|x_1 \ldots x_N} = \text{tr}_{1 \ldots N}\{\otimes_{k=1}^N |a_k\rangle \langle a_k| \otimes \mathbb{1}_B \, \Lambda_{1 \ldots N,B}[\otimes_{k=1}^n |x_k\rangle \langle x_k| \otimes |0\rangle \langle 0|]\}, \tag{6}$$

where the partial trace is taken over all N untrusted systems.

Note that the main difference between the correlations and assemblages from the point-of-view of channels is that one of the outputs of the channels is left unmeasured, and Bob has a fixed input state $|0\rangle$. We now relate channels to the families of assemblages presented in appendix C, starting with the LHS assemblages.

Proposition 10. *An assemblage $\{\sigma_{a_1 \ldots a_N|x_1 \ldots x_N}\}$ is a LHS assemblage if and only if there exists a local channel $\Lambda_{1 \ldots N,B}^{\mathsf{L}}: \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{\text{in}}}) \rightarrow \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{\text{out}}})$ such that the assemblage is channel-defined by $\Lambda_{1 \ldots N,B}^{\mathsf{L}}$.*

In the literature, most of the focus has been on detecting whether an assemblage has a LHS model, thus revealing entanglement in a shared resource. It is therefore reassuring that the LHS assemblages do not involve entanglement when viewed through the channels that define them. Note that it is possible to have an assemblage which does not have a LHS model, yet the correlations resulting any measurement Bob makes on the assemblage can be local. In other words, steering is a distinct phenomenon from non-locality.

Proposition 11. *An assemblage $\{\sigma_{a_1 \ldots a_N|x_1 \ldots x_N}\}$ is non-signalling if and only if there exists a causal channel $\Lambda_{1 \ldots N,B}^{\mathsf{C}}: \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{\text{in}}}) \rightarrow \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{\text{out}}})$ such that the assemblage is channel-defined by $\Lambda_{1 \ldots N,B}^{\mathsf{C}}$.*

Given this definition, if one is given a non-signalling assemblage then it is straightforward to find a causal channel that channel-defines the assemblage if the input states and output measurements are fixed. In fact it is an SDP that is a slight modification of the SDP that decides whether a channel is causal as outlined in appendix A. Given elements of the assemblage, since they are channel-defined, this just results in linear constraints made on the channel.

A consequence of the above proposition and the unitary representation of causal channels is the following theorem.

Theorem 12 (Unitary representation of non-signalling assemblages). *Let $\{\sigma_{a_1 \ldots a_N|x_1 \ldots x_N}\}$ be a non-signalling assemblage. Then, the assemblage is channel-defined by a channel $\Lambda_{1 \ldots N,B}^{\mathsf{C}}: \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{\text{in}}}) \rightarrow \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{\text{out}}})$ if and only if there exist*

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[11] For the specific purposes of studying steering, we could equally take $\mathcal{H}_{B_{\text{in}}}$ to be $\mathbb{C}$, i.e. the scalars, but for the sake of simplicity in our presentation, we have this more symmetric set-up.
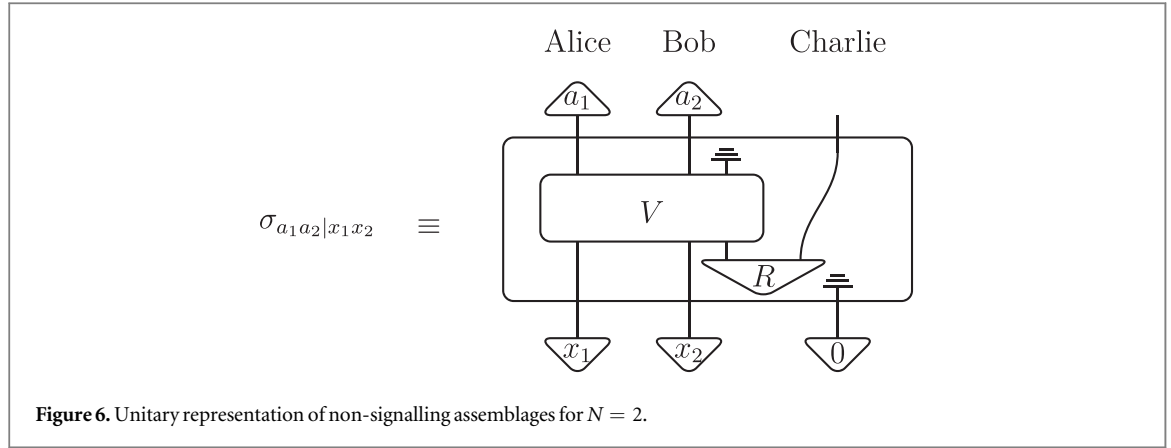
Figure 6. Unitary representation of non-signalling assemblages for $N = 2$.
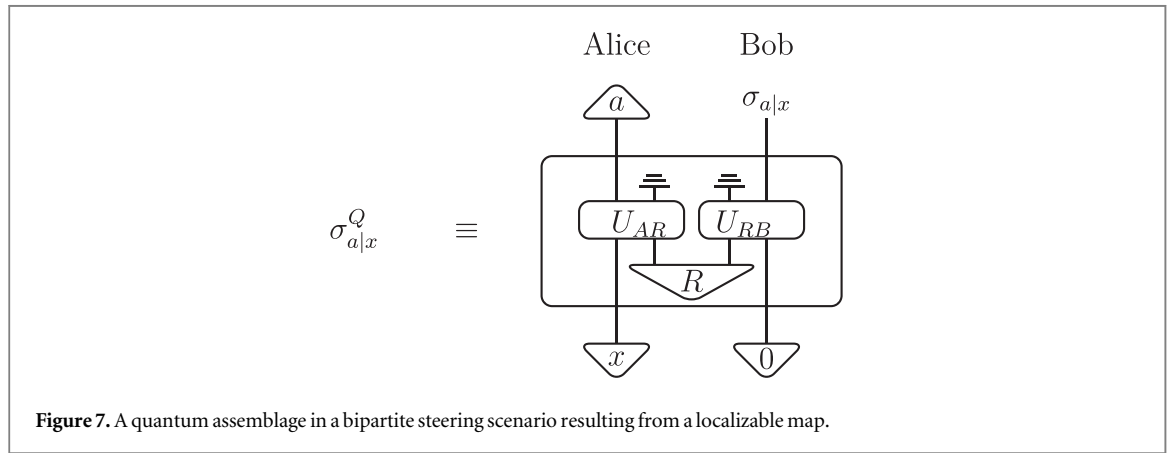

Figure 7. A quantum assemblage in a bipartite steering scenario resulting from a localizable map.

- *auxiliary systems E and E′ with input and output Hilbert spaces, $\mathcal{H}_{\mathrm{in}}^{E}$ and $\mathcal{H}_{\mathrm{out}}^{E}$ for E, with $\mathcal{H}_{\mathrm{in}}^{E'} = \mathcal{H}_{B}$ and $\mathcal{H}_{\mathrm{out}}^{E'} = \mathcal{H}_{B_{\mathrm{out}}}$ for E′, that is the output Hilbert space of E′ and B coincide;*

- *quantum state $|R\rangle \in \mathcal{H}_{\mathrm{in}}^{E} \otimes \mathcal{H}_{\mathrm{in}}^{E'}$;*

- *unitary operator $V : \mathcal{H}_{d}^{\otimes N} \otimes \mathcal{H}_{\mathrm{in}}^{E} \to \mathcal{H}_{d}^{\otimes N} \otimes \mathcal{H}_{\mathrm{out}}^{E}$,*

which produce a unitary representation of the channel $\Lambda_{1\dots N,B}^{\mathsf{C}}$ via

$$\Lambda_{1\dots N,B}^{\mathsf{C}}(\cdot) = \mathrm{tr}_{E_{\mathrm{out}}}\{V \otimes \mathbf{1}_{E'}(\mathrm{tr}_{B_{\mathrm{in}}}(\cdot) \otimes |R\rangle \langle R|_{E,E'})V^{\dagger} \otimes \mathbf{1}_{E'}\}.$$

Furthermore, the unitary V can be decomposed into a sequence of unitaries $U_{k,E} : \mathcal{H}_{m} \otimes \mathcal{H}_{1}^{E} \to \mathcal{H}_{d} \otimes \mathcal{H}_{2}^{E}$ for appropriately chosen Hilbert spaces $\mathcal{H}_{1}^{E}$ and $\mathcal{H}_{2}^{E}$, where for any given permutation π of the set $\{1, \dots, N\}$, we have that

$$V = U_{\pi(1),E}^{\pi} U_{\pi(2),E}^{\pi} \dots U_{\pi(N),E}^{\pi},$$

where $U_{k,E}^{\pi}$ is not necessarily the same as $U_{k,E}^{\pi'}$ for two different permutations π and π'.

A pictorial representation of this theorem for $N = 2$ is given in figure 6.

Given this characterization of the set of non-signalling assemblages, we now turn to the set of quantum assemblages.

Proposition 13. *An assemblage $\{\sigma_{a_1 \dots a_N | x_1 \dots x_N}\}$ is quantum if and only if there exists a localizable channel $\Lambda_{1\dots N,B}^{\mathsf{Q}} : \mathcal{L}(\mathcal{H}_{m}^{\otimes N} \otimes \mathcal{H}_{B_{\mathrm{in}}}) \to \mathcal{L}(\mathcal{H}_{d}^{\otimes N} \otimes \mathcal{H}_{B_{\mathrm{out}}})$ such that the assemblage is channel-defined by $\Lambda_{1\dots N,B}^{\mathsf{Q}}$.*

In figure 7 we give a pictorial representation of a channel-defined quantum assemblage. At this point we should point out the following corollary of this proposition along with the previous theorem, which was first proven by GHJW. We note that our proof is structurally very different from the previous proofs, and is a simple

consequence of the fact that, for $N = 1$, the unitary V in theorem 12 acts only on the input Hilbert space of the untrusted party and the ancillary register. The full proof of this corollary can be found in appendix F.

Corollary 14. *For $N = 1$, all non-signalling assemblages are also quantum assemblages.*

It is important to note that this is only true for the case of $N = 1$, i.e. a single untrusted party. In section 3.3, we use causal channels to give examples of post-quantum steering, i.e. non-signalling assemblages that are not quantum. We know that post-quantum correlations witness a non-localizable channel, then any assemblage that gives post-quantum correlations must have an associated non-localizable channel. However, there exist non-quantum assemblages that will never give rise to non-quantum correlations [18]: there are assemblages that cannot be channel-defined by a localizable channel, but for any measurement made on the Bob's system the corresponding Bell correlations are channel-defined by a local channel. This highlights that post-quantum steering is distinct from post-quantum non-locality, and indeed from non-locality itself.

For $N \geqslant 2$, as pointed out in [18], characterizing the set of quantum assemblages is difficult, and at least as hard as characterizing the set of quantum correlations. However, the almost quantum assemblages are a superset of the quantum assemblages, and for the former there is a characterization in terms of a semi-definite programme. In the next result we give a *physical interpretation* for the almost quantum assemblages.

Theorem 15. *An assemblage $\{\sigma_{a_1 \ldots a_N | x_1 \ldots x_N}\}$ is almost quantum if and only if there exists an almost localizable channel $\Lambda_{1 \ldots N,B}^{\tilde{Q}} : \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{in}}) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{out}})$ such that the assemblage is channel-defined by $\Lambda_{1 \ldots N,B}^{\tilde{Q}}$.*

The full proof is in appendix G, but is essentially a consequence of the following lemma, which is also proven in appendix G. Given this lemma, one can essentially use the proof of proposition 6 to obtain the result in theorem 15.

Lemma 16. *An assemblage $\{\sigma_{a_1 \ldots a_N | x_1 \ldots x_N}\}$ is almost quantum if and only if there exists a Hilbert space $\mathcal{H} \cong \mathcal{K} \otimes \mathcal{H}_B$, quantum state $|\psi\rangle \in \mathcal{H}$, and projective measurements $\{\Pi_{a_j | x_j} \in \mathcal{L}(\mathcal{K})\}$ for each jth party where $\sum_{a_j} \Pi_{a_j | x_j} = \mathbf{1}$ and for all permutations π of $\{1, \ldots, N\}$, $\prod_{j=1}^{N} \Pi_{a_{\pi(j)} | x_{\pi(j)}} |\psi\rangle = \prod_{j=1}^{N} \Pi_{a_j | x_j} |\psi\rangle$, such that*

$$\sigma_{a_1 \ldots a_N | x_1 \ldots x_N} = \mathrm{tr}_{\mathcal{K}} \left\{ \prod_{j=1}^{N} \Pi_{a_j | x_j} \otimes \mathbf{1}_B |\psi\rangle \langle \psi| \right\}.$$

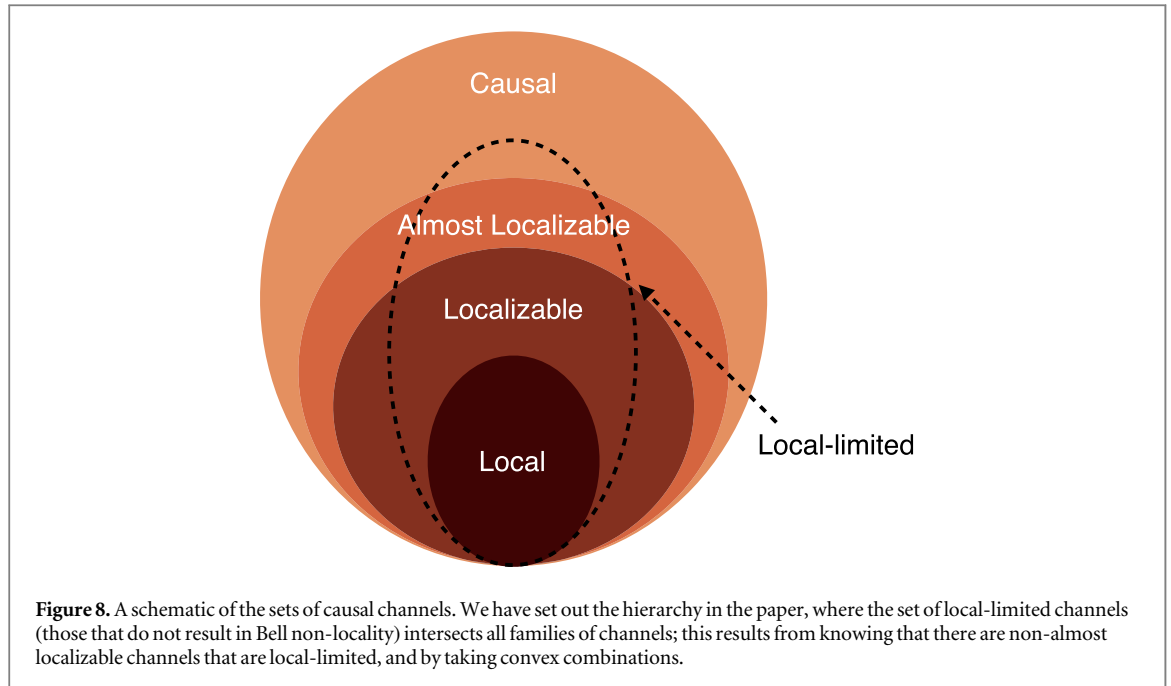### 3.2. Connections between channels and assemblages

In section 2.2, we indicated the general way to obtain correlations given a channel, and then we gave a canonical way of constructing a channel from correlations. In this section, we will do exactly the same for the case of steering.

In analogy with the case of Bell non-locality, we will first describe a general way to generate an assemblage from a channel. As in the case of Bell non-locality, the N untrusted parties can prepare a state $\rho_{x_j}^j \in \mathcal{D}(\mathcal{H}_{in}^j \otimes \mathcal{H}_{aux}^j)$ indexed by their input x_j for $j \in \{1, \ldots, N\}$, put one of its subsystems (living in $\mathcal{H}_{in}^j$) into the channel, and jointly measure the output of the channel and other subsystems associated with initial state $\rho_{x_j}^j$. The measurements then are the operators $M_{a_j} \in \mathcal{L}(\mathcal{H}_{out}^j \otimes \mathcal{H}_{aux}^j)$, which have outcomes a_j. The novelty in steering is the trusted party, and there is a potential ambiguity in how to generate an assemblage from a channel with N input port and N output ports. We could restrict to channels that trace out the input of the trusted party (or, equivalently, there is no input port), or the trusted party just always inputs the same quantum state into the channel. The second approach is more general when one considers the possibility that the trusted party has an auxiliary system with Hilbert space $\mathcal{H}_{B_{aux}}$, and prepares the state $\sigma_B \in \mathcal{D}(\mathcal{H}_{B_{in}} \otimes \mathcal{H}_{B_{aux}})$; there could be correlations between the input system and auxiliary system that would be destroyed by tracing out the input system. This more general approach results in the assemblage being a set of operators that act on the Hilbert space $\mathcal{H}_{B_{out}} \otimes \mathcal{H}_{B_{in}}$, and is in the spirit of *channel steering* [32], which we touch upon later.

To summarize this discussion, given a causal channel $\Lambda_{1 \ldots N,B}$, each jth untrusted party will prepare the state $\rho_{x_j}^j \in \mathcal{D}(\mathcal{H}_{in}^j \otimes \mathcal{H}_{aux}^j)$, and obtain measurement outcomes corresponding to the operators $M_{a_j} \in \mathcal{L}(\mathcal{H}_{out}^j \otimes \mathcal{H}_{aux}^j)$. The trusted party with Hilbert space $\mathcal{H}_B$, will prepare the state $\sigma_B \in \mathcal{D}(\mathcal{H}_{B_{in}} \otimes \mathcal{H}_{B_{aux}})$, and thus generate assemblage elements $\sigma_{a_1 \ldots a_N | x_1 \ldots x_N} \in \mathcal{D}(\mathcal{H}_{B_{out}} \otimes \mathcal{H}_{B_{aux}})$, which can be obtained as

$$\sigma_{a_1 \ldots a_N | x_1 \ldots x_N} = \mathrm{tr}_{1 \ldots N} \left\{ \bigotimes_{j=1}^{N} M_{a_j} \otimes \mathbf{1}_{B_{out}} \left(\Lambda_{1 \ldots N,B} \otimes \mathbf{1}_{aux} \left(\bigotimes_{j=1}^{N} \rho_{x_j} \otimes \sigma_B \right) \right) \right\}, \tag{7}$$

where $\mathbf{1}_{aux}$ is the identity operator acting on all Hilbert spaces $\mathcal{H}_{aux}^j$.

Figure 8. A schematic of the sets of causal channels. We have set out the hierarchy in the paper, where the set of local-limited channels (those that do not result in Bell non-locality) intersects all families of channels; this results from knowing that there are non-almost localizable channels that are local-limited, and by taking convex combinations.

Let us now move on to the case of constructing a generic channel from an assemblage. That is, given an assemblage with elements $\sigma_{a_1 \dots a_N | x_1 \dots x_N} \in \mathcal{D}(\mathcal{H}_B)$, we specify a canonical channel $\Sigma^c_{1 \dots N} : \mathcal{L}(\mathcal{H}^{\otimes N}_m \otimes \mathcal{H}_{B_{in}}) \to \mathcal{L}(\mathcal{H}^{\otimes N}_d \otimes \mathcal{H}_{B_{out}})$, with $\mathcal{H}_{B_{out}} = \mathcal{H}_B$, that will reproduce that assemblage, given appropriate choices of preparations and measurements. This canonical channel is defined as

$$\Sigma^c_{1 \dots N, B}(\cdot) = \sum_{x_1, \dots, x_N} \sum_{a_1, \dots, a_N} |a_1 \dots a_N\rangle \langle x_1 \dots x_N | \text{tr}_{B_{in}}(\cdot) | x_1 \dots x_N\rangle \langle a_1 \dots a_N| \otimes \sigma_{a_1 \dots a_N | x_1 \dots x_N},$$

and can be seen as a channel which completely decoheres the input and output systems with respect to a basis, traces out the trusted party's input, and then produces assemblage elements in the trusted party's output of the channel. Notice that the assemblage elements $\sigma_{a_1 \dots a_N | x_1 \dots x_N}$ are channel-defined by $\Sigma^c_{1 \dots N, B}$, as long as appropriate elements of an orthonormal basis are chosen.

The channel $\Sigma^c_{1 \dots N, B}$ can moreover be used to generate correlations, and not just assemblages. This is done by the method outlined in section 2.2, where the correlations are obtained as

$$p(a'_1, \dots, a'_N, a'_B | x'_1, \dots, x'_N, x'_B) = \text{tr}\left\{ \bigotimes_{j=1}^N M_{a'_j} \otimes M_{a'_B} \left(\Sigma^c_{1 \dots N, B} \otimes \mathbf{1}_{aux} \left(\bigotimes_{j=1}^N \rho_{x'_j} \otimes \rho_{x'_B} \right) \right) \right\},$$

from the local measurements $M_{a'_j}$ and states $\rho_{x'_j}$, where x'_B and a'_B represent the trusted party's inputs and outputs, respectively. We can now ask when this channel gives non-local correlations, or conversely, when is a channel $\Sigma^c_{1 \dots N, B}$ local-limited. The following result addresses this, and is proven in section E.3 of the appendix.
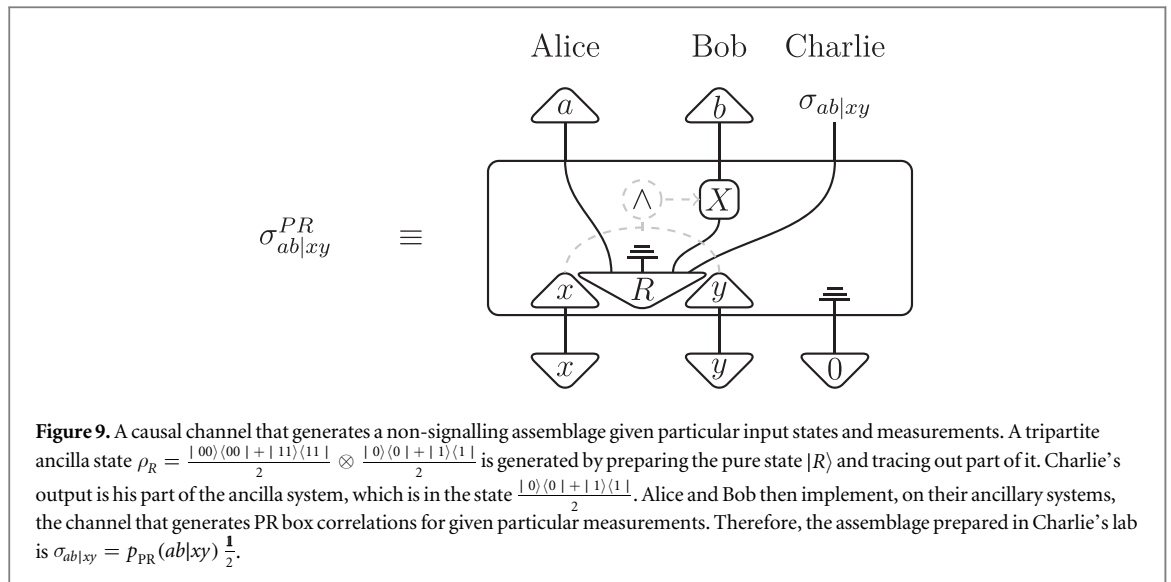
Proposition 17. *Given $\Sigma^c_{1 \dots N, B}(\cdot)$ from assemblage elements $\sigma_{a_1 \dots a_N | x_1 \dots x_N}$, this channel is local-limited if for all measurements $P_{a_B | x_B} \in \mathcal{L}(\mathcal{H}_B)$ indexed by the choice x_B and outcomes a_B, the correlations $p(a_1, \dots, a_N, a_B | x_1, \dots, x_N, x_B) := \text{tr}\{P_{a_B | x_B} \sigma_{a_1 \dots a_N | x_1 \dots x_N}\}$ are local.*

A direct consequence of this result is that the canonical channel that one would construct for the post-quantum assemblage given in [18] is a local-limited, yet non-localizable channel. Furthermore, this channel is actually not even almost localizable [18]. We summarize all of these observations in figure 8.

One can define moreover the set of channels restricted to producing only quantum correlations, and call them the *quantum-limited* channels, where these correlations can be non-local, therefore defining a larger set than the set of local-limited channels. We can then take, for instance, the post-quantum assemblages from [33] that can result in non-local but quantum correlations, and from their canonical channels give quantum-limited channels that are not almost localizable.

3.3. Examples of post-quantum steering

In this section we have outlined a constructive way to understand post-quantum steering: assemblages that cannot be channel-defined by localizable channels. We give a couple of examples of post-quantum steering that are a simple consequence of theorem 12.

Figure 9. A causal channel that generates a non-signalling assemblage given particular input states and measurements. A tripartite ancilla state $\rho_R = \frac{|00\rangle\langle 00| + |11\rangle\langle 11|}{2} \otimes \frac{|0\rangle\langle 0| + |1\rangle\langle 1|}{2}$ is generated by preparing the pure state $|R\rangle$ and tracing out part of it. Charlie's output is his part of the ancilla system, which is in the state $\frac{|0\rangle\langle 0| + |1\rangle\langle 1|}{2}$. Alice and Bob then implement, on their ancillary systems, the channel that generates PR box correlations for given particular measurements. Therefore, the assemblage prepared in Charlie's lab is $\sigma_{ab|xy} = p_{\mathrm{PR}}(ab|xy)\frac{1}{2}$.

The first example of post-quantum steering is depicted in figure 9, and is in a tripartite scenario where Alice and Bob steer Charlie, whose Hilbert space has dimension $d_C = 2$. There, an ancilla is initialized on state $\rho_R = \frac{|00\rangle\langle 00| + |11\rangle\langle 11|}{2}_{AB} \otimes \frac{|0\rangle\langle 0| + |1\rangle\langle 1|}{2}_C$, where AB and C denote Alice and Bob's and Charlie's share of the ancilla system. Then, the part of the ancilla shared by Alice and Bob is used as the ancilla in the channel that generates PR box correlations, while the qubit on state $\frac{|0\rangle\langle 0| + |1\rangle\langle 1|}{2}$ is output by Charlie. This map is causal since it has exactly the same form as described in theorem 12 (after one locally dilates all processes to be unitary). Once that Alice and Bob input qubits in the computational basis and measure their output systems, the following assemblage elements are then prepared in Charlie's lab:

$$\sigma_{ab|xy} = p_{\mathrm{PR}}(ab|xy)\frac{1}{2}.$$

This assemblage is a non-signalling one which has no quantum realization [18]. However, note that we can have post-quantum steering without any entanglement (across any of the bipartitions) in the shared ancilla state ρ_R in the causal channel. In our next example, the ancilla in the channel does consist of entanglement, and the channel generates pure state entanglement between three parties.
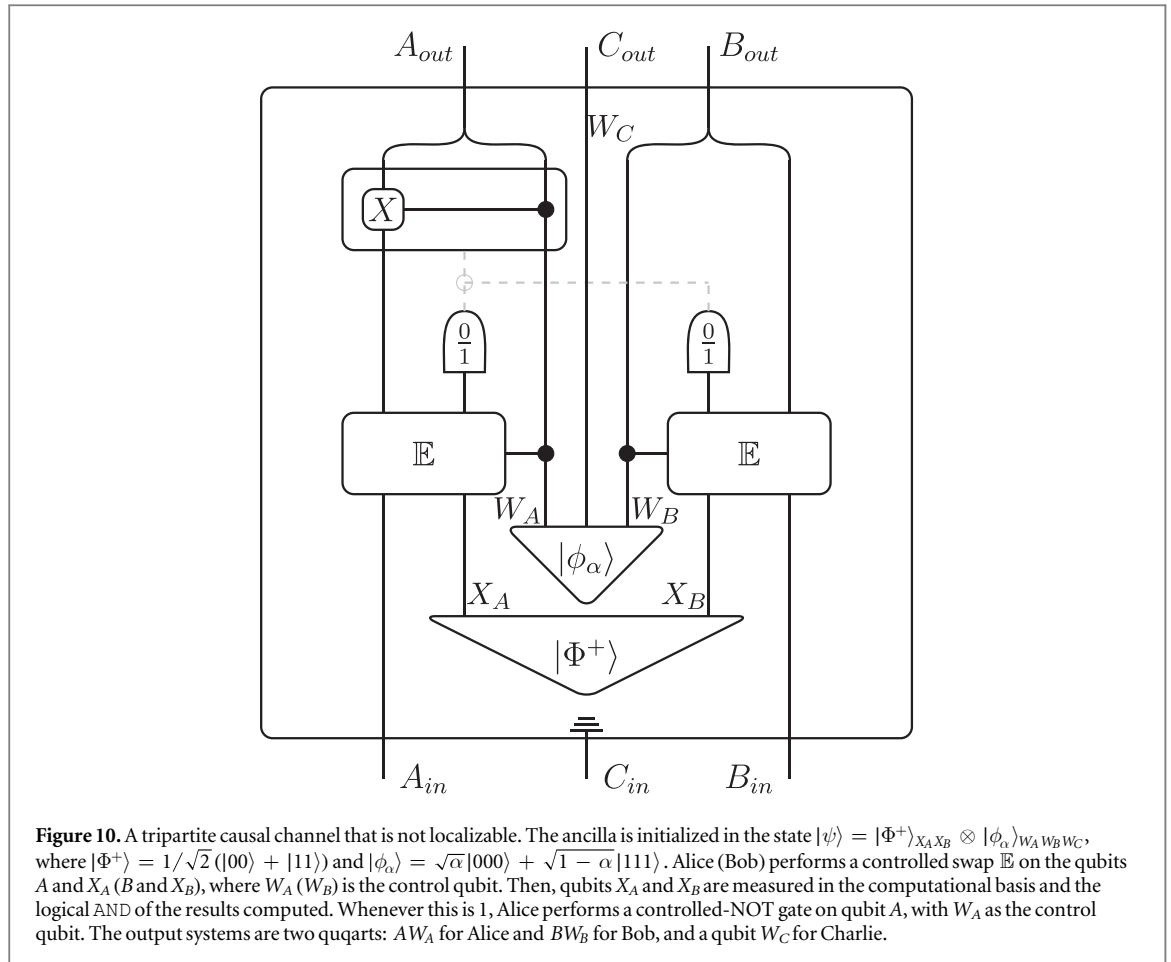
The second example of post-quantum steering also comprises a causal channel that is not localizable, and relies on the results of [26]. The steering scenario consists of Alice and Bob, who by performing two dichotomic measurements, steer Charlie, whose Hilbert space has dimension $d_C = 2$. The channel used by the three parties is depicted in figure 10. Each party's input system is given by a qubit labelled by A, B and C, respectively. Then, the channel makes use of a five qubit ancilla ($X_A W_A W_C W_B X_B$) initialized on the state:

$$|\psi\rangle = \frac{|00\rangle_{X_A X_B} + |11\rangle_{X_A X_B}}{\sqrt{2}} \otimes (\sqrt{\alpha}\,|000\rangle_{W_A W_B W_C} + \sqrt{1-\alpha}\,|111\rangle_{W_A W_B W_C}).$$

First, Alice (Bob) makes a controlled swap on the input and X_A (X_B) qubits, with W_A (W_B) being the control qubit. Then, both the X_A and X_B qubits are measured in the computational basis, and their logical AND computed. Finally, if this measurement result is 1, a controlled-NOT is performed by Alice on qubit A, with W_A as the control qubit. The output systems are then a ququart AW_A for Alice, another one BW_B for Bob, and a qubit W_C for Charlie. Since the marginal channel for Alice and Bob is causal [26], this tripartite extension is also causal. To construct an assemblage, $\{\sigma_{ab|xy}\}$ for $x, y, a, b \in \{0, 1\}$, from this channel, we prepare states $|x\rangle$ and $|y\rangle$ in the computational basis for Alice and Bob, respectively, and then measure in the computational basis, where Alice's outcome is a and Bob's is b. To check that this resulting assemblage is post-quantum, one can check that if we trace out Charlie's system, the correlations between Alice and Bob violate the CHSH inequality beyond Tsirelson's bound. Since we trace out Charlie's output system, the ancilla's state for Alice and Bob is given by

$$\rho_{\mathrm{ancilla}} = \frac{(|00\rangle_{X_A X_B} + |11\rangle_{X_A X_B})(\langle 00|_{X_A X_B} + \langle 11|_{X_A X_B})}{2} \otimes$$
$$(\alpha|00\rangle\langle 00|_{W_A W_B} + (1-\alpha)|11\rangle\langle 11|_{W_A W_B}).$$

For a choice of parameter $\alpha = \frac{1}{6}$, the correlations can be shown to give a value of 3 for the CHSH inequality, which is larger than Tsirelson's bound $2\sqrt{2}$. Therefore, the map is definitely not localizable for that choice of α. This channel can hence be used for Alice and Bob to channel-define a post-quantum assemblages on Charlie's

Figure 10. A tripartite causal channel that is not localizable. The ancilla is initialized in the state $|\psi\rangle = |\Phi^+\rangle_{X_A X_B} \otimes |\phi_\alpha\rangle_{W_A W_B W_C}$, where $|\Phi^+\rangle = 1/\sqrt{2}\,(|00\rangle + |11\rangle)$ and $|\phi_\alpha\rangle = \sqrt{\alpha}\,|000\rangle + \sqrt{1-\alpha}\,|111\rangle$. Alice (Bob) performs a controlled swap $\mathbb{E}$ on the qubits A and X_A (B and X_B), where W_A (W_B) is the control qubit. Then, qubits X_A and X_B are measured in the computational basis and the logical AND of the results computed. Whenever this is 1, Alice performs a controlled-NOT gate on qubit A, with W_A as the control qubit. The output systems are two quqarts: AW_A for Alice and BW_B for Bob, and a qubit W_C for Charlie.

subsystem. Not only this, but since almost quantum correlations cannot violate Tsirelson's bound either [25], then this assemblage is not even almost quantum, and thus the channel is not almost localizable.

Finally, we discuss how certifying the post-quantumness of the Bell correlations that are channel-defined by a causal map is not a necessary condition for such a channel to be non-localizable. For this, consider the post-quantum assemblage given in the main result of [18]. We can construct a canonical channel that is not localizable and that channel-defines this post-quantum assemblage. Now, this particular assemblage has the property that the Bell correlations it produces are quantum, or more precisely, local [18]. Hence, we can construct a provably non-localizable channel that can only channel-define local correlations in Bell scenarios.

4. Teleportation and Buscemi non-locality

Inspired by the connection between forms of non-locality and quantum channels, in this section we initiate the study of *post-quantum non-classical teleportation*, and *post-quantum Buscemi non-locality*. Non-classical teleportation [24] and Buscemi non-locality [23] (or semi-quantum non-locality[12]) have been introduced very recently within the quantum information community as generalizations of steering and Bell non-locality, respectively. We will review each of these notions, and then relate their study to our study of channels, and this will naturally give a framework in which to study their post-quantum generalizations.

4.1. Buscemi non-locality
The pioneering work by Buscemi consisted in defining a semi-quantum non-local game and arguing that any entangled state is more useful than a separable one for winning at it [23]. It should be noted that the kind of game Buscemi describes is subtly distinct to the one hinted by Leung, Toner and Watrous [34]. In this section, we will study the kind of non-locality that is witnessed in these games, and we begin by presenting the general setup.

Consider N parties, each of which has a quantum system with Hilbert space $\mathcal{K}_j$ and can prepare it in one out of m quantum states. For each $j \leqslant N$, the states in which party j may prepare their system are $\rho_{x_j} \in \mathcal{D}(\mathcal{K}_j)$, with $x_j \in \{1, \ldots, m\}$ being the classical label of the particular preparation. The parties then locally plug the system

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[12] We do not use this terminology so as not to confuse between semi-quantum and post-quantum.

into some device (it can be a black box in analogy with Bell non-locality), and then receive a classical output from the device. Let $a_j \in \{1, \ldots, d\}$ denote the classical output for the jth party, where d is the total number of possible outputs the device can locally produce.

Effectively, this whole process just described is a measurement on the preparations made by the N parties. By means of a set of tomographically-complete preparations at each site, the parties can hence generate a description of this measurement. For convenience, we now introduce a new piece of terminology to describe this measurement.

Definition 18. In a Buscemi non-locality experiment, for a set of classical outputs $\{(a_1, \ldots, a_N)\}$, a distributed measurement is $\{M_{a_1, \ldots, a_N} \in \mathcal{L}(\bigotimes_{j=1}^{N} \mathcal{K}_j)\}_{a_1, \ldots, a_N}$ where $M_{a_1, \ldots, a_N} \geqslant 0$ and $\sum_{a_1, \ldots, a_N} M_{a_1, \ldots, a_N} = \mathbb{1}_{1, \ldots, N}$.

Given this distributed measurement, it is straightforward to generate conditional probabilities from its elements and certain state preparations $\{\rho_{x_j}\}$ as

$$p(a_1, \ldots, a_N | x_1, \ldots, x_N) = \mathrm{tr}\{M_{a_1, \ldots, a_N} \rho_{x_1} \otimes \ldots \otimes \rho_{x_N}\}.$$

For the purposes of Buscemi's original work, we need to define the set of distributed measurements that result from the set of local operations and shared randomness, which we call the *local distributed measurements*.

Definition 19. A distributed measurement is local if there exist N auxiliary systems R_j for $j \in \{1, \ldots, N\}$ with Hilbert spaces $\bigotimes_{j=1}^{N} \mathcal{H}_{R_j}$ such that

$$M_{a_1, \ldots, a_N} = \mathrm{tr}_{R_1, \ldots, R_N}\left\{\rho_{R_1, \ldots, R_N} \bigotimes_{j=1}^{N} \Pi_{a_j}\right\},$$

where $\{\Pi_{a_j} \in \mathcal{L}(\mathcal{K}_j \otimes \mathcal{H}_{R_j})\}_{a_j}$ is a complete projective measurement, and $\rho_{R_1, \ldots, R_N} \in \mathcal{D}(\bigotimes \mathcal{H}_{R_j})$ is a separable state, i.e. for $|\phi_\lambda^j\rangle \in \mathcal{H}_{R_j}$

$$\rho_{R_1, \ldots, R_N} = \sum_\lambda p_\lambda |\phi_\lambda^1\rangle\langle\phi_\lambda^1|_{R_1} \otimes |\phi_\lambda^2\rangle\langle\phi_\lambda^2|_{R_2} \otimes \ldots \otimes |\phi_\lambda^N\rangle\langle\phi_\lambda^N|_{R_N},$$

with $p_\lambda \geqslant 0$ and $\sum_\lambda p_\lambda = 1$.

Without loss of generality the local measurements can be taken to be projective, since the dimension of the Hilbert spaces R_j is finite, but not constrained. Clearly, the state $\rho_{R_1, \ldots, R_N}$ could, in principle, be entangled, and thus we now define the set of *quantum distributed measurements*.

Definition 20. A distributed measurement is quantum if there exist N auxiliary systems R_j for $j \in \{1, \ldots, N\}$ with Hilbert spaces $\bigotimes_{j=1}^{N} \mathcal{H}_{R_j}$ such that

$$M_{a_1, \ldots, a_N} = \mathrm{tr}_{R_1, \ldots, R_N}\left\{\rho_{R_1, \ldots, R_N} \bigotimes_{j=1}^{N} \Pi_{a_j}\right\},$$

where $\{\Pi_{a_j} \in \mathcal{L}(\mathcal{K}_j \otimes \mathcal{H}_{R_j})\}_{a_j}$ is a complete projective measurement, and $\rho_{R_1, \ldots, R_N} \in \mathcal{D}(\bigotimes \mathcal{H}_{R_j})$ is any quantum state, entangled or otherwise.

The main result of Buscemi in [23] can then be restated as: for every non-separable state $\rho_{R_1, \ldots, R_N}$, there exists a set of projective measurements $\{\Pi_{a_j} \in \mathcal{L}(\mathcal{K}_j \otimes \mathcal{H}_{R_j})\}_{a_j}$ such that the distributed measurement is not local. A corollary of this is that the set of local distributed measurements is strictly contained in the set of quantum distributed measurements.

In complete analogy with the study of Bell non-locality and steering, we can ask what are the most general distributed measurements that do not permit superluminal signalling. The following definition formalizes the answer to this.

Definition 21. Given a bipartition $S_1 \cup S_2 = \{1, \ldots, N\}$ of N parties where $S_1 = \{i_1, \ldots, i_s\}$ and $S_2 = \{i_{s+1}, \ldots, i_N\}$, a distributed measurement $\{M_{a_1 \ldots a_N}\}$ does not permit signalling across this bipartition if there exist sets of complete measurements $\{M_{a_{i_1}, \ldots, a_{i_s}} \in \mathcal{L}(\mathcal{K}_{i_1} \otimes \ldots \otimes \mathcal{K}_{i_s})\}$ and $\{M_{a_{i_{s+1}}, \ldots, a_{i_N}} \in \mathcal{L}(\mathcal{K}_{i_{s+1}} \otimes \ldots \otimes \mathcal{K}_{i_N})\}$ such that

$$\sum_{a_{i_1}, \ldots, a_{i_s}} M_{a_{i_1}, \ldots, a_{i_s}} = \mathbb{1}_{i_1, \ldots, i_s} \tag{8}$$

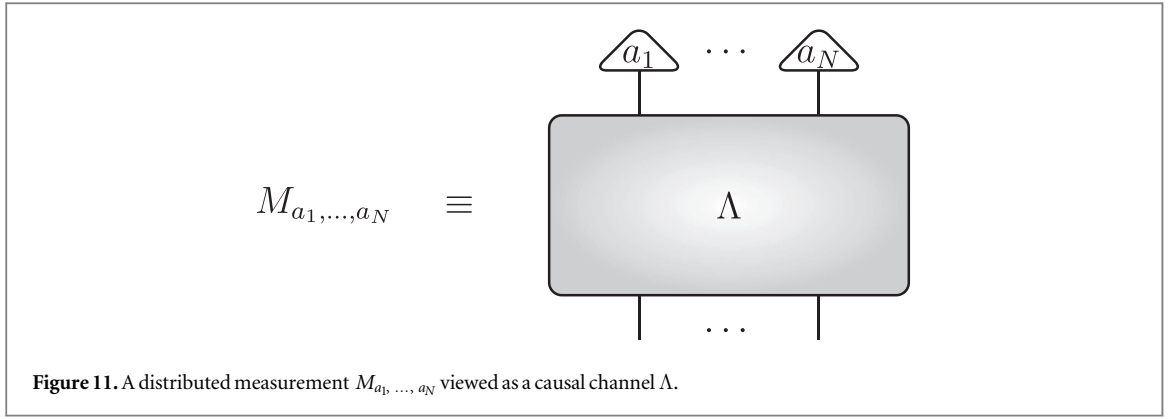$$\sum_{a_{i_{s+1}}, \ldots, a_{i_N}} M_{a_{i_{s+1}}, \ldots, a_{i_N}} = \mathbb{1}_{i_{s+1}, \ldots, i_N} \tag{9}$$

Figure 11. A distributed measurement $M_{a_1, \ldots, a_N}$ viewed as a causal channel Λ.

$$\sum_{a_{i_{s+1}}, \ldots, a_{i_N}} M_{a_1, \ldots, a_N} = M_{a_{i_1}, \ldots, a_{i_s}} \otimes \mathbf{1}_{i_{s+1}, \ldots, i_N} \tag{10}$$

$$\sum_{a_{i_1}, \ldots, a_{i_s}} M_{a_1, \ldots, a_N} = \mathbf{1}_{i_1, \ldots, i_s} \otimes M_{a_{i_{s+1}}, \ldots, a_{i_N}}. \tag{11}$$

A distributed measurement $\{M_{a_1 \ldots a_N}\}$ belongs to the set of non-signalling distributed measurements if and only if it does not permit signalling across any bipartition of the N parties.

If a distributed measurement is non-signalling but not quantum then we refer to this as *post-quantum Buscemi non-locality*. We are not the first to describe the set of non-signalling distributed measurements, Šupić *et al* [35] defined this set in the bipartite setting, although the terminology 'distributed measurement' is of our creation. We believe we are, however, the first to point out the possibility of post-quantum Buscemi non-locality. Indeed, in the next section we point this out in a clear fashion.

4.2. Buscemi non-locality via quantum channels

In this section we take our channels-based perspective and apply it to the study of Buscemi non-locality. This indeed proceeds similarly to the study of steering and Bell non-locality. The Buscemi non-locality scenario consists of N parties, where party jth (for each $j \leqslant N$) acts on the Hilbert space $\mathcal{K}_j$, and outputs data $a_j \in \{1, \ldots, d\}$. To study such a Buscemi scenario, consider N parties to have input Hilbert spaces $\mathcal{H}_{\text{in}}^j = \mathcal{K}_j$, and output Hilbert spaces $\mathcal{H}_{\text{out}}^j = \mathcal{H}_d$ for all j, where $\mathcal{H}_d$ has dimension d. Denote by $\{|a\rangle\}_{a=1:d}$ an orthonormal basis of $\mathcal{H}_d$. We now consider channels $\Lambda_{1 \ldots N} : \mathcal{L}(\bigotimes_j \mathcal{H}_{in}^j) \to \mathcal{L}(\mathcal{H}_d^{\otimes N})$ and relate channels of this form to distributed measurements.

Definition 22. A distributed measurement $\{M_{a_1, \ldots, a_N} \in \mathcal{L}(\bigotimes_{j=1}^N \mathcal{K}_j)\}$ is channel-defined if there exists a channel $\Lambda_{1 \ldots N} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j) \to \mathcal{L}(\bigotimes_{j=1}^n \mathcal{H}_{\text{out}}^j = \mathcal{H}_d^{\otimes N})$, and some choice of orthonormal bases $\{|a_j\rangle \in \mathcal{H}_d\}_{a=1:d}$ for each party, such that

$$M_{a_1, \ldots, a_N} = \Lambda_{1 \ldots N}^\dagger \left(\bigotimes_{k=1}^N |a_k\rangle \langle a_k| \right),$$

where $\Lambda_{1 \ldots N}^\dagger$ is the dual of $\Lambda_{1 \ldots N}$.

Figure 11 presents a pictorial representation of distributed measurements as quantum channels. Given this definition, as before, we can now give alternative definitions of local, quantum, and non-signalling distributed measurements.

Proposition 23. *A distributed measurement is local if and only if there exists a local channel* $\Lambda_{1 \ldots N}^{\mathsf{L}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j) \to \mathcal{L}(\mathcal{H}_d^{\otimes N})$ *such that the distributed measurement is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{L}}$.

Proposition 24. *A distributed measurement is quantum if and only if there exists a localizable channel* $\Lambda_{1 \ldots N}^{\mathsf{Q}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j) \to \mathcal{L}(\mathcal{H}_d^{\otimes N})$ *such that the distributed measurement is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{Q}}$.

Proposition 25. *A distributed measurement is non-signalling if and only if there exists a causal channel* $\Lambda_{1 \ldots N}^{\mathsf{C}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j) \to \mathcal{L}(\mathcal{H}_d^{\otimes N})$ *such that the distributed measurement is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{C}}$.

Given these definitions of distributed measurements, it is straightforward to see that if each party were to prepare pure states from an orthonormal basis, then we recover the Bell non-locality setting. This then implies that local, quantum, and post-quantum non-locality implies a local, quantum, and post-quantum distributed measurement. A simple consequence of this is that the set of non-signalling distributed measurements is strictly larger than the set of quantum distributed measurements. For example, we can take the channel that produces the PR box correlations, and generate a post-quantum distributed measurement. However, do there exist post-quantum distributed measurements that will never produce post-quantum correlations? We leave this question as open, but in this direction we now define the set of almost quantum distributed measurements as the analogue of almost quantum correlations and assemblages.

Definition 26. A distributed measurement is almost quantum if there exists an almost localizable channel $\Lambda_{1\ldots N}^{\tilde{Q}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j) \to \mathcal{L}(\mathcal{H}_d^{\otimes N})$ such that the distributed measurement is channel-defined by $\Lambda_{1\ldots N}^{\tilde{Q}}$.

Given that the set of almost quantum correlations is larger than the set of quantum correlations, it follows that the set of almost quantum distributed measurements is larger than the set of quantum distributed measurements. In future work we will investigate whether this set has a useful characterization in terms of semi-definite programming. Given such a characterization, we should be able to address the question of whether post-quantum Buscemi non-locality implies post-quantum Bell non-locality.

4.3. Non-classical teleportation

The final scenario we consider is a generalization of the steering scenario as first outlined by Cavalcanti *et al* [24]. In this scenario the original motivation was to consider two parties, and have one party 'teleport' quantum information to the other, even if their resources are noisy. In particular, Alice is given one out of m possible quantum states ρ_j for $j \in \{1, \ldots, m\}$, and produces some classical data (using a measurement on this input state and some other shared resource with Bob), and Bob has a quantum system upon which he can perform state tomography. Importantly, the set of states $\{\rho_j\}$ is known to all parties, unlike in conventional teleportation where there is a single unknown state that is to be teleported. Once Bob knows the choice of state ρ_j and the classical data (that resulted from Alice's measurement), Bob can deduce their (subnormalised) state conditioned upon this information. This is analogous to the assemblage in a steering scenario, which is a collection of (subnormalised) states conditioned on the classical information generated by the untrusted party. In the case where Alice and Bob share a maximally entangled state, Alice can make an entangled measurement on her input state ρ_j and her half of this maximally entangled state. Conditioned on the outcome of the measurement, the state in Bob's laboratory will be ρ_j with some unitary applied that depends on the outcome. In general, given Bob's conditional (subnormalised) quantum state, they wish to establish if 'non-classical teleportation' took place.

We now extend this scenario to mimic closer the case of multipartite steering experiments. Consider N parties, each of which has a quantum system with Hilbert space $\mathcal{K}_j$ and can prepare it in one out of m quantum states. For each $j \leqslant N$, the states in which party j may prepare their system are $\rho_{x_j} \in \mathcal{D}(\mathcal{K}_j)$, with $x_j \in \{1, \ldots, m\}$ being the classical label of the particular preparation. In addition, consider another party, Bob, who has a quantum system with Hilbert space $\mathcal{K}_B$ and can perform quantum state tomography on his part of his system. The first N parties generate classical data locally from their system, and we denote by $a_j \in \{1, \ldots, d\}$ the classical output obtained by party j, for each $j \leqslant N$.

Since the first N parties could prepare their input system in an arbitrary state before plugging it into their unknown device, they could each choose states from a tomographically-complete set of states, as with the Buscemi non-locality scenario. That is, enough states that span the space $\mathcal{D}(\mathcal{K}_j)$. The difference now in this scenario from the Buscemi non-locality scenario is that we have Bob's quantum system with Hilbert space $\mathcal{K}_B$, and upon which he can perform any quantum operation he likes. Therefore if we consider the whole process in terms of known quantum systems, we have the input Hilbert spaces $\mathcal{K}_j$ and an 'output' Hilbert space $\mathcal{K}_B$ in Bob's laboratory. Therefore the process of producing classical data and a (subnormalised) quantum state in Bob's laboratory can be described in terms of an object, which we call a *teleportage*. This object can be characterized as a map from space of operators over $\bigotimes_{j=1}^N \mathcal{K}_j$ to the space of operators on $\mathcal{K}_B$, and it is characterized by the fact that a tomographically-complete set of input states can be generated, and a tomographically-complete measurement can be made on Bob's system.

Definition 27. For a set of classical outputs $\{(a_1, \ldots, a_N)\}$ and the Hilbert space $\mathcal{K}_B$ of Bob's system, a teleportage is an instrument $\{T_{a_1, \ldots, a_N} \in \mathcal{L}(\bigotimes_{j=1}^N \mathcal{K}_j) \to \mathcal{L}(\mathcal{K}_B)\}_{a_1, \ldots, a_N}$ such that, for all quantum states $\rho \in \mathcal{D}(\bigotimes_{j=1}^N \mathcal{K}_j)$, $\text{tr}\{\sum_{a_1, \ldots, a_N} T_{a_1, \ldots, a_N}(\rho)\} = \text{tr}\{\rho\}$.

We note that one can obtain an assemblage $\{\sigma_{a_1, \ldots, a_N | x_1, \ldots, x_N} \in \mathcal{D}(\mathcal{K}_B)\}$ for any set of input quantum states $\{\rho_{x_j}\}$, in the following way:

$$\sigma_{a_1, \ldots, a_N | x_1, \ldots, x_N} = T_{a_1, \ldots, a_N}\left(\bigotimes_{j=1}^N \rho_{x_j}\right).$$

This is actually slightly distinct from the assemblages in the standard steering scenario, since the states ρ_{x_j} have some quantum information, and so the classical labels x_j do not capture everything.

As with the study of Buscemi non-locality and steering, we can define the physically meaningful sets of teleportages.

Definition 28. A teleportage is local if there exist N auxiliary systems R_j for $j \in \{1, \ldots, N\}$ with Hilbert spaces $\bigotimes_{j=1}^N \mathcal{H}_{R_j}$ such that

$$T_{a_1, \ldots, a_N} = \mathrm{tr}_{1, \ldots, N, R_1, \ldots, R_N}\left\{\rho_{R_1, \ldots, R_N, B} \bigotimes_{j=1}^N \Pi_{a_j} \otimes \mathbb{1}_B\right\},$$

where $\{\Pi_{a_j} \in \mathcal{L}(\mathcal{K}_j \otimes \mathcal{H}_{R_j})\}_{a_j}$ is a complete projective measurement, and $\rho_{R_1, \ldots, R_N, B} \in \mathcal{D}(\bigotimes \mathcal{H}_{R_j} \otimes \mathcal{K}_B)$ is a separable state, i.e. for $|\phi_\lambda^j\rangle \in \mathcal{H}_{R_j}$ and $|\phi_\lambda^B\rangle \in \mathcal{K}_B$

$$\rho_{R_1, \ldots, R_N, B} = \sum_\lambda p_\lambda |\phi_\lambda^1\rangle\langle\phi_\lambda^1|_{R_1} \otimes |\phi_\lambda^2\rangle\langle\phi_\lambda^2|_{R_2} \otimes \ldots \otimes |\phi_\lambda^N\rangle\langle\phi_\lambda^N|_{R_N} \otimes |\phi_\lambda^B\rangle\langle\phi_\lambda^B|,$$

with $p_\lambda \geq 0$ and $\sum_\lambda p_\lambda = 1$.

Definition 29. A teleportage is quantum if there exist N auxiliary systems R_j for $j \in \{1, \ldots, N\}$ with Hilbert spaces $\bigotimes_{j=1}^N \mathcal{H}_{R_j}$ such that

$$T_{a_1, \ldots, a_N} = \mathrm{tr}_{1, \ldots, N, R_1, \ldots, R_N}\left\{\rho_{R_1, \ldots, R_N, B} \bigotimes_{j=1}^N \Pi_{a_j} \otimes \mathbb{1}_B\right\},$$

where $\{\Pi_{a_j} \in \mathcal{L}(\mathcal{K}_j \otimes \mathcal{H}_{R_j})\}_{a_j}$ is a complete projective measurement, and $\rho_{R_1, \ldots, R_N, B} \in \mathcal{D}(\bigotimes \mathcal{H}_{R_j} \otimes \mathcal{K}_B)$ is any quantum state, entangled or otherwise.

Definition 30. Given a bipartition $S_1 \cup S_2 = \{1, \ldots, N\}$ of N parties where $S_1 = \{i_1, \ldots, i_s\}$ and $S_2 = \{i_{s+1}, \ldots, i_N\}$, a teleportage $\{T_{a_1, \ldots, a_N} : \mathcal{L}(\bigotimes_{j=1}^N \mathcal{K}_j) \to \mathcal{L}(\mathcal{K}_B)\}$ does not permit signalling across this bipartition if there exist further teleportages $\{T_{a_{i_1}, \ldots, a_{i_s}} : \mathcal{L}(\bigotimes_{j=1}^s \mathcal{K}_j) \to \mathcal{L}(\mathcal{K}_B)\}$, $\{T_{a_{i_{s+1}}, \ldots, a_{i_N}} : \mathcal{L}(\bigotimes_{j=s+1}^N \mathcal{K}_j) \to \mathcal{L}(\mathcal{K}_B)\}$, and quantum state $\rho_B \in \mathcal{D}(\mathcal{H}_B)$ such that

$$\sum_{a_{i_1}, \ldots, a_{i_s}} T_{a_{i_1}, \ldots, a_{i_s}}(\cdot) = \rho_B$$

$$\sum_{a_{i_{s+1}}, \ldots, a_{i_N}} T_{a_{i_{s+1}}, \ldots, a_{i_N}}(\cdot) = \rho_B$$

$$\sum_{a_{i_{s+1}}, \ldots, a_{i_N}} T_{a_1, \ldots, a_N}(\cdot) = T_{a_{i_1}, \ldots, a_{i_s}}(\cdot)$$

$$\sum_{a_{i_1}, \ldots, a_{i_s}} T_{a_1, \ldots, a_N}(\cdot) = T_{a_{i_{s+1}}, \ldots, a_{i_N}}(\cdot).$$

A teleportage $\{T_{a_1 \ldots a_N}\}$ belongs to the set of non-signalling teleportages if and only if it does not permit signalling across any bipartition of the N parties.

We say a teleportage demonstrates *post-quantum non-classical teleportation* if it is a non-signalling teleportage that is not a quantum teleportage. As far as we know, we are the first to define the set of non-signalling teleportages, in addition to introducing the nomenclature.

4.4. Non-classical teleportation via quantum channels

Finally, we look at non-classical teleportation through the lens of channels. Recall that we have a black box device with N inputs for quantum systems, and N classical outputs $a_j \in \{1, \ldots, d\}$ for $j \in \{1, \ldots, N\}$. For the N ports of the black box device we associate input and output Hilbert spaces with each party such that $\mathcal{H}_{in}^j = \mathcal{K}_j$ and $\mathcal{H}_{out}^j = \mathcal{H}_{out}^{j'} = \mathcal{H}_d$ for all $j, j' \in \{1, \ldots, N\}$, where $\mathcal{H}_d$ is a Hilbert space of dimension d. For Bob, we have an input and output Hilbert space denoted $\mathcal{H}_{in}^B$ and $\mathcal{H}_{out}^B$, respectively, where we have that $\mathcal{H}_{in}^B = \mathcal{H}_{out}^B = \mathcal{K}_B$. Therefore the channels of interest will be $\Lambda_{1 \ldots N} : \mathcal{L}(\bigotimes_j \mathcal{H}_{in}^j \otimes \mathcal{K}_B) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$, and as before we can define teleportages in terms of these channels.

Figure 12. A Teleportage $T_{a_1, \ldots, a_N}$ viewed as a causal channel Λ.
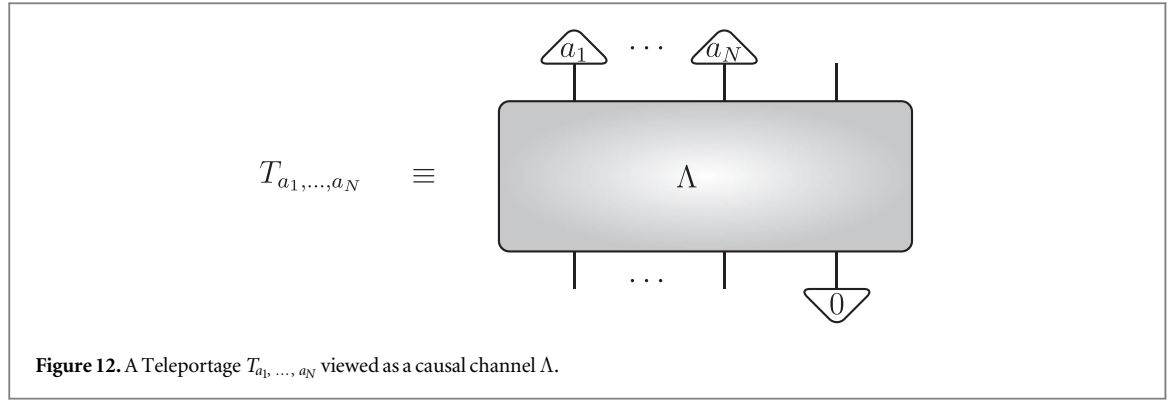
Definition 31. A teleportage $\{ T_{a_1, \ldots, a_N} : T_{a_1, \ldots, a_N} \in \mathcal{L}(\bigotimes_{j=1}^N \mathcal{K}_j) \to \mathcal{L}(\mathcal{K}_B) \}_{a_1, \ldots, a_N}$ is channel-defined if there exists a channel $\Lambda_{1 \ldots N} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j \otimes \mathcal{K}_B) \to \mathcal{L}(\bigotimes_{j=1}^N \mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$, and some choice of orthonormal bases $\{ |a_j\rangle \in \mathcal{H}_d \}_{a=1:d}$ for each party, such that

$$T_{a_1, \ldots, a_N}(\cdot) = \text{tr}_{\text{out}_1, \ldots, \text{out}_N, B} \{ \otimes_{k=1}^N |a_k\rangle \langle a_k| \, \Lambda_{1 \ldots N}(\cdot \otimes |0\rangle \langle 0|_B) \},$$

where $|0\rangle_B \in \mathcal{K}_B$.

Figure 12 depicts a teleportage as a quantum channel. Given this definition, as before we obtain the following results:

Proposition 32. *A teleportage is local if and only if there exists a local channel* $\Lambda_{1 \ldots N}^{\mathsf{L}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j \otimes \mathcal{K}_B) \to \mathcal{L}(\bigotimes_{j=1}^n \mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ *such that the teleportage is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{L}}$.

Proposition 33. *A teleportage is quantum if and only if there exists a localizable channel* $\Lambda_{1 \ldots N}^{\mathsf{Q}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j \otimes \mathcal{K}_B) \to \mathcal{L}(\bigotimes_{j=1}^n \mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ *such that the teleportage is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{Q}}$.

Proposition 34. *A teleportage is non-signalling if and only if there exists a causal channel* $\Lambda_{1 \ldots N}^{\mathsf{C}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j \otimes \mathcal{K}_B) \to \mathcal{L}(\bigotimes_{j=1}^n \mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ *such that the teleportage is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{C}}$.

It should be clear that post-quantum steering implies post-quantum non-classical teleportation, since if an assemblage is post-quantum then it is channel-defined by a non-localizable channel, this non-localizable channel will then channel-define a teleportage that is post-quantum.

For the study of steering we had an alternative characterization of non-signalling assemblages in terms of a unitary representation. This result can be generalized to the set of non-signalling teleportages as follows.

Theorem 35. *Unitary representation of non-signalling teleportages Let* $\{ T_{a_1 \ldots a_N} \}$ *be a non-signalling teleportage. Then, the teleportage is channel-defined by a channel* $\Lambda_{1 \ldots N, B}^{\mathsf{C}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j \otimes \mathcal{K}_B) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ *if and only if there exist*

- *auxiliary systems* E *and* E' *with input and output Hilbert spaces,* $\mathcal{H}_{\text{in}}^E$ *and* $\mathcal{H}_{\text{out}}^E$ *for* E*, with* $\mathcal{H}_{\text{in}}^{E'} = \mathcal{K}_B$ *and* $\mathcal{H}_{\text{out}}^{E'} = \mathcal{K}_B$ *for* E'*, that is the output Hilbert space of* E' *and* B *coincide;*

- *quantum state* $|R\rangle \in \mathcal{H}_{\text{in}}^E \otimes \mathcal{H}_{\text{in}}^{E'}$;

- *unitary operator* $V : \bigotimes_j \mathcal{H}_{\text{in}}^j \otimes \mathcal{H}_{\text{in}}^E \to \mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{\text{out}}^E$,

which produce a unitary representation of the channel $\Lambda_{1 \ldots N, B}^{\mathsf{C}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j \otimes \mathcal{K}_B) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ *via*

$$\Lambda_{1 \ldots N, B}^{\mathsf{C}}(\cdot) = \text{tr}_{E_{\text{out}}} \{ V(\text{tr}_{B_{\text{in}}}(\cdot) \otimes |R\rangle \langle R|_{E,E'}) V^\dagger \}.$$

Furthermore, the unitary V *can be decomposed into a sequence of unitaries* $U_{k,E} : \mathcal{H}_m \otimes \mathcal{H}_1^E \to \mathcal{H}_d \otimes \mathcal{H}_2^E$ *for appropriately chosen Hilbert spaces* $\mathcal{H}_1^E$ *and* $\mathcal{H}_2^E$*, where for any given permutation* π *of the set* $\{1, \ldots, N\}$*, we have that*

$$V = U_{\pi(1),E}^{\pi} U_{\pi(2),E}^{\pi} \dots U_{\pi(N),E}^{\pi},$$

where $U_{k,E}^{\pi}$ is not necessarily the same as $U_{k,E}^{\pi'}$ for two different permutations π and π'.

Given this last result about non-signalling teleportages, we can actually generalize the GHJW theorem from the case of steering to the study of non-classical teleportation.

Corollary 36. *For $N = 1$, all non-signalling teleportages are also quantum teleportages.*

That is, for the original context in which non-classical teleportation was studied, the bipartite setting, the no-signalling principle is already enough to characterize exactly *everything* that can be done quantum mechanically in the experiment.

Finally, in analogy with everything that has gone before, we can define the set of almost quantum teleportages as follows.

Definition 37. A teleportage is almost quantum if there exists an almost localizable channel $\Lambda_{1\dots N}^{\tilde{Q}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j \otimes \mathcal{K}_B) \to \mathcal{L}(\bigotimes_{j=1}^n \mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ such that the teleportage is channel-defined by $\Lambda_{1\dots N}^{\tilde{Q}}$.
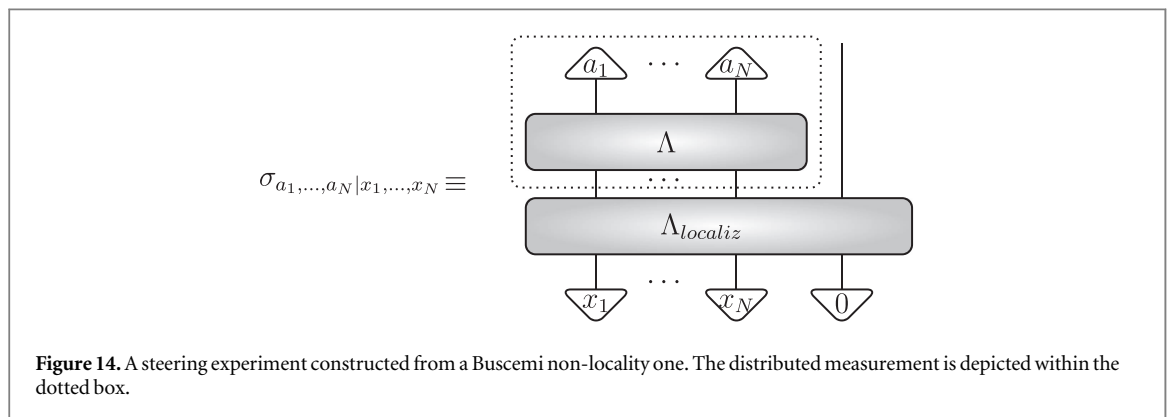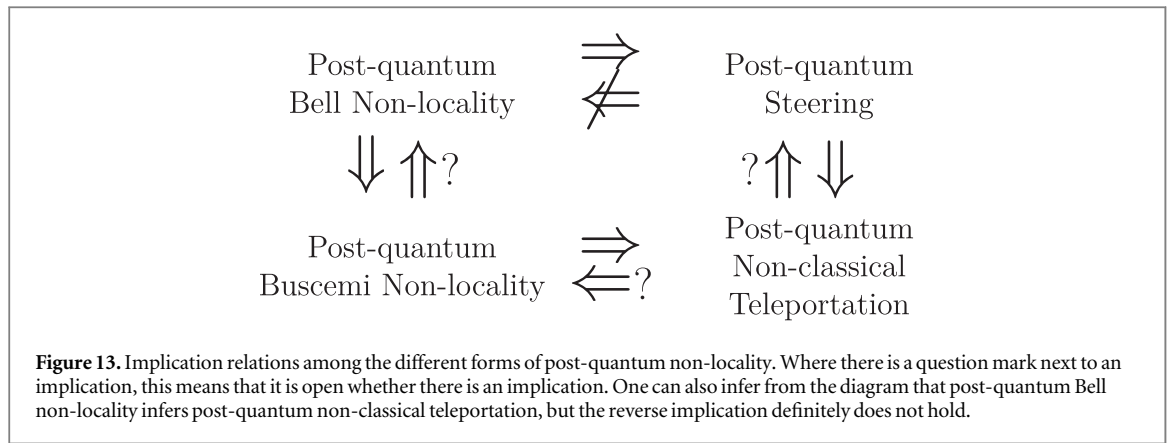
4.5. Connections between all forms of post-quantum non-locality

The relationship between entanglement, steering, and non-locality is now well-studied within the scope of quantum states. Since non-locality implies steering the non-trivial question is which entangled states demonstrate steering, but not non-locality. It has been shown that for all possible measurements on a quantum state, entanglement, steering, and non-locality are all inequivalent [30]. In post-quantum non-locality, obviously we cannot automatically associate a process with measurements on a quantum state. Furthermore, due to the GHJW theorem, post-quantum steering cannot be demonstrated when there are only two parties, although post-quantum non-locality can be demonstrated with only two parties. Therefore, the relationship between post-quantum non-locality and post-quantum steering is somewhat subtle. The resolution is, of course, to consider a steering scenario with two (or more) uncharacterized parties and then generate correlations by making a measurement on Bob's system. If these correlations demonstrate post-quantum non-locality then this implies post-quantum steering, since the whole process cannot be associated with local measurements on a quantum system. However, post-quantum steering does not imply post-quantum non-locality, as demonstrated in [18].

The relationship between post-quantum non-locality and post-quantum Buscemi non-locality was discussed at length in the previous section. In particular, if we take a distributed measurement and for a combination of local preparations of states, we obtain post-quantum correlations, then this implies post-quantum Buscemi non-locality. As mentioned above, we leave it open whether there are post-quantum distributed measurements that do not result in post-quantum non-locality for all possible preparations.

The next point to consider is the relationship between post-quantum Buscemi non-locality and post-quantum non-classical teleportation. As in the relationship between non-locality and steering, if we take a teleportage and make a measurement on Bob's system, we obtain a distributed measurement. If the distributed measurement is post-quantum, then clearly the teleportage was itself post-quantum. Likewise, one can obtain an assemblage from a teleportage by preparing certain quantum systems for each of the uncharacterized parties. If the assemblage demonstrates post-quantum steering then the teleportage was post-quantum. We see then that all these different forms of post-quantum non-locality are somehow related to each other as summarized in figure 13.

What is the relationship between post-quantum steering and post-quantum Buscemi non-locality? At first sight it seems difficult to relate the two, since in one scenario measurements are made, but preparations are made in the other. However, given our picture of non-locality from the perspective of quantum channels we can find a resolution. One way of generating an assemblage from a distributed measurement would be the following (see figure 14): encode the classical inputs $\{|x_j\rangle\}$ as elements of an orthonormal basis $\mathcal{H}_m$ for $j \in \{1, \dots, N\}$, and take a localizable channel $\Lambda : \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_B) \to \bigotimes_{j=1}^N \mathcal{K}_j \otimes \mathcal{H}_B$ where $\mathcal{K}_j$ is the Hilbert space associated with the jth party's input to a distributed measurement, and $\mathcal{H}_B$ is an auxiliary Hilbert space associated with Bob's system; apply the channel Λ to the input states $\{|x_1 x_2, \dots, x_N\rangle\}$, and then apply the distributed measurement to the systems now living in the Hilbert space $\bigotimes_{j=1}^N \mathcal{K}_j$, which results in an assemblage $\{\sigma_{a_1 \dots a_N | x_1 \dots x_N}\}$, i.e. operators acting on Bob's system. Since this extra element is a localizable channel, it will not introduce any post-quantum elements in its own right. Therefore, if we take a distributed measurement and turn it into an assemblage in this fashion, if the assemblage is post-quantum then the original distributed measurement itself was post-quantum.

Figure 13. Implication relations among the different forms of post-quantum non-locality. Where there is a question mark next to an implication, this means that it is open whether there is an implication. One can also infer from the diagram that post-quantum Bell non-locality infers post-quantum non-classical teleportation, but the reverse implication definitely does not hold.

Figure 14. A steering experiment constructed from a Buscemi non-locality one. The distributed measurement is depicted within the dotted box.

Given figure 13, we immediately see that post-quantum non-classical teleportation cannot imply post-quantum non-locality, since post-quantum steering does not imply post-quantum non-locality. That is, if post-quantum non-classical teleportation and post-quantum non-locality were equivalent then, post-quantum steering would imply post-quantum non-locality, which is not true. Furthermore, this also implies that either post-quantum Buscemi non-locality does not imply post-quantum non-locality, or post-quantum non-classical teleportation does not imply post-quantum Buscemi non-locality, or both. In figure 13 we indicate these main open questions between all forms of post-quantum non-locality with a question mark next to the implication. To prove, for example, that post-quantum Buscemi non-locality does not imply post-quantum Bell non-locality, one would need to find a distributed measurement that cannot be realized via a localizable channel, yet this channel does not give post-quantum correlations, e.g. it could be local-limited. We conjecture that all four notions of post-quantum non-locality are inequivalent.

5. Discussion

In this work we have shown that the study of post-quantum non-locality and steering can be seen as two facets of the study of quantum channels that do not permit superluminal signalling. We further showed that other scenarios can be readily approached within this scope, and hence initiated the study of post-quantum Buscemi non-locality and post-quantum non-classical teleportation. This general perspective allows us generate new examples of post-quantum steering, and allow us to generate novel kinds of non-signalling, but non-localizable channels. Furthermore, we have expanded the definition of almost quantum correlations to the domain of quantum channels (with no reference made to measurements), allowing us to recover almost quantum correlations and almost quantum assemblages in an appropriate domain.

Another channel-based perspective on the study of steering has led to so-called *channel steering* [32], as briefly mentioned in section 3.2. Channel steering is a generalization of standard bipartite steering (involving Alice and Bob), where now there is a third party, Charlie, that inputs a quantum system into a channel, and Alice and Bob have systems that are the outputs of this channel. In a sense, this channel is then a broadcast channel. Alice can perform a measurement on her system to demonstrate to Bob that she can steer his output of the channel. Channel steering is distinct from the forms of non-locality considered here, but we can extend our

channels to include this third party, and then study causal, but non-localizable, channels for post-quantum channel steering. We leave this for future work.

The characterization of quantum non-locality is not only of foundational interest, but it is also of use in quantum information theory. In particular, characterizing the set of quantum correlations is useful for device-independent quantum information, since it allows for a way to practically constrain what, say, a malicious agent can do in the preparation of devices. The study of Buscemi non-locality is of relevance to measurement-device-independent quantum information, hence this paradigm may profit from the characterization of what is quantum mechanically allowed in the setup, with direct consequences regarding randomness certification and entanglement quantification [35, 36].

One of the main open problems of this work is to further probe the relationships between the different kinds of post-quantum non-locality and steering. For example, as discussed, we know that post-quantum steering does not always imply post-quantum non-locality, but does post-quantum Buscemi non-locality imply post-quantum Bell non-locality? In figure 13 we summarized all the known relationships between all forms of post-quantum non-locality. We conjecture that all of these different notions of post-quantum non-locality are inequivalent, just as post-quantum non-locality is inequivalent to post-quantum steering.

On the way to proving our conjecture, it may be relevant to first study possible characterizations of all forms of almost quantum non-locality in terms of a semi-definite programme. Such a connection was crucial in [18] when showing that post-quantum steering does not imply post-quantum non-locality. Indeed, in this work we gave an interpretation in terms of quantum channels to the original SDP characterization of almost quantum asseblages, thus giving a physical underpinning of this set. This interpretation allowed us to generalize the notion of almost quantum nonclassicality, hence now it would be interesting to relate back these general notions to SDPs when possible.

Since we have shown that post-quantum non-locality and steering are two aspects of a more general study of quantum channels, we hope this work motivates a *resource theory of post-quantumness*. This resource theory could be approached from the point-of-view quantum channels, where the non-localizability of a channel is a resource. This relates directly to the study of zero-error communication with quantum channels [37]. Given this connection, we expect to find applications of post-quantum steering, just as we find that post-quantum non-locality can be used to trivialize communication complexity. Furthermore, we might be able to find applications of post-quantum Buscemi non-locality and post-quantum non-classical teleportation. Going further, there are other possibilities for non-locality scenarios. In particular, one can consider scenarios where all parties' outputs are quantum systems.

Our work could fit neatly within the study of quantum combs [38], quantum strategies [39], quantum causal models [40–43] and process matrices [20]. Indeed, since in certain scenarios in our work it is assumed that one party has access to a quantum system but the global system may be incompatible with quantum mechanics, it has a similar motivation to the study of indefinite causal order [20]. It would be interesting to see how our non-signalling processes interact with processes that could include signalling, and whether this interaction could be used to understand the structure of post-quantum non-locality.

Last but not least, the resource theory of non-locality has been studied by only thinking of systems as black boxes. That is, one does not need to consider Hilbert spaces, or other features of quantum mechanics, but only consider the correlations associated with particular devices. The approach in this paper has been couched in the language of quantum theory. Can we consider generalizing our framework further to consider trusted (and characterized) devices that may not be quantum, but are objects that can be described within a broad family of, say, generalized probabilistic theories [44]? Indeed, steering has already been studied within the broad framework of these theories [45, 46]. The study of non-signalling channels in general theories is left for future work and could then shed insight onto what is so special about quantum theory.

Acknowledgments

Appendix A. Relevant concepts from quantum channels

In this section we review families of multipartite quantum channels which are pertinent when discussing non-locality and steering. The general scenario we consider is that outlined in section 1.

Beckman *et al* [22] considered quantum channels in such a set-up of space-like separated laboratories, especially those channels that are compatible with relativistic causality. That is, if two parties are space-like separated, the channel mapping their input states to their output states do not permit communication between them, called the *causal channels*. There are multiple equivalent mathematical definitions capturing this concept [47], and we shall present the definition of semicausal and causal channels.

Definition 38 (Semicausal and Causal channels). Given a multipartite system and a bipartition $S_A \cup S_B = \{1, \ldots, N\}$ of its elements, a map $\Lambda : \mathcal{L}(\mathcal{H}_{\text{in}}^{S_A} \otimes \mathcal{H}_{\text{in}}^{S_B}) \to \mathcal{L}(\mathcal{H}_{\text{out}}^{S_A} \otimes \mathcal{H}_{\text{out}}^{S_B})$ is semicausal from S_B to S_A, denoted $S_B \not\rightarrow S_A$, there exists a channel $\Gamma : \mathcal{L}(\mathcal{H}_{\text{in}}^{S_A}) \to \mathcal{L}(\mathcal{H}_{\text{out}}^{S_A})$, such that for all states $\rho_{\text{in}} \in \mathcal{D}(\mathcal{H}_{\text{in}}^{S_A} \otimes \mathcal{H}_{\text{in}}^{S_B})$, $\text{tr}_{S_B}\{\Lambda[\rho_{\text{in}}]\} = \Gamma(\text{tr}_{S_B}[\rho_{\text{in}}])$. A map that is semicausal for all bipartitions is called causal.

For every channel there exists a unitary operator U acting on a system and ancilla E, such that $\Lambda[\rho] = \text{tr}_E\{U(\rho \otimes |0\rangle\langle 0|_E)U^{\dagger}\}$ for some state $|0\rangle_E \in \mathcal{H}_E$. What is the form of a unitary dilation of (semi)causal channels? For bipartite semicausal maps, works by Schumacher and Westmoreland [47], D'Ariano *et al* [26], and Piani *et al* [27] provide the following characterization:

Theorem 39 (Unitary representation for bipartite semicausal channels [26, 27, 47]). *Let Λ_{AB} be a CPTP map, with $B \not\rightarrow A$. Then, there exists an auxiliary system E with input, intermediate and output Hilbert spaces, $\mathcal{H}_{\text{in}}^E$, $\mathcal{H}_{\text{int}}^E$, and $\mathcal{H}_{\text{out}}^E$, respectively, and quantum state $|0\rangle \in \mathcal{H}_{\text{in}}^E$, producing a unitary representation of the map as $\Lambda_{AB}[\rho] = \text{tr}_E\{U(\rho \otimes |0\rangle\langle 0|_E)U^{\dagger}\}$, where the unitary U can be decomposed as a unitary acting on A and E followed by a unitary acting on B and E alone. That is,*

$$U = U_{EB}\, U_{AE},$$

for U_{EB}: $\mathcal{H}_{\text{int}}^E \otimes \mathcal{H}_{\text{in}}^B \to \mathcal{H}_{\text{out}}^E \otimes \mathcal{H}_{\text{out}}^B$ and U_{AE}: $\mathcal{H}_{\text{in}}^A \otimes \mathcal{H}_{\text{in}}^E \to \mathcal{H}_{\text{out}}^A \otimes \mathcal{H}_{\text{int}}^E$.

The statement of this theorem is depicted in figure A1. When the map is fully causal, then there exist both a decomposition in terms of $U = U_{EB}\, U_{AE}$ and one in terms of $V = V_{AE}\, V_{EB}$, where the U_j and V_j are unitaries, for $j \in \{AE, EB\}$.

This result can be generalized to multipartite causal channels. We now use notation where parties are labelled by numbers going from 1 to N. Given a multipartite causal map $\Lambda_{1\ldots N}$, there exist unitary operators $\{U_{\pi}\}_{\pi}$ acting on local systems plus a global auxiliary system E such that $\Lambda_{1\ldots N}[\rho] = \text{tr}_E\{U_{\pi}(\rho \otimes |0\rangle\langle 0|_E)U_{\pi}^{\dagger}\}$, where the unitary U_{π} has the form of $U_{\pi} = \prod_{k=1}^{N} U_{\pi(k)E}$ and π is a permutation of the parties $\{1, 2, \ldots, N\}$. The proof is presented in appendix D.1.

A particular class of causal channels is the class of *localizable channels* [22]. These are channels implemented by local operations performed by each party on their input and a share of a quantum ancilla (see figure A2). We formalize this definition below.

Definition 40 (Localizable channels). A causal channel $\Lambda_{1\ldots N}$ is localizable if there exists an N-partite ancilla system R, with Hilbert spaces $\mathcal{H}_R := \mathcal{H}_{R_1} \otimes \mathcal{H}_{R_2} \otimes \ldots \otimes \mathcal{H}_{R_N}$ with R_j labelling the jth subsystem of R, and state $\sigma_R \in \mathcal{D}(\mathcal{H}_R)$ such that, for all states $\rho \in \mathcal{D}(\bigotimes_{j=1}^{N} \mathcal{H}_{\text{in}}^j)$,

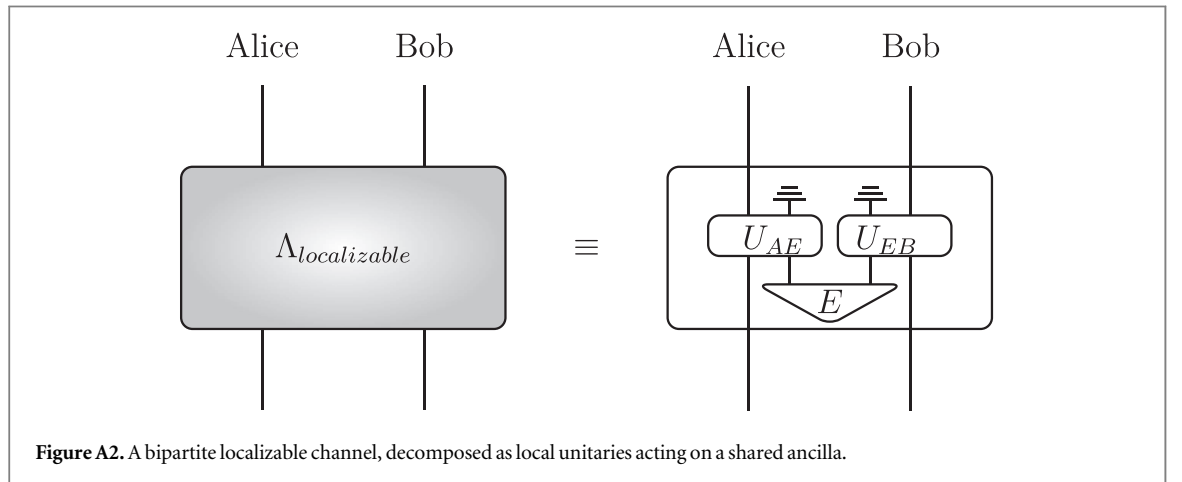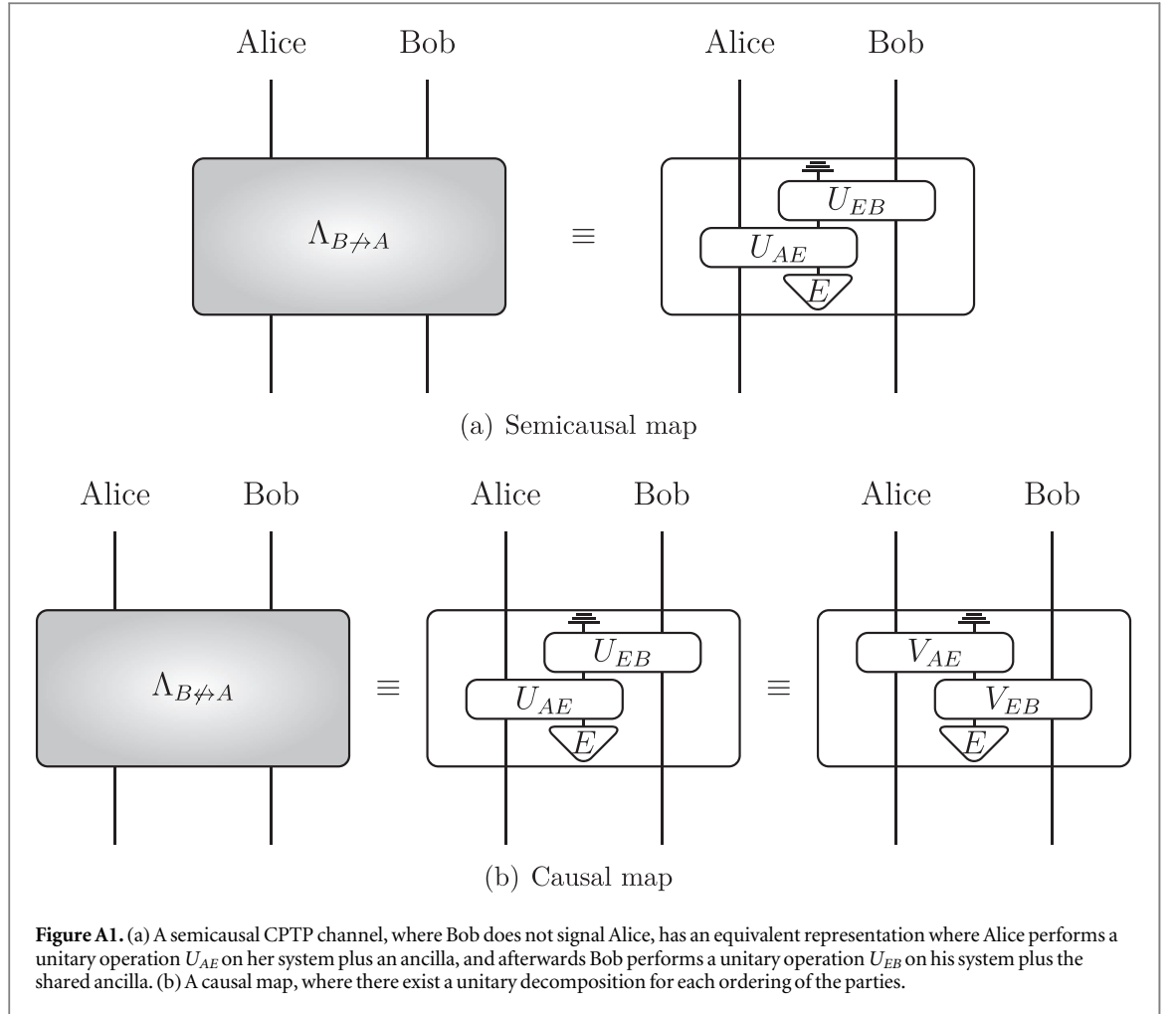$$\Lambda_{1\ldots N}[\rho] = \otimes_{k=1}^{N} \Lambda_{R_k}[\rho \otimes \sigma_R],$$

where $\Lambda_{R_k} : \mathcal{L}(\mathcal{H}_{\text{in}}^k \otimes \mathcal{H}_{R_k}) \to \mathcal{L}(\mathcal{H}_{\text{out}}^k)$.

Notice that in this definition for localizable channels, the ancilla σ_R is the same for all inputs to the channel $\Lambda_{1\ldots N}$.

It is known that, already for bipartite systems, there exist channels that are causal but not localizable [22]. Furthermore, there are examples that are not entanglement-breaking [26], unlike the example given in [22].

Just as we considered the causal channels in terms of unitaries, we can consider localizable channels in terms of unitary operators. Since there are only local maps in the localizable channels, it is straightforward to dilate each of these maps if we increase the Hilbert space dimension of local systems R_j in the ancilla. This gives the following equivalent definition of localizable channels.

Definition 41 (Unitary representation of localizable channels). A causal channel $\Lambda_{1\ldots N}$ is localizable if and only if there exists an N-partite ancilla system E with input and output Hilbert spaces

Figure A1. (a) A semicausal CPTP channel, where Bob does not signal Alice, has an equivalent representation where Alice performs a unitary operation U_{AE} on her system plus an ancilla, and afterwards Bob performs a unitary operation U_{EB} on his system plus the shared ancilla. (b) A causal map, where there exist a unitary decomposition for each ordering of the parties.

Figure A2. A bipartite localizable channel, decomposed as local unitaries acting on a shared ancilla.

$\mathcal{H}_{\text{in}}^{E} := \mathcal{H}_{\text{in}}^{E_1} \otimes \mathcal{H}_{\text{in}}^{E_2} \otimes ... \otimes \mathcal{H}_{\text{in}}^{E_N}$ and $\mathcal{H}_{\text{out}}^{E} := \mathcal{H}_{\text{out}}^{E_1} \otimes \mathcal{H}_{\text{out}}^{E_2} \otimes ... \otimes \mathcal{H}_{\text{out}}^{E_N}$, with j labelling the jth subsystem of E, and state $|\psi\rangle_E \in \mathcal{H}_{\text{in}}^{E}$ such that, for all states $\rho \in \mathcal{D}(\bigotimes_{j=1}^{N} \mathcal{H}_{\text{in}}^{j})$,

$$\Lambda_{1...N}[\rho] = \text{tr}_E\{V (\rho \otimes |\psi\rangle \langle\psi|_E) V^{\dagger}\},$$

where $V = \otimes_{k=1}^{N} U_{kE_k}$ for unitary operators $U_{kE_k}: \mathcal{H}_{\text{in}}^{k} \otimes \mathcal{H}_{\text{in}}^{E_k} \to \mathcal{H}_{\text{out}}^{k} \otimes \mathcal{H}_{\text{out}}^{E_k}$.

In addition to the above unitary representation, there exists another equivalent representation. This representation does not make reference to a tensor product structure in the ancilla, instead the unitaries in a unitary representation of a causal channel are independent of each other, in a particular sense. Now we have a global ancilla living in Hilbert space $\mathcal{H}_E$ and local ancillae E_k for each kth party, with input and output Hilbert

spaces $\mathcal{H}_{\text{in}}^{E_k}$ and $\mathcal{H}_{\text{out}}^{E_k}$, respectively. The local ancillae are introduced so that everything can remain unitary. Therefore, the total input and output Hilbert spaces of all the ancillae are $\mathcal{H}_E \otimes \mathcal{H}_{\text{in}}^{E_1} \otimes \mathcal{H}_{\text{in}}^{E_2} \otimes \ldots \otimes \mathcal{H}_{\text{in}}^{E_N}$ and $\mathcal{H}_E \otimes \mathcal{H}_{\text{out}}^{E_1} \otimes \mathcal{H}_{\text{out}}^{E_2} \otimes \ldots \otimes \mathcal{H}_{\text{out}}^{E_N}$, respectively.

Definition 42 (Commuting unitary representation of localizable channels). A causal channel $\Lambda_{1\ldots N}$ is localizable if and only if there exists a global ancilla system E with Hilbert space $\mathcal{H}_E$, a local ancilla system E_k for each kth party, with input and output Hilbert spaces $\mathcal{H}_{\text{in}}^{E_k}$ and $\mathcal{H}_{\text{out}}^{E_k}$, respectively, and state $|\psi\rangle_E \in \mathcal{H}_E \otimes \mathcal{H}_{\text{in}}^{E_1} \otimes \mathcal{H}_{\text{in}}^{E_2} \otimes \ldots \otimes \mathcal{H}_{\text{in}}^{E_N}$ such that, for all states $\rho \in \mathcal{D}(\bigotimes_{j=1}^N \mathcal{H}_{\text{in}}^j)$,

$$\Lambda_{1\ldots N}[\rho] = \text{tr}_{EE_1\ldots E_N}\left\{ \prod_{j=1}^N U_{jE} \left(\rho \otimes |\psi\rangle\langle\psi|_E\right) \prod_{k=0}^{N-1} U_{N-kE}^\dagger \right\},$$

where $U_{kE}: \mathcal{H}_{\text{in}}^k \otimes \mathcal{H}_E \otimes \mathcal{H}_{\text{in}}^{E_k} \to \mathcal{H}_{\text{out}}^k \otimes \mathcal{H}_E \otimes \mathcal{H}_{\text{out}}^{E_k}$ is a unitary operator for all k, such that, for any permutation π on the set $\{1, 2, \ldots, N\}$, $\prod_{k=1}^N U_{kE} = \prod_{k=1}^N U_{\pi(k)E}$.

Since all the Hilbert spaces in this work are taken to be finite dimensional, these two unitary representations of localizable channels are equivalent. This can be shown by a straightforward extension of lemma 4.1 in [48]. It should be remarked upon that if we were to allow for infinite dimensional Hilbert spaces, then these two definitions will not be equivalent, as pointed out by Cleve, Liu and Paulsen [49]. In full generality, since the first unitary representation implies the commuting unitary representation, one could then take the commuting unitary representation to be the most general definition of localizable channels when allowing for infinite dimension Hilbert spaces.

Finally, from the point-of-view of non-locality and steering, the set of *local* channels is of interest.

Definition 43 (Local channel). A channel is local if it is localizable, but with the additional constraint that the ancilla state σ_R is a separable state, i.e.

$$\sigma_R = \sum_\lambda p_\lambda |\phi_\lambda^1\rangle\langle\phi_\lambda^1|_{R_1} \otimes |\phi_\lambda^2\rangle\langle\phi_\lambda^2|_{R_2} \otimes \ldots \otimes |\phi_\lambda^N\rangle\langle\phi_\lambda^N|_{R_N},$$

for $|\phi_\lambda^j\rangle_{R_j} \in \mathcal{H}_{R_j}$.

It can be readily seen that localizable channels are more general than local channels. For example, the latter set of channels cannot be used to generate entanglement between two parties, but a localizable channel can. We can now summarize all the information about causal channels in the following theorem.

Theorem 44. *Let* C, $\tilde{\mathsf{Q}}$, Q, *and* L *be the set of causal, almost localizable, localizable, and local channels, respectively, then we have that* $\mathsf{C} \supsetneq \tilde{\mathsf{Q}} \supsetneq \mathsf{Q} \supsetneq \mathsf{L}$.

A natural question is given a channel Λ, can we decide if it belongs to C, $\tilde{\mathsf{Q}}$, Q, or L. We first restrict to the bipartite setting, and we suppose that Λ is given a convenient representation, such as the Choi–Jamiołkowski representation [50]. That is, for a bipartite setting channel Λ, the *Choi state* $\Omega \in \mathcal{D}(\mathcal{H}_{\text{in}}^1 \otimes \mathcal{H}_{\text{out}}^1 \otimes \mathcal{H}_{\text{in}}^2 \otimes \mathcal{H}_{\text{out}}^2)$, such that $\Omega := (\Lambda \otimes \mathbb{1}_{\mathcal{H}_{\text{in}}^1} \otimes \mathbb{1}_{\mathcal{H}_{\text{in}}^2})(|\Phi^+\rangle\langle\Phi^+|)$, for

$$|\Phi^+\rangle = \frac{1}{\sqrt{d_{\mathcal{H}_{\text{in}}^1} d_{\mathcal{H}_{\text{in}}^2}}} \sum_j |j\rangle_{\mathcal{H}_{\text{in}}^1} |j\rangle_{\mathcal{H}_{\text{in}}^1} \sum_k |k\rangle_{\mathcal{H}_{\text{in}}^2} |k\rangle_{\mathcal{H}_{\text{in}}^2},$$

with $d_{\mathcal{H}}$ being the dimension of the Hilbert space $\mathcal{H}$. The Choi state Ω is positive semi-definite if only if Λ is a channel. In order to decide if the channel Λ is causal, we can use the definition of causal channels along with this property of Choi states to state the following result.

Proposition 45. *A channel* Λ *is causal if and only if its Choi state* Ω *satisfies the following:*

1. *$\exists$ a density matrix* $\Sigma_1 \in \mathcal{D}(\mathcal{H}_{\text{in}}^1 \otimes \mathcal{H}_{\text{out}}^1)$ *such that* $\text{tr}_{\mathcal{H}_{\text{out}}^2} \Omega = \frac{1}{d_{\mathcal{H}_{\text{in}}^2}} \Sigma_1 \otimes \mathbb{1}_{\mathcal{H}_{\text{in}}^2}$, *and*

2. *$\exists$ a density matrix* $\Sigma_2 \in \mathcal{D}(\mathcal{H}_{\text{in}}^2 \otimes \mathcal{H}_{\text{out}}^2)$ *such that* $\text{tr}_{\mathcal{H}_{\text{out}}^1} \Omega = \frac{1}{d_{\mathcal{H}_{\text{in}}^1}} \mathbb{1}_{\mathcal{H}_{\text{in}}^1} \otimes \Sigma_2$.

A consequence of this result is that there exist positive semi-definite matrices Σ_1 and Σ_2 such that the conditions of the proposition are satisfied. In other words, one can decide whether a channel is in the set C using a semi-definite programme (see [51]), and so one can efficiently decide this problem.

Deciding whether a channel belongs in the sets of Q and L is not as easily resolved as the case for C. For example, while channels in L will have a Choi state Ω that is not entangled across the partition of party 1 and

party 2's respective Hilbert spaces, the converse is not true. That is, there are channels in C (but not in Q) whose Choi state is also not entangled across this partition [22]. Furthermore, deciding membership in L (even up to some error) is NP-hard, although it is possible to find conditions to test whether a channel is in the set L [51].

For case of deciding membership in Q (even up to some error), this is problem is also NP-hard [51]. In addition to this, the set of localizable channels is not closed [34]. Gutoski has given a criterion for deciding if a channel is in Q, which is somewhat analogous to the condition of complete positivity for channels. However, in general, there is no known way of deciding in finite time whether this criterion is satisfied. Indeed, as we will point out, this problem is deeply related to the problem of deciding whether certain correlations in a Bell test can be realized by local measurements on a quantum state; a problem with deep connections to open problems in mathematics [52]. For the case of deciding if a channel belongs to the set $\tilde{Q}$, we leave this to future work.

Appendix B. Bell non-locality

A traditional Bell experiment (sometimes called a 'Bell scenario') consists of N distant parties, each of them having access to a share of a physical system. These parties input (in a space-like separated manner) classical data into their device (labelled as $x_i \in \{1, ..., m\}$ for party i), and obtain outputs (labelled as $a_i \in \{1, ..., d\}$ for party i) from the device. For simplicity, in a bipartite setting (i.e. for $N = 2$), we will use the notation of inputs being x and y instead of x_1 and x_2, with outputs being a and b instead of a_1 and a_2.

The objects of interest in these Bell experiments are the correlations observed in the generated classical data, i.e. the conditional probability distribution $p(a_1, ..., a_N|x_1, ..., x_N)$. Depending on the type of device that the parties use (i.e. classical, quantum, possibly post-quantum), different correlations may be feasible in the experiment. The sets of correlations that have been of main interest in the literature are the following.

Definition 46. Classical correlations, also referred to as 'locally causal' [1], are those allowing for shared random variables $\lambda \in \Lambda$, and take the form

$$p(a_1, ..., a_N|x_1, ..., x_N) = \sum_{\lambda \in \Lambda} D_\lambda^1(a_1|x_1) \ ... \ D_\lambda^N(a_N|x_N) \, p(\lambda), \tag{12}$$

where $D_\lambda^j(\cdot) \in \{0, 1\}$ is a deterministic response function given λ for the jth party, and $p(\lambda)$ is the distribution over the variables λ such that $\sum_\lambda p(\lambda) = 1$.

Definition 47. Quantum correlations, arise if there exists a Hilbert space $\mathcal{H}$, a state $|\psi\rangle \in \mathcal{H}$, and (complete) projective measurements $\{\Pi_{a_i|x_i}^{(i)}\}_{a_i, x_i}$ for each party i, such that the conditional probability distribution is given by the Born rule:

$$p(a_1, ..., a_N|x_1, ..., x_N) = \langle\psi|\Pi_{a_1|x_1}^{(1)} \ ... \ \Pi_{a_N|x_N}^{(N)}|\psi\rangle, \tag{13}$$

and such that $\prod_{i=1}^N \Pi_{a_i|x_i}^{(i)} = \prod_{j=1}^N \Pi_{a_{\pi(j)}|x_{\pi(j)}}^{(\pi(j))}$, for any permutation π of the parties $\{1, 2, ..., N\}$.

Definition 48. Given a bipartition $S_1 \cup S_2 = \{1, ..., N\}$ of N parties $S_1 = \{i_1, ..., i_s\}$ and $S_2 = \{i_{s+1}, ..., i_N\}$, a conditional probability distribution $p(a_1, ..., a_N|x_1, ..., x_N)$ does not permit signalling across this bipartition if

$$p(a_{i_1}, ..., a_{i_s}|x_{i_1}, ..., x_{i_s}) = \sum_{a_{i_{s+1}}, ..., a_{i_N}} p(a_1, ..., a_N|x_1, ..., x_N) \tag{14}$$

$$p(a_{i_{s+1}}, ..., a_{i_N}|x_{i_{s+1}}, ..., x_{i_N}) = \sum_{a_{i_1}, ..., a_{i_s}} p(a_1, ..., a_N|x_1, ..., x_N), \tag{15}$$

for all inputs $(x_1, ..., x_N)$. A distribution $p(a_1, ..., a_N|x_1, ..., x_N)$ belongs to the set of non-signalling correlations if and only if it does not permit signalling across all bipartitions of the N parties.

There exist non-signalling correlations that do not have a quantum realization [17]. A relevant set of post-quantum yet non-signalling correlations is that of the *almost quantum correlations* [25]. It is notable that, as mentioned, the almost quantum correlations happen to comply with the physical information-theoretic principles that have been proposed so far to characterize the quantum set [19]. We now present the definition of the set of almost quantum correlations.

Definition 49. Almost quantum correlations, arise if there exists a Hilbert space $\mathcal{H}$, a state $|\psi\rangle \in \mathcal{H}$, and (complete) projective measurements $\{\Pi_{a_i|x_i}^{(i)}\}_{a_i, x_i}$ for each party i, such that the conditional probability distribution is given by:

$$p(a_1, ..., a_N|x_1, ..., x_N) = \langle\psi|\Pi_{a_1|x_1}^{(1)} \ ... \ \Pi_{a_N|x_N}^{(N)}|\psi\rangle, \tag{16}$$

such that $\prod_{i=1}^N \Pi_{a_i|x_i}^{(i)}|\psi\rangle = \prod_{j=1}^N \Pi_{a_{\pi(j)}|x_{\pi(j)}}^{(\pi(j))}|\psi\rangle$, for any permutation π of the N parties.

Appendix C. Einstein–Podolsky–Rosen steering

In analogy with the study of non-locality, in steering scenarios there are four sets of assemblages of particular interest [18].

Definition 50. LHS assemblages (a.k.a. *unsteerable* assemblages) are those that take the form:

$$\sigma_{a_1 \dots a_N | x_1 \dots x_N} = \sum_\lambda \prod_{j=1}^N D_\lambda(a_j | x_j) \sigma_\lambda,$$

where λ is a shared random variable, $D_\lambda(a_j | x_j) \in \{0, 1\}$ a deterministic response function for the jth party, and $\sigma_\lambda \in \mathcal{L}(\mathcal{H}_B)$ a subnormalised quantum state prepared by Bob as a function of λ, such that $\sum_\lambda \mathrm{tr}\{\sigma_\lambda\} = 1$.

Definition 51. Quantum assemblages arise when all N untrusted parties perform local measurements on a shared (possibly entangled) quantum state that is also shared with Bob. That is, the elements of the assemblage are

$$\sigma_{a_1 \dots a_N | x_1 \dots x_N} = \mathrm{tr}_{A_1 \dots A_N}[(\Pi_{a_1|x_1} \otimes \Pi_{a_2|x_2} \otimes \dots \otimes \Pi_{a_N|x_N} \otimes \mathbf{1}_B) \rho], \tag{17}$$

where for each jth party $\sum_{a_j} \Pi_{a_j|x_j} = \mathbf{1}$ forms a complete projective measurement for each x_j, and $\rho \in \mathcal{D}(\mathcal{H}_1 \otimes \dots \otimes \mathcal{H}_N \otimes \mathcal{H}_B)$ is the state of the shared system between N parties (the jth party having Hilbert space $\mathcal{H}_j$, for each $j \leqslant N$) and Bob (with Hilbert space $\mathcal{H}_B$).

Definition 52. Given a bipartition $S_1 \cup S_2 = \{1, \dots, N\}$ of the N untrusted parties where $S_1 = \{i_1, \dots, i_s\}$ and $S_2 = \{i_{s+1}, \dots, i_N\}$, an assemblage $\{\sigma_{a_1 \dots a_N | x_1 \dots x_N}\}$ does not permit signalling across this bipartition if its elements satisfy

$$\sigma_{a_{i_1}, \dots, a_{i_s} | x_{i_1}, \dots, x_{i_s}} = \sum_{a_{i_{s+1}}, \dots, a_{i_N}} \sigma_{a_1, \dots, a_N | x_1, \dots, x_N} \tag{18}$$

$$\sigma_{a_{i_{s+1}}, \dots, a_{i_N} | x_{i_{s+1}}, \dots, x_{i_N}} = \sum_{a_{i_1}, \dots, a_{i_s}} \sigma_{a_1, \dots, a_N | x_1, \dots, x_N}, \tag{19}$$

for all inputs $(x_1, \dots, x_N)$. An assemblage $\{\sigma_{a_1, \dots, a_N | x_1, \dots, x_N}\}$ belongs to the set of non-signalling assemblages if and only if it does not permit signalling across any bipartition of the N untrusted parties, and

$$\rho_B = \sum_{a_1, \dots, a_N} \sigma_{a_1, \dots, a_N | x_1, \dots, x_N}$$

for all inputs $(x_1, \dots, x_N)$, where $\rho_B \in \mathcal{D}(\mathcal{H}_B)$ is Bob's reduced quantum state.

In complete analogy with the study of non-locality, we call *post-quantum assemblages* those assemblages that are non-signalling yet are not quantum, and *post-quantum steering* is the demonstration that an assemblage is post-quantum. Furthermore, one can now study more specific relaxations of the set of quantum assemblages; this not only allows us to generate post-quantum assemblages, but if an assemblage does not belong to a set that is a relaxation of the quantum set, it is definitely not quantum. A relevant set is that of *almost quantum assemblages* [18], inspired by almost quantum correlations, and defined in [18] in terms of a semi-definite programme (see [31]).

Before presenting the definition, we give a bit of simplifying notation. Given a subset $S \subseteq \{1 \dots N\}$, we define the strings $(\vec{a}_S | \vec{x}_S) := (a_j \dots a_k \dots a_l | x_j \dots x_k \dots x_l)$ such that $j, k, l \in S$, i.e. strings of outputs given inputs for parties in the subset S. The string of all N parties is denoted $(\vec{a} | \vec{x})$. To refer to particular inputs and outputs in the string $(\vec{a}_S | \vec{x}_S)$, if $j \in S$, then $[(a_S | x_S)]_j := (a_j | x_j)$, $[a_S]_j := a_j$, and $[x_S]_j := x_j$. We then take the set of such strings, or *words*, to be $W := \{(\vec{a}_S | \vec{x}_S)\}_S$ (which also includes the empty string $\varnothing$ for when S is the empty set). In addition to this notation we also define two words $(\vec{a}_S | \vec{x}_S)$ and $(\vec{a}'_{S'} | \vec{x}'_{S'})$, to be *orthogonal*, denoted as $(\vec{a}_S | \vec{x}_S) \perp (\vec{a}'_{S'} | \vec{x}'_{S'})$ if there is a $j \in S \cap S'$ such that $[x_S]_j = [x'_{S'}]_j$ but $[a_S]_j \neq [a'_{S'}]_j$. We can now present the definition of almost quantum assemblages [18].

Definition 53. An assemblage $\{\sigma_{a_1 \dots a_N | x_1 \dots x_N}\}$ is an almost quantum assemblage if for the set of words $W := \{(\vec{a}_S | \vec{x}_S): S \subseteq \{1 \dots N\}\}$ with $|W|$ elements, there exists a matrix Γ of dimension $|W| \times |W|$, whose elements are $d_B \times d_B$ matrices $\Gamma_{(\vec{a}_S | \vec{x}_S), (\vec{a}'_{S'} | \vec{x}'_{S'})}$ indexed by words $(\vec{a}_S | \vec{x}_S)$ and $(\vec{a}'_{S'} | \vec{x}'_{S'})$, such that

(i) $\Gamma \geqslant 0$,

(ii) $\Gamma_{(\vec{a}_S | \vec{x}_S), (\vec{a}'_{S'} | \vec{x}'_{S'})} = 0$ if $(\vec{a}_S | \vec{x}_S) \perp (\vec{a}'_{S'} | \vec{x}'_{S'})$,
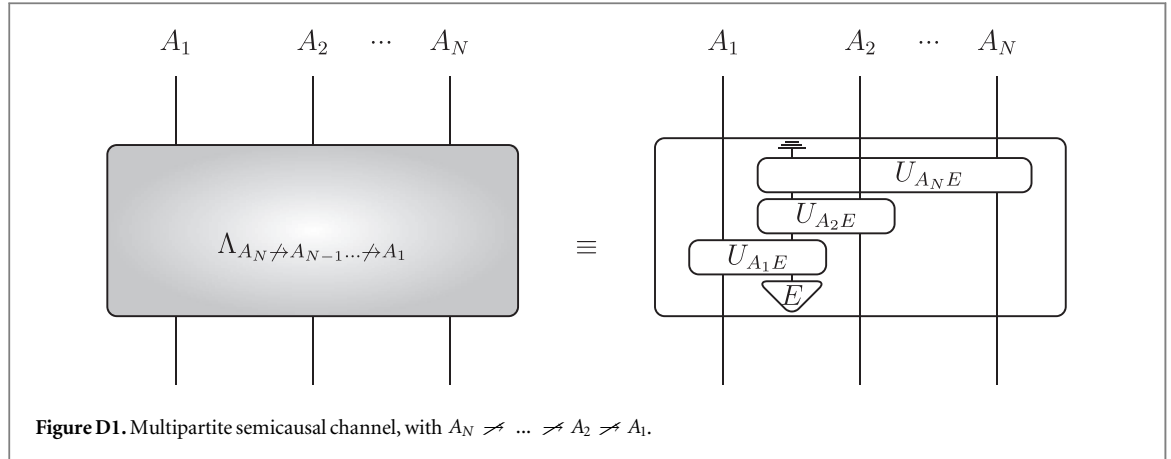
(iii) $\Gamma_{\varnothing, \varnothing} = \sigma_R$,

Figure D1. Multipartite semicausal channel, with $A_N \not\rightarrow ... \not\rightarrow A_2 \not\rightarrow A_1$.

(iv) $\Gamma_{\varnothing,(\vec{a}_S|\vec{x}_S)} = \sigma_{(\vec{a}_S|\vec{x}_S)}, \quad \forall \quad S,$

(v) $\Gamma_{(\vec{a}_T\vec{a}'_S|\vec{x}_T\vec{x}'_S),(\vec{a}_T\vec{a}''_{S'}|\vec{x}_T\vec{x}''_{S'})} = \Gamma_{(\vec{a}'_S|\vec{x}'_S),(\vec{a}_T\vec{a}''_{S'}|\vec{x}_T\vec{x}''_{S'})} = \Gamma_{(\vec{a}_T\vec{a}'_S|\vec{x}_T\vec{x}'_S),(\vec{a}''_{S'}|\vec{x}''_{S'})}$ for all subsets S, S', and T such that $T \cap S$ and $T \cap S'$ are both the empty set.

While this last definition is somewhat technical, section 3.1 presents an equivalent physical definition of the almost quantum assemblages.

Appendix D. Proofs of the statements in section 1 and appendix A

In this section we provide the proofs of the theorems and propositions of section 1 and appendix A.

D.1. Unitary representation of multipartite causal maps
In this section we characterize multipartite causal channels (see figure D1). Consider an N-partite map that acts on $\mathcal{H}_{in} = \otimes_{k=1}^{N} \mathcal{H}_{in}^{k}$, and denote by A_k the N input systems (a.k.a. parties). The main theorem that we will prove is the following.

Theorem 54 (Unitary representation for multipartite semicausal maps). *Let $\Lambda_{1...N}$ be a CP map, with $A_N \not\rightarrow ... \not\rightarrow A_2 \not\rightarrow A_1$. Then, there exist unitary operators U_{A_kE}: $\mathcal{H}_{in}^{k} \otimes \mathcal{H}_{int_k}^{E} \rightarrow \mathcal{H}_{out}^{k} \otimes \mathcal{H}_{int_{k+1}}^{E}$ for $2 \leqslant k \leqslant N-1$, U_{A_1E}: $\mathcal{H}_{in}^{1} \otimes \mathcal{H}_{in}^{E} \rightarrow \mathcal{H}_{out}^{1} \otimes \mathcal{H}_{int_1}^{E}$, and U_{A_NE}: $\mathcal{H}_{in}^{N} \otimes \mathcal{H}_{int_N}^{E} \rightarrow \mathcal{H}_{out}^{N} \otimes \mathcal{H}_{out}^{E}$, with $\mathcal{H}_{int_k}^{E}$ being the kth Hilbert space of the system E between unitaries $U_{A_{(k-1)}E}$ and U_{A_kE} acting on the system plus an ancilla E such that*

$$\Lambda_{1...N}[\rho] = \text{tr}_E\{U(\rho \otimes |0\rangle\langle 0|_E)U^\dagger\},$$

where $U = U_{A_NE}...U_{A_1E}$.

For bipartite maps this reduces to the result by [47]. To prove the multipartite statement, we need the following lemma:

Lemma 55. *Unitary representation for unitary semicausal maps [47]. Let Λ_{ABC} be a unitary tripartite CPTP map. If $C \not\rightarrow A$ the map can be decomposed as*

$$\Lambda_{ABC}[\rho_{ABC}] = U\rho_{ABC}U^\dagger,$$

with $U = U_{BC}\,U_{AB}$, where U_{BC} and U_{AB} are unitaries.

Now we can present the proof of the main theorem of the section.

Proof. Consider an N-partite map $\Lambda_{1...N}$ where $A_N \not\rightarrow A_{N-1} \not\rightarrow ... \not\rightarrow A_1$. By theorem 39, this channel can be decomposed as:

$$\Lambda_{1...N}[\rho] = \text{tr}_E\{U(\rho \otimes |0\rangle\langle 0|_E)U^\dagger\},$$

with $U = U_{A_NE}U_{A_{N-1}...A_1E}$.
Now, the situation is the following: we have $N-1$ input systems plus an ancilla E interacting via the unitary $U_{A_{N-1}...A_1E}$, and afterwards a unitary transformation U_{A_NE} is applied to the ancilla E and the last system A_N. The

first part of this protocol, by lemma 55, can be decomposed as a unitary between $A_1 \ldots A_{N-2}$ and E, followed by a unitary between A_{N-1} and E, i.e. $U_{A_{N-1} \ldots A_1 E} = U_{A_{N-1} E} U_{A_{N-2} \ldots A_1 E}$. By applying this way lemma 55 recursively, one gets $U_{A_{N-1} \ldots A_1 E} = U_{A_{N-1} E} U_{A_{N-2} E} \ldots U_{A_1 E}$. It follows that the channel $\Lambda_{1 \ldots N}$ has a unitary decomposition with $U = U_{A_N E} U_{A_{N-1} E} \ldots U_{A_1 E}$. □

A multipartite map is then causal if it is semicausal for all possible orderings of the parties.

Appendix E. Assemblages, correlations, distributed measurements, and teleportages, and the channels that define them

E.1. Proofs for correlations and assemblages

In this subsection we gives proofs of propositions 10, 11–13. Essentially the same proofs apply for the propositions 3, 4, 5, since one can take Bob's Hilbert space to be empty in a steering scenario, and for $N \geqslant 2$, we recover a Bell scenario.

Proposition 10. *An assemblage is a LHS assemblage if and only if there exists a local channel* $\Lambda_{1 \ldots N, B}^{\mathsf{L}}$: $\mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{\text{in}}}) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{\text{out}}})$ *such that the assemblage is channel-defined by* $\Lambda_{1 \ldots N, B}^{\mathsf{L}}$.

Proof. First we prove that if an assemblage has a LHS model, then it is channel-defined by a local channel. The first thing to note is that any LHS assemblage can be reproduced by the N untrusted parties making local measurements $\{M_{a_j|x_j} \in \mathcal{L}(\mathcal{H}_j)\}$ (for the jth party) on a separable state $\rho \in \bigotimes_{j=1}^N \mathcal{H}_j \otimes \mathcal{H}_B$, for $\mathcal{H}_j$ being the jth party's local Hilbert space. While some entangled states will only produce a LHS model assemblage, this means that there exists another separable state that can produce the same LHS model assemblage. Without loss of generality we can model this measurement as a projective measurement. This choice of local measurement $\{M_{a_j|x_j}\}$ can then be simulated by preparing the input choice as the state $|x_j\rangle \in \mathcal{H}_m$ where $x_j \in \{1, \ldots, m\}$. The outcome of the measurement will be translated into a register with Hilbert space $\mathcal{H}_d$, with outcomes $|a_j\rangle \in \mathcal{H}_m$ described by elements of an orthonormal basis; the register is initially prepared in the state $|0\rangle \in \mathcal{H}_d$. Therefore, each jth party is associated with the Hilbert space $\mathcal{H}_j \otimes \mathcal{H}_d \otimes \mathcal{H}_m$, and to this system we apply the unitary

$$U_j = \sum_{x_j, a_j} M_{a_j|x_j} \otimes |a_j\rangle \langle 0| \otimes |x_j\rangle \langle x_j|. \tag{20}$$

After applying this unitary, the systems in the register associated with $\mathcal{H}_j \otimes \mathcal{H}_m$ are traced out for each jth system, leaving the system in $\mathcal{H}_d$, which is then measured in the orthonormal basis $|a_j\rangle \in \mathcal{H}_m$. Therefore, the assemblage can be channel-defined by this whole local channel including the unitaries as the separable state ρ.

In the other direction, if given a local channel then it channel-defines a LHS assemblage. This should be clear since the local unitaries followed by a measurement acting on a share of a separable state, will only produce local measurements on a separable state, and therefore the assemblage has a LHS model. □

Proposition 11. *An assemblage* $\{\sigma_{a_1 \ldots a_N | x_1 \ldots x_N}\}$ *is non-signalling if and only if there exists a causal channel* $\Lambda_{1 \ldots N, B}^{\mathsf{C}} : \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{\text{in}}}) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{\text{out}}})$ *such that the assemblage is channel-defined by* $\Lambda_{1 \ldots N, B}^{\mathsf{C}}$.

Proof. First we prove that an assemblage that is channel-defined by a causal channel is a non-signalling assemblage. This follows immediately from the definition of causal channel. Given this channel-defined assemblage $\{\sigma_{a_1 \ldots a_N | x_1 \ldots x_N}\}$, when we take a sum over outcomes a_j, then this is equivalent to tracing out the output system of a causal channel. This thus results in a new assemblage that is channel-defined by a causal channel (with fewer output systems), and thus the initial assemblage is a non-signalling assemblage. It is also straightforward to see, given the definition of a causal channel, that when tracing out the N untrusted parties, we obtain a reduced quantum state for Bob that is independent of the inputs $(x_1, \ldots, x_N)$.

We now proceed to the converse statement that given a non-signalling assemblage, then there exists a causal channel such that the assemblage is channel-defined by it. First, if we fix orthonormal bases for the input and output Hilbert spaces as $\{|x_j\rangle\}$ and $\{|a_j\rangle\}$, respectively, then we construct the channel with Kraus decomposition

$$\Gamma(\cdot) = \sum_{\mathbf{a} := (a_1, \ldots, a_N), \mathbf{x} := (x_1, \ldots, x_N)} K_{\mathbf{a}, \mathbf{x}} (\text{tr}_B\{\cdot\} \otimes \sigma_{\mathbf{a}|\mathbf{x}}) K_{\mathbf{a}, \mathbf{x}}^\dagger, \tag{21}$$

such that $\sigma_{\mathbf{a}|\mathbf{x}} \in \mathcal{D}(\mathcal{H}_B)$, and

$$K_{\mathbf{a}, \mathbf{x}} = |a_1\rangle \langle x_1| \otimes |a_2\rangle \langle x_2| \otimes \ldots \otimes |a_N\rangle \langle x_N| \otimes \mathbf{1}_B. \tag{22}$$

It can be readily verified that this is a channel, and when we prepare $|x_1 \ldots x_N\rangle \langle x_1 \ldots x_N| \otimes |0\rangle \langle 0|_B$ as input, act on the input with Γ, and then measure in the basis $\{|a_1 \ldots a_N\rangle\}$ on the N untrusted parties, we obtain an assemblage.

It remains to show that channel Γ is itself a causal channel given a non-signalling assemblage. This can be shown inductively first tracing out the output system of party 1 as follows:

$$\text{tr}_1\{\Gamma(\cdot)\} = \sum_{a_1} \langle a_1 | \left(\sum_{\mathbf{a}:=(a_1, \ldots, a_N), \mathbf{x}:=(x_1, \ldots, x_N)} K_{\mathbf{a},\mathbf{x}}(\text{tr}_B\{\cdot\} \otimes \sigma_{\mathbf{a}|\mathbf{x}}) K_{\mathbf{a},\mathbf{x}}^\dagger \right) | a_1 \rangle \tag{23}$$

$$= \sum_{x_1} \langle x_1 | \left(\sum_{\mathbf{a}:=(a_2, \ldots, a_N), \mathbf{x}:=(x_2, \ldots, x_N)} K_{\mathbf{a},\mathbf{x}}(\text{tr}_B\{\cdot\} \otimes \sigma_{\mathbf{a}|\mathbf{x}}) K_{\mathbf{a},\mathbf{x}}^\dagger \right) | x_1 \rangle \tag{24}$$

$$= \Gamma'(\text{tr}_1\{\cdot\}), \tag{25}$$

where $\Gamma' : \mathcal{L}(\mathcal{H}_m^{N-1} \otimes \mathcal{H}_B) \to \mathcal{L}(\mathcal{H}_d^{N-1} \otimes \mathcal{H}_B)$ is another channel corresponding to parties 2 to N. The second line above results from the fact that the assemblages are non-signalling, and new channel Γ' is written as

$$\Gamma'(\cdot) = \sum_{\mathbf{a}:=(a_2, \ldots, a_N), \mathbf{x}:=(x_2, \ldots, x_N)} L_{\mathbf{a},\mathbf{x}}(\text{tr}_B\{\cdot\} \otimes \sigma_{\mathbf{a}|\mathbf{x}}) L_{\mathbf{a},\mathbf{x}}^\dagger, \tag{26}$$

for

$$L_{\mathbf{a},\mathbf{x}} = |a_2\rangle \langle x_2| \otimes |a_3\rangle \langle x_3| \otimes \ldots \otimes |a_N\rangle \langle x_N| \otimes \mathbf{1}_B. \tag{27}$$

The same argument works for any party j, and then given the new channel Γ', one can trace out one or more of the remaining parties' outputs to get another channel, and so on. In this way, the channel Γ channel-defines the non-signalling assemblage, and is causal, thus concluding the proof. $\qquad\square$

Proposition 13. *An assemblage* $\{\sigma_{a_1 \ldots a_N | x_1 \ldots x_N}\}$ *is quantum if and only if there exists a localizable channel* $\Lambda_{1 \ldots N, B}^{\mathsf{Q}} : \mathcal{L}(\mathcal{H}_m^{\otimes N} \otimes \mathcal{H}_{B_{\text{in}}}) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{B_{\text{out}}})$ *such that the assemblage is channel-defined by* $\Lambda_{1 \ldots N, B}^{\mathsf{Q}}$.

Proof. The proof of this is exactly the same as the proof for proposition 23, except the separable state in the proof is replaced with an entangled state. $\qquad\square$

As mentioned above, the proofs above easily generalize to the study of correlations. Indeed, one can run through the above arguments and just have Bob's system be the empty system, thus recovering the Bell scenario for $N \geqslant 2$.

E.2. Proofs for distributed measurements and teleportages

In this subsection we gives proofs of propositions 32–34. Essentially the same proofs apply for the propositions 23–25 since, as with the connection between steering and Bell scenarios, one can take Bob's Hilbert space to be empty in a non-classical teleportation scenario, and for $N \geqslant 2$ we recover a Buscemi non-locality scenario.

Proposition 32. *A teleportage is local if and only if there exists a local channel* $\Lambda_{1 \ldots N}^{\mathsf{L}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j \otimes \mathcal{K}_B) \to \mathcal{L}(\bigotimes_{j=1}^n \mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ *such that the teleportage is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{L}}$.

Proof. The proof that a teleportage is channel-defined by a local channel is a local teleportage is immediate from the definitions, i.e. a local channel sequentially combined with a local measurement is again a local measurement. For the converse statement that given a local teleportage, there exists a local channel that channel-defines the teleportage, the channel is constructed by having local unitaries that 'copy' the outcome of a local measurement to a local output register (with Hilbert space $\mathcal{H}_d$) into an orthonormal basis of this register, which is then measured in this basis. $\qquad\square$

Proposition 33. *A teleportage is quantum if and only if there exists a localizable channel* $\Lambda_{1 \ldots N}^{\mathsf{Q}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j \otimes \mathcal{K}_B) \to \mathcal{L}(\bigotimes_{j=1}^n \mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ *such that the teleportage is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{Q}}$.

Proof. The proof of this follows the proof of proposition 44, except now with an entangled state instead of a separable state. $\qquad\square$

Proposition 34. *A teleportage is non-signalling if and only if there exists a causal channel* $\Lambda_{1 \ldots N}^{\mathsf{C}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\text{in}}^j = \bigotimes_j \mathcal{K}_j \otimes \mathcal{K}_B) \to \mathcal{L}(\bigotimes_{j=1}^n \mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ *such that the teleportage is channel-defined by* $\Lambda_{1 \ldots N}^{\mathsf{C}}$.

Proof. Given a teleportage that is channel-defined by a causal channel, the teleportage is non-signalling essentially by definition: taking a sum over outcomes a_j is equivalent to tracing out the jth output system of the channel, resulting in a new teleportage for all systems not including j. For the other direction, of given a non-signalling teleportage, we can construct a causal channel that channel-defines the teleportage. First, given the elements $T_{a_1, \ldots, a_N}$ of the teleportage, since it is forms an instrument in general, we can straightforwardly construct a channel from an instrument: we introduce output registers $\mathcal{H}_d$ for each jth party and thus define a channel $\Gamma : \mathcal{L}(\bigotimes_j \mathcal{K}_j \otimes \mathcal{H}_B) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_B)$ in the following way:

$$\Gamma(\cdot) = \sum_{\mathbf{a} := (a_1, \ldots, a_N)} |\mathbf{a}\rangle \langle \mathbf{a}| \otimes T_\mathbf{a}(\mathrm{tr}_B\{\cdot\}), \tag{28}$$

with $|\mathbf{a}\rangle \in \mathcal{H}_d^{\otimes N}$. Now with measurements on the register of the outputs $\mathcal{H}_d^{\otimes N}$, the teleportage is then channel-defined by Γ. It remains to show that Γ is causal. This follows from the definition of non-signalling teleportages. $\qquad \square$

E.3. Constructing channels from correlations and assemblages

Let us begin with proposition 8's statement and proof.

Proposition 8. *Given* $\Lambda_{1 \ldots N}^c(\cdot)$ *from* $p(a_1 \ldots a_N | x_1 \ldots x_N)$, *for all measurements* $M_{a_j'}$, *and all states* $\rho_{x_j'}$, *the correlations*

$$p(a_1', a_2', \ldots, a_N' | x_1', x_2', \ldots, x_N') = \mathrm{tr}\left\{ \bigotimes_{j=1}^N M_{a_j'} \left(\Lambda_{1 \ldots N}^c \otimes \mathbf{1}_{\mathrm{aux}} \left(\bigotimes_{j=1}^N \rho_{x_j'} \right) \right) \right\} \tag{29}$$

are local if the correlations $p(a_1 \ldots a_N | x_1 \ldots x_N)$ *are local.*

Proof. Given a distribution $p(a_1' \ldots a_N' | x_1' \ldots x_N')$ of the form in equation (5) with $\Lambda_{1 \ldots N}(\cdot)$ being the channel in (4), we first observe that this probability is invariant if for each state $\rho_{x_j'} \in \mathcal{D}(\mathcal{H}_{\mathrm{in}}^j \otimes \mathcal{H}_{\mathrm{aux}}^j)$, we replace it with $\rho_{x_j'}' = \sum_{x_j} (\mathbf{1}_{\mathrm{aux}}^j \otimes |x_j\rangle \langle x_j|) \rho_{x_j'} (\mathbf{1}_{\mathrm{aux}}^j \otimes |x_j\rangle \langle x_j|)$, where $\mathbf{1}_{\mathrm{aux}}^j$ is the identity acting on $\mathcal{H}_{\mathrm{aux}}^j$. Therefore, for whatever state $\rho_{x_j'}$, the probabilities $p(a_1' \ldots a_N' | x_1' \ldots x_N')$ are unchanged if we replace it with a state $\rho_{x_j'}' = \sum_{x_j} \sigma_{x_j, x_j'} \otimes |x_j\rangle \langle x_j|$, where $\sigma_{x_j, x_j'} \in \mathcal{D}(\mathcal{H}_{\mathrm{aux}}^j)$ and is equal to $\mathrm{tr}_{\mathrm{in}}(\rho_{x_j'} \mathbf{1}_{\mathrm{aux}}^j \otimes |x_j\rangle \langle x_j|)$. Furthermore, we apply a similar argument to the general measurement $M_{a_j'}$, such that the distribution is conserved by replacing it with

$$M_{a_j'}' = \sum_{a_j} Q_{a_j, a_j'} \otimes |a_j\rangle \langle a_j|, \tag{30}$$

such that $Q_{a_j, a_j'} = \mathrm{tr}_{\mathrm{in}}((\mathbf{1}_{\mathrm{aux}}^j \otimes \langle a_j|) M_{a_j'})$, which lives in $\mathcal{D}(\mathcal{H}_{\mathrm{aux}}^j)$. Taking $\rho_{a_j'}'$ and $\{M_{a_j'}'\}_j$ as our state and measurements, we obtain the correlations

$$p(a_1', a_2', \ldots, a_N' | x_1', x_2', \ldots, x_N') = \sum_{a_1 \ldots a_N} \sum_{x_1 \ldots x_N} p(a_1 \ldots a_N | x_1 \ldots x_N) \mathrm{tr}\left\{ \bigotimes_{j=1}^N Q_{a_j, a_j'} \sigma_{x_j, x_j'} \right\} \tag{31}$$

$$= \sum_{a_1 \ldots a_N} \sum_{x_1 \ldots x_N} p(a_1 \ldots a_N | x_1 \ldots x_N) \prod_{j=1}^N q(a_{j'} | a_j, x_j, x_j'). \tag{32}$$

Hence, the correlations $p(a_1', a_2', \ldots, a_N' | x_1', x_2', \ldots, x_N')$ can be seen as the correlations $p(a_1 \ldots a_N | x_1 \ldots x_N)$ acted on by some local stochastic processing in the form of a conditional probability $q(a_{j'} | a_j, x_j, x_j')$. Since local stochastic processing cannot turn local correlations into non-local ones, the distribution $\{p(a_1', a_2', \ldots, a_N' | x_1', x_2', \ldots, x_N')\}$ will be local provided that $\{p(a_1 \ldots a_N | x_1 \ldots x_N)\}$ is. $\qquad \square$

We now present proposition 17, and its proof.

Proposition 17. *Given* $\Sigma_{1 \ldots N, B}^c(\cdot)$ *from assemblage elements* $\sigma_{a_1 \ldots a_N | x_1 \ldots x_N}$, *this channel is local-limited if for all measurements* $P_{a_B | x_B} \in \mathcal{L}(\mathcal{H}_B)$ *indexed by the choice* x_B *and outcomes* a_B, *the correlations* $p(a_1, \ldots, a_N, a_B | x_1, \ldots, x_N, x_B) := \mathrm{tr}\{P_{a_B | x_B} \sigma_{a_1 \ldots a_N | x_1 \ldots x_N}\}$ *are local.*

Proof. Firstly, since the channel $\Sigma_{1 \ldots N, B}^c$ traces out the part of $\rho_{x_B'}$ that is input into the channel, without loss of generality, we can replace $\rho_{x_B'}$ with $\sigma_0 \otimes \sigma_{x_B'}$, where $\sigma_0 \in \mathcal{D}(\mathcal{H}_{B_{\mathrm{in}}})$ is some fixed state, and $\sigma_{x_B'} \in \mathcal{D}(\mathcal{H}_{B_{\mathrm{aux}}})$. Therefore, the preparation of $\sigma_{x_B'}$ following by the measurement $M_{a_B'}$ can be incorporated into a single measurement $\{P_{a_B' | x_B'} \in \mathcal{L}(\mathcal{H}_{B_{\mathrm{out}}})\}$, such that $\sum_{a_B'} P_{a_B' | x_B'} = \mathbf{1}_{B_{\mathrm{out}}}$. This then simplifies the correlations to be

$$p(a_1', \ldots, a_N', a_B' | x_1', \ldots, x_N', x_B') = \mathrm{tr}\left\{ \bigotimes_{j=1}^{N} M_{a_j'} \otimes P_{a_B'|x_B'} \left(\Sigma_{1\ldots N,B}^c \otimes \mathbb{1}_{\mathrm{aux}} \left(\bigotimes_{j=1}^{N} \rho_{x_j'} \otimes \rho_0 \right) \right) \right\}. \tag{33}$$

By expanding out this expression and using identical reasoning to the proof of proposition 8, we arrive at

$$p(a_1', \ldots, a_N', a_B' | x_1', \ldots, x_N', x_B') = \sum_{a_1 \ldots a_N} \sum_{x_1 \ldots x_N} \mathrm{tr}\{ P_{a_B'|x_B'} \sigma_{a_1 \ldots a_N | x_1 \ldots x_N} \} \prod_{j=1}^{N} q(a_{j'} | a_j, x_j, x_j'), \tag{34}$$

where $q(a_{j'} | a_j, x_j, x_j')$ is a local conditional probability. Therefore, if $\mathrm{tr}\{ P_{a_B'|x_B'} \sigma_{a_1 \ldots a_N | x_1 \ldots x_N} \}$ is local for all measurements $P_{a_B'|x_B'}$, the correlations are local, and the channel is local-limited. $\qquad \square$

Appendix F. Unitary representation of non-signalling assemblages and teleportages

In this section we discuss the unitary representations of non-signalling assemblages and teleportages as outlined in theorems 12 and 35. As outlined in the main text, the GHJW theorem [53, 54] can also be seen as a corollary of theorem 12, and our generalization of the GHJW theorem is a corollary of theorem 35. We will only present the proof of theorem 35 since theorem 12 is a special case.

Theorem 35 Unitary representation of non-signalling teleportages. *Let $\{ T_{a_1 \ldots a_N} \}$ be a non-signalling teleportage. Then, the teleportage is channel-defined by a channel $\Lambda_{1 \ldots N,B}^{\mathsf{C}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\mathrm{in}}^j \otimes \mathcal{K}_B) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ if and only if there exist*

- *auxiliary systems E and E' with input and output Hilbert spaces, $\mathcal{H}_{\mathrm{in}}^E$ and $\mathcal{H}_{\mathrm{out}}^E$ for E, with $\mathcal{H}_{\mathrm{in}}^E = \mathcal{K}_B$ and $\mathcal{H}_{\mathrm{out}}^{E'} = \mathcal{K}_B$ for E', that is the output Hilbert space of E' and B coincide;*

- *quantum state $|R\rangle \in \mathcal{H}_{\mathrm{in}}^E \otimes \mathcal{H}_{\mathrm{in}}^{E'}$;*

- *unitary operator $V : \bigotimes_j \mathcal{H}_{\mathrm{in}}^j \otimes \mathcal{H}_{\mathrm{in}}^E \to \mathcal{H}_d^{\otimes N} \otimes \mathcal{H}_{\mathrm{out}}^E$,*

which produce a unitary representation of the channel $\Lambda_{1 \ldots N,B}^{\mathsf{C}} : \mathcal{L}(\bigotimes_j \mathcal{H}_{\mathrm{in}}^j \otimes \mathcal{K}_B) \to \mathcal{L}(\mathcal{H}_d^{\otimes N} \otimes \mathcal{K}_B)$ via

$$\Lambda_{1 \ldots N,B}^{\mathsf{C}}(\cdot) = \mathrm{tr}_{E_{\mathrm{out}}} \{ V(\mathrm{tr}_{B_{\mathrm{in}}}(\cdot) \otimes |R\rangle \langle R|_{E,E'}) V^\dagger \}.$$

Furthermore, the unitary V can be decomposed into a sequence of unitaries $U_{k,E} : \mathcal{H}_m \otimes \mathcal{H}_1^E \to \mathcal{H}_d \otimes \mathcal{H}_2^E$ for appropriately chosen Hilbert spaces $\mathcal{H}_1^E$ and $\mathcal{H}_2^E$, where for any given permutation π of the se $\{1, \ldots, N\}$, we have that

$$V = U_{\pi(1),E}^\pi U_{\pi(2),E}^\pi \cdots U_{\pi(N),E}^\pi,$$

where $U_{k,E}^\pi$ is not necessarily the same as $U_{k,E}^{\pi'}$ for two different permutations π and π'.

Proof. First of all, we can treat all N untrusted parties as a single party called 'Alice' (or A for short) that produces an N-length string $\mathbf{a} := (a_1, \ldots, a_N)$. This treatment of N untrusted parties, and 1 trusted party is summarised in figure D2, where the N parties are treated as a single Alice. This allows us to then consider a bipartite non-classical teleportation scenario and a causal map Λ_{AB}. First, write the map in its unitary representation with $A \nrightarrow B$, this first equivalence then gives part of the statement of the theorem where V is the unitary V_{AE}, and the state on $\mathcal{K}_B$ is the state of the right-hand system of $|R\rangle$. The second equivalence shows that the whole channel can be seen as a localizable channel across the bipartition between Alice and Bob.
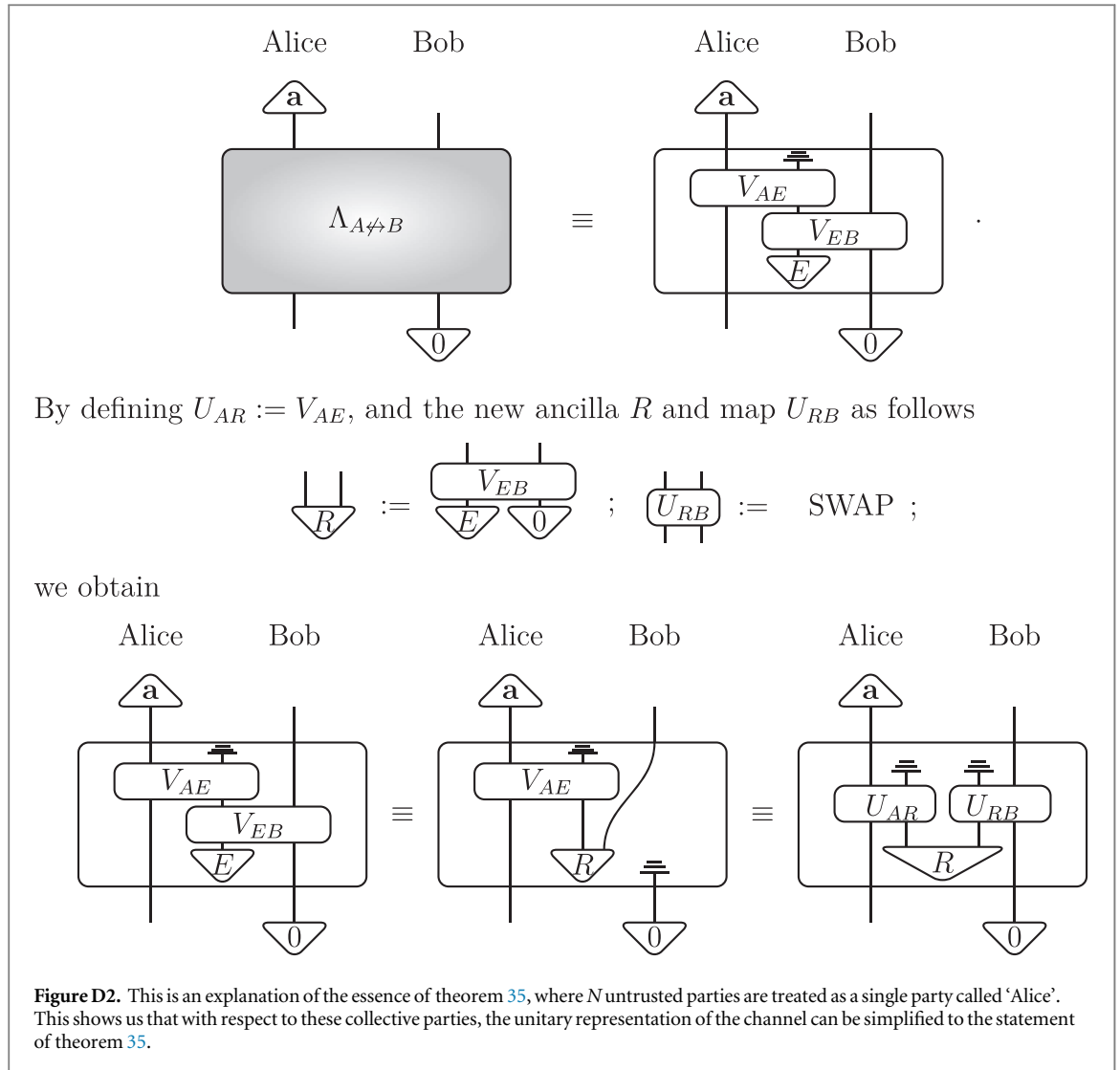
In order to prove the remainder of the theorem, it remains to decompose V into a sequence of unitaries. To do this, we use the fact that among the N parties (that make up Alice) there cannot be any signalling. Therefore, using theorem 54 we arrive at the full statement of the theorem. $\qquad \square$

Appendix G. Almost quantum assemblages

Here we provide the proofs of lemma 16 and theorem 15, respectively.

Lemma. *An assemblage $\{ \sigma_{a_1 \ldots a_N | x_1 \ldots x_N} \}$ is almost quantum if and only if there exists a Hilbert space $\mathcal{H} \cong \mathcal{K} \otimes \mathcal{H}_B$, quantum state $|\psi\rangle \in \mathcal{H}$, and projective measurements $\{ \Pi_{a_j|x_j} \in \mathcal{L}(\mathcal{K}) \}$ for each jth party where $\sum_{a_j} \Pi_{a_j|x_j} = \mathbb{1}$ and for all permutations π of $\{1, \ldots, N\}$, $\prod_{j=1}^{N} \Pi_{a_{\pi(j)}|x_{\pi(j)}} |\psi\rangle = \prod_{j=1}^{N} \Pi_{a_j|x_j} |\psi\rangle$, such that*

$$\sigma_{a_1 \ldots a_N | x_1 \ldots x_N} = \mathrm{tr}_{\mathcal{K}} \left\{ \prod_{j=1}^{N} \Pi_{a_j|x_j} \otimes \mathbb{1}_B |\psi\rangle \langle\psi| \right\}. \tag{35}$$

Figure D2. This is an explanation of the essence of theorem 35, where *N* untrusted parties are treated as a single party called 'Alice'. This shows us that with respect to these collective parties, the unitary representation of the channel can be simplified to the statement of theorem 35.

Proof. Consider a steering scenario where *N* parties steer one, by performing *m* measurements of *d* outcomes each. Let *D* be the dimension of the Hilbert space of the characterized party, denoted by $\mathcal{H}_B$.

We will first prove the 'if' direction.

Take an assemblage $\{\sigma_{\vec{a}|\vec{x}}\}$, generated by the uncharacterized parties performing measurements $\Pi^{(k)}_{a_k|x_k}$ on the state ρ, which without loss of generality we can consider to be a pure state $\rho = |\psi\rangle\langle\psi|$. Let *W* be the set of words as in definition 53. Now define the matrix Γ, with size $|W| \times |W|$ and whose elements are $D \times D$ matrices, as follows:

$$\Gamma_{(\vec{a}_S|\vec{x}_S),(\vec{a}'_{S'}|\vec{x}'_{S'})} := \mathrm{tr}_{1, ..., N}\left\{ \prod_{i=1:|S|} \Pi^\dagger_{a_i|x_i} \prod_{j=1:|S'|} \Pi_{a'_j|x'_j} \rho \right\},$$

$$\Gamma_{\varnothing,(\vec{a}'_{S'}|\vec{x}'_{S'})} := \mathrm{tr}_{1, ..., N}\left\{ \prod_{j=1:size(x')} \Pi_{a'_j|x'_j} \rho \right\},$$

$$\Gamma_{(\vec{a}_S|\vec{x}_S),\varnothing} := \mathrm{tr}_{1, ..., N}\left\{ \prod_{i=1:size(x)} \Pi^\dagger_{a_i|x_i} \rho \right\},$$

$$\Gamma_{\varnothing,\varnothing} := \sigma_R.$$

By definition, this Γ satisfies (iii) and (iv) of definition 53. Since the measurements $\Pi_{a_k|x_k}$ are projective, Γ also satisfies condition (ii). The commutation relation of the projective measurements $\Pi_{a_k|x_k}$ on the state ρ further imply (v).

We still need to show that $\Gamma \geqslant 0$. We hence need to show that Γ as an element of $M_{|W|}(M_D)$ (i.e. of the set of $|W| \times |W|$ matrices whose entries are $D \times D$ matrices) is positive semi-definite, which is equivalent to showing that Γ as an element of $M_{|W|D}$ (i.e. of the set of $|W|D \times |W|D$ matrices whose entries are complex numbers) is

positive semi-definite [55]. We will hence show that $\Gamma \in M_{|W|D}$ is a Gramian matrix (i.e. it can be written as $\Gamma = V^\dagger V$ for some matrix V), and since all Gramian matrices are positive semi-definite the claim follows.

First observe that entries of Γ are of the form

$$\Gamma_{[(\vec{a}_S|\vec{x}_S)]_j,[(\vec{a}_{S'}'|\vec{x}_{S'}')]_l} = \langle j| \, \Gamma_{(\vec{a}_S|\vec{x}_S),(\vec{a}_{S'}'|\vec{x}_{S'}')} |l \rangle,$$

where $[(\vec{a}_S|\vec{x}_S)]_j$, $[(\vec{a}_{S'}'|\vec{x}_{S'}')]_l$ denotes the (j, l) component of the matrix $\Gamma_{(\vec{a}_S|\vec{x}_S),(\vec{a}_{S'}'|\vec{x}_{S'}')}$, with $\Gamma_{(\vec{a}_S|\vec{x}_S),(\vec{a}_{S'}'|\vec{x}_{S'}')} \in \{\Gamma_{u,v}: \, u, v \in W\}$, and $|j\rangle$, $|l\rangle$ elements of an orthonormal basis of $\mathcal{H}_B$.

By cyclicity of the partial trace we can also write

$$\Gamma_{(\vec{a}_S|\vec{x}_S),(\vec{a}_{S'}'|\vec{x}_{S'}')} = \mathrm{tr}_{1\ldots N}(F_{(\vec{a}_{S'}'|\vec{x}_{S'}')}|\psi\rangle\langle\psi| \, G^\dagger_{(\vec{a}_S|\vec{x}_S)})$$

for $F_{(\vec{a}_{S'}'|\vec{x}_{S'}')}$, $G_{(\vec{a}_S|\vec{x}_S)} \in \{\mathbf{1}_A\} \cup \{\Pi_v: \, v \in S\}$, where we have identified $\Pi_{(\vec{a}_S|\vec{x}_S)} := \prod_{i=1:|S|} \Pi_{a_i|x_i}$.

By defining $\mathcal{H}_A$ as the Hilbert space of the uncharacterized parties, note that

$$\langle j| \, \Gamma_{(\vec{a}_S|\vec{x}_S),(\vec{a}_{S'}'|\vec{x}_{S'}')} |l\rangle = \sum_{|y\rangle \in \mathcal{H}_A} \langle j|\langle y| F_{(\vec{a}_{S'}'|\vec{x}_{S'}')}|\psi\rangle \langle\psi| \, G^\dagger_{(\vec{a}_S|\vec{x}_S)} |y\rangle |l\rangle$$

$$= \sum_{|y\rangle \in \mathcal{H}_A} \langle\psi| \, G^\dagger_{(\vec{a}_S|\vec{x}_S)}|y\rangle|l\rangle \langle j|\langle y| F_{(\vec{a}_{S'}'|\vec{x}_{S'}')}|\psi\rangle$$

$$= \left(\sum_{|y'\rangle \in \mathcal{H}_A} \langle\psi| \, G^\dagger_{(\vec{a}_S|\vec{x}_S)}|y'\rangle|l\rangle \langle y'|\right)\left(\sum_{|y\rangle \in \mathcal{H}_A} \langle j|\langle y| F_{(\vec{a}_{S'}'|\vec{x}_{S'}')}|\psi\rangle|y\rangle\right)$$

$$= \sum_{y'} \alpha^*_{y',[(\vec{a}_S|\vec{x}_S)]_l} \langle y'| \sum_y \alpha_{y,[(\vec{a}_{S'}'|\vec{x}_{S'}')]_j}|y\rangle$$

$$= \langle u|_{[(\vec{a}_S|\vec{x}_S)]_l}|v\rangle_{[(\vec{a}_{S'}'|\vec{x}_{S'}')]_j},$$

where $\{|y\rangle\}$ is an orthonormal basis in $\mathcal{H}_A$ such that $\langle y'|y\rangle = \delta^y_{y'}$ and $\alpha_{y,[(\vec{a}_{S'}'|\vec{x}_{S'}')]_j} = \langle j|\langle y| F(\vec{a}_{S'}'|\vec{x}_{S'}')|\psi\rangle$ is some scalar. Now we can further define the set of vectors

$$\left\{ |v\rangle_{[(\vec{a}_{S'}'|\vec{x}_{S'}')]_j}: \, (\vec{a}_{S'}'|\vec{x}_{S'}') \in W, \, j = 1 \ldots D, \, |v\rangle_{[(\vec{a}_{S'}'|\vec{x}_{S'}')]_j} = \sum_y \alpha_{y,[(\vec{a}_{S'}'|\vec{x}_{S'}')]_j}|y\rangle \right\},$$

and define the matrix V as that whose columns are each one of these vectors. We see then that $\Gamma = V^\dagger V$, i.e. the elements of Γ are all the inner product of vectors associated with a row and column. Therefore, Γ is Gramian.

This proves the first part of the claim: an assemblage that arises from performing those types of measurements on a quantum state are almost quantum assemblages.

For the converse, take an almost quantum assemblage $\{\sigma_{\vec{a}|\vec{x}}\}$. Let $\Gamma \in M_{|W|}(M_D)$ be its moment matrix from definition 53. Let $\{|i\rangle: \, i \in W\}$ be the Gram vectors of that matrix, i.e. $\Gamma_{i,j} = \langle i|j\rangle$. Note that these vectors have as entries elements of M_D (i.e. $D \times D$ matrices on complex numbers). For each uncharacterized party k, define the subspace $V^{(k)}_{a_k|x_k} := \mathrm{span}\{|(a_k\vec{a}'|x_k\vec{x}')\rangle: \, (\vec{a}'|\vec{x}') \in S_{\bar{k}}\}$, where $S_{\bar{k}} \subset W$ is the subset of words that do not involve party k. Note for clarity we have dropped the index indicating the subset of the parties. Now define:

$$E^{(k)}_{a_k|x_k} := \mathrm{proj}(V^{(k)}_{a_k|x_k}), \quad \forall \, x = 1:m, \quad \forall \, a = 1:d-1, \tag{36}$$

$$E^{(k)}_{d|x_k} := \mathbf{1} - \sum_{a=1}^{d-1} E^{(k)}_{a_k|x_k}. \tag{37}$$

By definition, $E^{(k)}_{a_k|x_k}$ are projection operators. Condition (ii) of definition 53 implies that $E^{(k)}_{a_k|x_k} E^{(k)}_{a_k'|x_k} = 0$ if $a \neq a'$. Hence, $\{E^{(k)}_{a_k|x_k}\}_a$ defines a complete projective measurement for each x_k for each party.

Now we will see the action of any sequence $\prod_{k=1}^N E^{(k)}_{a_k|x_k}$ on $|\varnothing\rangle$. Let us start with just one projector:

$$E^{(k)}_{a_k|x_k}|\varnothing\rangle = E^{(k)}_{a_k|x_k}|(a_k|x_k)\rangle + E^{(k)}_{a_k|x_k}(|\varnothing\rangle - |(a_k|x_k)\rangle)$$

$$= |(a_k|x_k)\rangle$$

since by definition $E^{(k)}_{a_k|x_k}|(a_k|x_k)\rangle = |(a_k|x_k)\rangle$ and condition (v) implies that

$$\langle (a_k\vec{a}'|x_k\vec{x}'), \varnothing\rangle = \langle (a_k\vec{a}'|x_k\vec{x}'), (a_k|x_k)\rangle$$

for all $(\vec{a}'|\vec{x}') \in S_{\bar{k}}$.

The same reasoning can be applied to $E^{(k)}_{a|x}$ acting on an arbitrary $|v\rangle$ with $v \in S_{\bar{k}}$:

$$E^{(k)}_{a_k|x_k}|v\rangle = E^{(k)}_{a_k|x_k}|(a_k|x_k) \, v\rangle + E^{(k)}_{a_k|x_k}(|v\rangle - |(a_k|x_k) \, v\rangle)$$

$$= |(a_k|x_k) \, v\rangle \tag{38}$$

since $|(a_k|x_k) \, v\rangle \in V^{(k)}_{a_k|x_k}$ and condition (v) implies $\langle (a_k|x_k) \, v, v\rangle = \langle (a_k|x_k) \, v, (a_k|x_k) \, v\rangle$.

Therefore,

$$\sigma_{\vec{a}_S|\vec{x}_S} = \langle\varnothing|\prod_{k=1:n}E^{(k)}_{a_k|x_k}|\varnothing\rangle = \mathrm{tr}\left\{\prod_{k=1:n}E^{(k)}_{a_k|x_k}|\varnothing\rangle\langle\varnothing|\right\}. \tag{39}$$

This motivates the following definition: $\rho := |\varnothing\rangle\langle\varnothing|$.

The next ingredient is to check that commutation relations of the projective measurements $\{E^{(k)}_{a_k|x_k}\}$ on the state ρ satisfy the conditions stated in the Lemma. This follows immediately from (38).

To conclude the proof, notice that the trace in equation (39) can actually be interpreted as a partial trace on the Hilbert space of the uncharacterized parties $\mathcal{H}_A$. This follows from understanding the set $M_{|S|}(M_D)$ as the tensor product algebra $M_{|S|}\otimes M_D$ [55]. First, invoke the isomorphism between projection operators $E \in M_{|S|}(M_D)$ and projection operators $E\otimes\mathbf{1}_D$ with $E\in M_{|S|}$ [13]. Then, notice that an orthonormal basis $\{|\phi\rangle\}$ for $M_{|S|}(M_D)$ can be seen as $\{|\phi\rangle\langle\phi| = |\varphi\rangle\langle\varphi|\otimes\mathbf{1}_D\}$, where $\{|\varphi\rangle\}$ in an orthonormal basis for $\mathcal{H}_A$. □

Theorem. *An assemblage* $\{\sigma_{a_1\ldots a_N|x_1\ldots x_N}\}$ *is almost quantum if and only if there exists an almost localizable channel* $\Lambda^{\tilde{Q}}_{1\ldots N,B} : \mathcal{L}(\mathcal{H}_m^{\otimes N}\otimes\mathcal{H}_{B_{in}}) \to \mathcal{L}(\mathcal{H}_d^{\otimes N}\otimes\mathcal{H}_{B_{out}})$ *such that the assemblage is channel-defined by* $\Lambda^{\tilde{Q}}_{1\ldots N,B}$.

Proof. Given an assemblage channel-defined by an almost localizable channel, it is immediate that is an almost quantum assemblage due to lemma 29. To show that an assemblage as defined in lemma 29 can be channel-defined by an almost localizable channel, we can use exactly the same constructive argument as in proposition 17. That is, given projectors as in lemma 29, we can construct local unitaries that act on a register in the state $|\psi\rangle$ as in the proof of proposition 17. □

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[13] For the reader who is not familiar with operator algebras, the main idea of the identification goes as follows. Let E be a projection operator onto a linear subspace of a vector space of dimension $|S|$ on the space of matrices M_D. E is hence an element of $M_{|S|}(M_D)$. Now perform a change of basis of $M_{|S|}(M_D)$ so that E is diagonal. That is, in this new basis $\{|\phi\rangle\}$, either $E|\phi\rangle = |\phi\rangle$ or $E|\phi\rangle = 0$. Remember that these $|\phi\rangle$ are vectors with entries in M_D. Hence, $\langle\phi|E|\phi\rangle = \mathbf{1}_D$ or $\langle\phi|E|\phi\rangle = \mathbb{O}_D$. It follows that we can think of E as $E = E_{\mathcal{H}_A}\otimes\mathbf{1}_D$, where $E_{\mathcal{H}_A}$ is a projector in $M_{|S|}$.

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